

An integrated environment for mathematical public transport optimization

Documentation

Version 2024.12

Currently developed at

Department of Mathematics and Systems Analysis
Aalto University, Espoo

and

Optimization Research Group
RPTU Kaiserslautern-Landau

and

Fraunhofer ITWM
Kaiserslautern

Originally developed at

Institute for Numerical and Applied Mathematics
University of Göttingen

Contributors

Head

- Philine Schiewe (current head)
- Anita Schöbel (former head and founder)

Technical Lead

- Sven Jäger

Researchers

- Sebastian Albert
- Christine Biedinger
- Thorsten Dahlheimer
- Vera Grafe
- Olli Herrala
- Klara Hoffmann
- Sarah Roth
- Alexander Schiewe
- Moritz Stinzendörfer
- Reena Urban

Student Assistants

- Ricardo Reicherz

Former Researchers

- Urs Baumgart
- Rasmus Fuhse
- Konstantinos Gkoumas
- Marc Goerigk
- Jonas Harbering
- Jonas Ide
- Julius Pätzold
- Michael Schachtebeck
- Jochen Schulz
- Michael Siebert
- Felix Spühler
- Anke Uffmann

Former Student Assistants

- Florentin Hildebrandt
- Jonas Hürter
- Jarkko Jalovaara
- David Kaiser
- Eero Ketola
- Benjamin Lieser
- Kim Reece
- Michael Rihlmann
- Leevi Rönty
- Mridul Roy
- Lisa Sandig
- Christopher Scholl
- Linda Sieber
- Vitali Telezki

Contents

1	Introduction	7
1.1	What is LINTIM?	7
1.2	Installation and Requirements	8
1.2.1	Connecting LINTIM with a solver	8
1.3	Installation Script	10
1.4	Typical Usage: A Hands-On-Example	10
2	Overview on the Planning Steps	14
2.1	Stop Location	15
2.2	Line Pool Generation	16
2.3	Line Planning	16
2.4	Ridepool Generation	18
2.5	Ride Concept	19
2.6	Periodic Timetabling	20
2.7	Tariff Planning	21
2.8	Vehicle Scheduling	22
2.9	Delay Management	25
2.10	Integrated Planning	27
3	Detailed Description of Algorithms	29
3.1	Stop Location	29
3.1.1	Without a given infrastructure network	29
3.1.2	For a given infrastructure network	30
3.2	Line Pool Generation	30
3.2.1	Creating a new line pool with the tree based heuristic	31
3.2.2	Creating a line pool while restricting the duration of the lines	32
3.2.3	Creating a line pool by k shortest paths	33
3.2.4	Terminal-to-terminal	33
3.2.5	Center-Periphery	33
3.2.6	Line costs	35
3.3	Line Planning	35
3.3.1	Cost	36
3.3.2	Direct	37
3.3.3	Cost direct weighted sum	39
3.3.4	Traveling time without frequencies	39
3.3.5	Traveling time with frequencies	40
3.3.6	Cost with traveling time bound	41
3.3.7	Minchanges	41
3.3.8	Game	42
3.4	Ridepool Generation	43
3.4.1	All	43
3.4.2	Demand heuristic	43

3.4.3	Tree based	44
3.4.4	Adapt Line Pool	44
3.5	Ride Concept	44
3.5.1	Lineplanning and ridepooling with fixed demand factors	44
3.5.2	Lineplanning and ridepooling including demand factors as variables	45
3.6	Timetabling	45
3.6.1	Modulo network simplex algorithms	45
3.6.2	Constraint propagation	47
3.6.3	Abscon	47
3.6.4	MATCH	47
3.6.5	PESP-IP	47
3.6.6	Cycle-based IP	48
3.6.7	Phase 1 simplex	48
3.6.8	Adaptions	48
3.7	Tariff Planning	48
3.7.1	General Remarks	49
3.7.2	Flat Tariff	50
3.7.3	Distance Tariffs	51
3.7.4	Zones	51
3.8	Vehicle Scheduling	54
3.8.1	Mdm1	54
3.8.2	Mdm2	54
3.8.3	Assignment model	54
3.8.4	Transportation model	54
3.8.5	Network flow model	55
3.8.6	Canal model	55
3.8.7	Line-based	55
3.8.8	Simple	55
3.8.9	IP model	55
3.9	Delay Management	56
3.9.1	Propagate	56
3.9.2	Integer-Linear-Programming based methods	56
3.10	Integrated Planning	58
3.10.1	Integrated timetabling and passenger routing	59
3.10.2	Integrated timetabling and aperiodic vehicle scheduling	59
3.10.3	Integrated line planning and timetabling	60
3.10.4	Integrated line planning, timetabling and vehicle scheduling	60
3.10.5	Robust Timetabling and Vehicle Scheduling Using Machine Learning	61
3.10.6	Eigenmodel	63
4	Auxiliary Algorithms	65
4.1	Dataset Generation	65
4.1.1	Input	65
4.1.2	Output	65
4.1.3	Algorithms	65
4.2	OD Matrix Creation	67
4.2.1	Input	67
4.2.2	Output	67
4.2.3	Algorithms	68
4.2.4	Distribute from node demand	68
4.3	Load distribution	68
4.3.1	Input	69
4.3.2	Output	69
4.3.3	Algorithms	69

4.3.4	Using the EAN	70
4.3.5	Using spanner MIPs	71
4.4	Headway creation	72
4.4.1	Input	72
4.4.2	Output	72
4.4.3	Algorithm	72
4.5	PTN to EAN	72
4.5.1	Input	72
4.5.2	Output	72
4.5.3	Algorithm	72
4.6	EAN buffer activities	74
4.7	EAN reroute passengers	75
4.8	Tariff (Reference) Price Matrix	75
4.8.1	Input	75
4.8.2	Output	76
4.8.3	Algorithms	76
4.9	Rollout	77
4.9.1	Input	78
4.9.2	Output	78
4.9.3	Algorithm	78
4.9.4	Requirements and caveats	79
4.9.5	Generating trips	79
4.10	Delay generation	80
4.11	Visualization	81
4.11.1	PTN	81
4.11.2	OD	82
4.11.3	Loads	83
4.11.4	Line pool	84
4.11.5	Line concept	84
4.11.6	Timetable	85
4.11.7	Disposition timetable	85
4.11.8	Tariff	85
4.11.9	mapgui	87
4.12	Interaction with VISUM	87
4.12.1	Writing files for VISUM	87
4.12.2	Reading a config file	87
4.12.3	Reading the infrastructure	88
4.12.4	Reading the PTN	89
4.12.5	Reading the demand	89
4.12.6	Reading stops and lines	90
4.12.7	Reading a timetable	90
4.12.8	Reading fixed lines	91
4.12.9	Reading fixed times	91
5	Evaluation	93
5.1	Evaluation of the PTN Created by Stop Location	93
5.2	Evaluation of the PTN	94
5.3	Evaluation of the OD Matrix	94
5.4	Evaluation of the Line Pool	94
5.5	Evaluation of the Line Concept	95
5.6	Evaluation of the EAN	97
5.7	Evaluation of the Timetable	98
5.7.1	Capacitated Routing	99
5.8	Evaluation of the Tariff created by Tariff Planning	100

5.9	Evaluation of the Trips	101
5.10	Evaluation of the Vehicle Schedules	101
5.11	Evaluation of the Disposition Timetable	102
6	Overview of Supported Integer Programming Solvers	104
7	Configuration Parameters	106
7.1	General	106
7.2	Stop Location	106
7.3	OD	107
7.4	PTN	107
7.5	Line Planning	107
7.6	Ridepooling	108
7.7	Load Generation	109
7.8	Load Visualization	110
7.9	Periodic EAN	111
7.10	Debug	112
7.11	Timetabling	113
7.12	Tariff Planning	113
7.13	Vehicle Scheduling	115
7.14	Delay Management	115
7.15	Dataset Generation	116
7.16	Integrated Models	116
7.16.1	General	116
7.16.2	LinTimPass	117
7.16.3	LinTimPassVeh	117
7.16.4	TimPass	117
7.16.5	TimVeh	117
7.16.6	TimVehToLin	117
7.17	TimPassLib	117
8	In- and Output Data	118
8.1	Config	118
8.2	Statistic	119
8.3	Basis	119
8.3.1	Additional load	119
8.3.2	Change station	120
8.3.3	Demand	120
8.3.4	Demand geo	120
8.3.5	Edge	120
8.3.6	Edge forbidden	121
8.3.7	Edge infrastructure	121
8.3.8	Edge infrastructure forbidden	121
8.3.9	Edge walking	122
8.3.10	Existing stop	122
8.3.11	Existing stop geo	122
8.3.12	Existing edge	123
8.3.13	Headway	123
8.3.14	Load	123
8.3.15	Node	123
8.3.16	OD	124
8.3.17	OD node	124
8.3.18	Pool	124
8.3.19	Pool cost	124

8.3.20	Reference Price Matrix	124
8.3.21	Restricted turns	125
8.3.22	Restricted turns infrastructure	125
8.3.23	Ridepool	125
8.3.24	Routings	125
8.3.25	Station limits	126
8.3.26	Stop	126
8.3.27	Stop geo	126
8.3.28	Terminals	126
8.4	Line Planning	126
8.4.1	Line concept	127
8.4.2	Fixed lines	127
8.4.3	Line capacities	127
8.4.4	Rideconcept	127
8.5	Timetabling	127
8.5.1	Activities periodic	128
8.5.2	Events periodic	128
8.5.3	Fixed times	128
8.5.4	Initial duration assumptions	129
8.5.5	Timetable periodic	129
8.5.6	Timetable for VISUM	129
8.6	Tariff Planning	129
8.6.1	Price Matrix	129
8.6.2	Zones	130
8.6.3	Zone Prices	130
8.7	Vehicle Scheduling	130
8.7.1	Vehicle schedules	130
8.8	Delay Management	130
8.8.1	Events expanded	131
8.8.2	Activities expanded	131
8.8.3	Timetable expanded	131
8.8.4	Timetable disposition	132
8.8.5	Delays events	132
8.8.6	Delays activities	132
8.8.7	Trips	132
8.9	GTFS	133
9	Datasets	134
9.1	Configuration Parameters for Datasets	134
9.2	Artificial Datasets	135
9.2.1	Toy	135
9.2.2	Grid	135
9.2.3	Ring	135
9.3	Datasets based on real world data	135
9.3.1	Sioux Falls	135
9.3.2	Lowersaxony	137
9.3.3	Goevb	137
9.3.4	Athens	138
9.3.5	Bahn-01	138
9.3.6	Bahn-02	138
9.3.7	Bahn-03	140
9.3.8	Bahn-04	140
9.3.9	Bahn-equal-frequencies	140
9.3.10	BOMHarbour	142

9.3.11	Mandl	142
9.4	Adding new datasets	142
9.4.1	Adding a dataset from PESPlib	143
9.4.2	Adding a dataset from TimPassLib	143
9.4.3	Dataset generator	145
10	LinTim Core	146
10.1	Model	146
10.1.1	Interfaces	146
10.1.2	Classes	146
10.1.3	Enumerations	147
10.1.4	Package model.impl	148
10.2	Input and Output	148
10.3	Algorithm	148
10.4	Utility	148
10.5	Solver	148
10.6	Exceptions	149
11	Introduction to extending LinTim	153
11.1	Logging	153
11.1.1	Output from LinTim programs	153
11.1.2	Log levels	153
11.1.3	Error messages	153
11.1.4	Info messages	154
11.2	Cleaning	154
12	Continuous Integration	155
12.1	Running the tests	155
12.2	Adding test cases	155
13	Changelog	156

Chapter 1

Introduction

1.1 What is LINTIM?

LINTIM is an academic algorithm and dataset library for mathematical public transport optimization. Problems in public transport optimization range from finding suitable locations for stations over calculating passenger-friendly timetables to handling unexpected delays. As it would be too complicated (though best in theory) to handle all these problems at the same time, they are split up and solved sequentially.

However, what seems to be best for one particular problem may have devastating influence on a different problem: For example a good timetable might not be well suited for delay management. **LinTim** (standing for **Line**planning and **Tim**etabling) addresses this issue by integrating the various public transport optimization problems and algorithms into one single environment. It hence gives the possibility to go back and forth in the sequence of public transport optimization problems in order to find solutions that work well on a greater scope and not only for the respective problem.

The data files are based on simple plain text formats that allow the implementation of algorithms in whatever programming language the developer likes to use. Thus, it is made easy to extend the current LINTIM-library and keep up to date with new developments and ideas.

LINTIM is designed for the use in UNIX, and will not work flawlessly in a native Windows environment.

Throughout the documentation, we will use some markers to indicate what certain teletyped texts mean:

- `[Fo]` foldername (relative paths w.r.t. the current dataset),
- `[Fi]` filename (relative paths w.r.t. the current dataset),
- `[R]` command that can run in some shell,
- `[C]` config entry with key and value,
- `[CK]` config key,
- `[CV]` config value,
- `[S]` statistic entry with key and value,
- `[SK]` statistic key,
- `[SV]` statistic value.
- `[CK]([Fi])` a config key for a filename, followed by the default value

1.2 Installation and Requirements

LINTIM uses many different programming languages. For the most parts, it is enough to have Java (≥ 11 with `ant` $\geq 1.9.8$ and `maven` ≥ 4), C, C++ and Python3 (≥ 3.5) installed on your system. There may be some special algorithms requiring additional programming languages, but if this is the case this is noted in the respective section of the documentation.

Using Windows 10 The easiest way to run LINTIM under Windows 10 is using a WSL installation. For installation instructions, see <https://docs.microsoft.com/en-us/windows/wsl/install-win10>. Using the WSL you can follow the installation notes listed for Linux below.

Using macOS Although macOS is a Unix-based operating systems, some of the below mentioned installation commands need to be adjusted when using macOS. The most important difference is the unavailability of `apt-get` for package management. Please check the different installed packages for the best way to install on macOS but for most of them, there are easy installation procedures using Homebrew, see <https://brew.sh/>. With that, see the installation notes for Linux below for more information.

Using a Linux distribution In this section, we list the commands to install all dependencies available in the Debian GNU/Linux Package index using `apt-get`. If you use another package manager, you need to adapt the corresponding commands.

To install all package manager dependencies of LINTIM, run

```
[R] sudo apt-get install build-essential openjdk-11-jdk ant graphviz python3-pip
```

To install the python package dependencies using `pip`, run

```
[R] sudo pip3 install numpy networkx pulp holoviews weightedstats pandas matplotlib seaborn
```

Also for using all of LINTIM, you will have to fulfill other third-party dependencies. For more information, have a look at [\[Fi\]](#) `/libs/README.md`. For a list of supported integer programming solvers and how to connect them with LINTIM, see the next section.

1.2.1 Connecting LINTIM with a solver

Some programs make use of integer programming solvers like Xpress, CPLEX and Gurobi. However, for each of the planning stages line planning, timetabling and vehicle scheduling there are also algorithms working without a solver installed. See the instructions of the respective algorithms for configuring LINTIM to use your chosen solver and Chapter 6 for a general overview which methods support which solver. If you want to use an integer programming solver, make sure to install it using the corresponding documentation and to set the environment variables correctly. In the following, we give a short overview which environment variables need to be set for LINTIM to find the corresponding solver. We suggest adding the below code snippets to your `/.bashrc`-file (or your equivalent, depending on your used environment), for automatic environment variable setting.

Gurobi For Gurobi, the `CLASSPATH` and `LD_LIBRARY_PATH` variables need to be set. On your machine, this might mean to run

```
[R] export GUROBI_HOME=/opt/gurobi/linux64
```

```
[R] export CLASSPATH=${GUROBI_HOME}/lib/gurobi.jar:${CLASSPATH}
```

```
[R] export LD_LIBRARY_PATH=${GUROBI_HOME}/lib/:${LD_LIBRARY_PATH}
```

For more information, check the Gurobi documentation.

Xpress For Xpress, source the `xpvars.sh` script provided with the installation. On your machine, this might mean to run

```
R source /opt/xpressmp/bin/xpvars.sh
```

This will take care of setting the appropriate environment variables for Xpress. For more information, check the Xpress documentation.

CPLEX For CPLEX, the `PATH`, `CLASSPATH` and `LD_LIBRARY_PATH` variables need to be set. On your machine, this might mean to run

```
R export CPLEX_HOME=/opt/ibm/ILOG/CPLEX_Studio201/cplex
```

```
R export CLASSPATH=${CPLEX_HOME}/lib/cplex.jar:${CLASSPATH}
```

```
R export LD_LIBRARY_PATH=${CPLEX_HOME}/bin/x86-64_linux/${LD_LIBRARY_PATH}
```

```
R export PATH=${CPLEX_HOME}/bin/x86-64_linux:${PATH}
```

Additionally, make sure to run the python installation script provided with the CPLEX installation to install the CPLEX python package. On your machine, this might mean to run

```
R sudo python3 /opt/ibm/ILOG/CPLEX_Studio201/python/setup.py install
```

For more information, check the CPLEX documentation.

SCIP For SCIP, the `PATH` and `LD_LIBRARY_PATH` variables need to be set. On your machine, this might mean to run

```
R export SCIPDIR=/opt/scipoptsuite-7.0.2
```

```
R export LD_LIBRARY_PATH=${SCIPDIR}/build/lib/${LD_LIBRARY_PATH}
```

```
R export PATH=${SCIPDIR}/build/bin:${PATH}
```

If you want to use SCIP from a Java program, make sure to install JSCIPOpt as well, see <https://github.com/scipopt/JSCIPOpt>. After installing, extend the above environment variables with

```
R export JSCIPDIR=/opt/scipoptsuite-7.0.2
```

```
R export LD_LIBRARY_PATH=${JSCIPDIR}/build/Release:${LD_LIBRARY_PATH}
```

```
R export CLASSPATH=${JSCIPDIR}/build/Release/scip.jar:${CLASSPATH}
```

For more information, check the SCIP and JSCIPOpt documentation.

GLPK To use GLPK as a solver in `LINTIM`, only the binary `glpsol` needs to be in the `PATH`. You can install GLPK e.g. with

```
R sudo apt-get install glpk-utils
```

COIN and CBC The coin and cbc solver are both bundled with the PuLP python package. Therefore you don't need to install anything additionally here.

1.3 Installation Script

The installation script is a Python script which leads you through the most parts of the installation of LINTIM. By following the instructions of the script you install the required system dependencies, LINTIM, the LINTIM dependencies, Gurobi and the Python dependencies. If you want to use the installation script you have to start it from the shell by running

```
R python3 install.py
```

after downloading `Fi install.py` and `Fi util.py`. If you already downloaded LINTIM you can find the installation file in `Fo src/installation`. The installation script attempts to install the required python packages via pip. It should therefore be executed within the desired (virtual) environment. Note that certain installations require sudo access where you will be prompted for your password.

1.4 Typical Usage: A Hands-On-Example

In the following we describe the typical usage of LINTIM and give an overview over the structure of the repository.

Its root directory consists of the following:

- `Fo /ci`
Folder for continuous integration tests.
- `Fo /datasets`
The LINTIM instances and their customized configuration files.
- `Fo /doc`
All documents regarding the LINTIM project (e.g. this documentation).
- `Fo /libs`
A folder to place dependencies. If necessary, the dependency will be described in the corresponding algorithm section.
- `Fo /src`
The source code of the LINTIM algorithms.

In `Fo /datasets` you can see all the datasets which are implemented in LINTIM for the time being. For further information on these datasets see Chapter 9, including information on how to add your own datasets to LINTIM.

Our goal in this example will be to calculate a disposition timetable for the “toy”-dataset and describe several of the in- and output files that you can find during the process. Note that in general, LINTIM provides the capability to configure all file paths. For simplicity, we will only provide the default values for this config keys in this chapter. For more information, see the following chapters.

Change into the folder

```
Fo /datasets/toy
```

in order to run algorithms on the “toy”-dataset. You find an exemplary folder-structure of a dataset folder:

- `Fo basis`
Contains all the data describing the instance like OD matrix, edges, loads, line pool, headways, etc.
- `Fo delay-management`
Will contains all the data related to delay-management and aperiodic planning.
- `Fo graphics`
Will contain all graphical output of the algorithms you might use.

- `Fo` `line-planning`
Will contains all the data related to line planning.
- `Fo` `statistic`
Will contain all output of evaluations you might run (may not exist yet, will be created automatically on evaluation).
- `Fo` `timetabling`
Will contain all data related to periodic timetabling.
- `Fo` `vehicle-scheduling`
Will contain all data related to vehicle scheduling.

As you can see, the folder names (and thus the contents) are related to the different steps of mathematical public transport optimization.

Every output you produce will by default be written into the respective folders.

This means, if you somehow produce an output regarding e.g. the delay-management, it will be written to `Fo` `delay-management`.

Also, each dataset folder contains a Makefile.

LinTim algorithms are used by calling make.

For instance typing

```
R make lc-line-concept
```

while being located in the “toy”-folder will compile all necessary files, calculate a line concept for the “toy”-instance and write it into `Fi` `line-planning/Line-concept.lin`.

Note that by default, this will use Xpress as an integer optimization solver. Therefore to successfully run this step, Xpress needs to be installed. See Chapter 1.2 for more information.

For calculating a line concept, LINTIM uses the data given in `Fo` `basis`.

Having a look into the makefile the line

```
line-concept :
${SRC_DIR}/line-planning/line-planning.sh ${FILENAME_CONFIG}
```

tells us, that the line concept is calculated using the algorithms from

`Fo` `/src/line-planning` with the configurations given in

`Fi` `${FILENAME_CONFIG}`, which is `Fi` `basis/Config.cnf` by default.

For detailed instructions on configuration files and how to change them see Section 8.1.

If you want to use different algorithms see Chapter 2 to know which are already implemented, Chapter 3 for detailed information on the implemented algorithms and Chapter 11 for instructions on how to implement your own into LINTIM.

So let’s have a look at what we got from our call

```
R make lc-line-concept
```

The file `Fi` `line-planning/Line-concept.lin` should contain something like this:

```
# line-id; edge-order; edge-id; frequency
1;1;1;0
1;2;6;0
1;3;7;0
2;1;2;3
2;2;6;3
...
```

LINTIM usually works with text files structured similarly (# comments a line). The advantage of this concept is that they are very independent of the programming language used.

In the most text files, like in this example, an explanation will be given on how to read them.

So now we got ourselves a first line concept for our “toy”-example. Next thing to do would be calculating a feasible timetable. For this we first have to provide an Event-Activity-Network (EAN). We can make LINTIM calculate this by calling

```
[R] make ean
```

Note that in order to calculate this EAN LINTIM of course needs a public transportation network (PTN), given by the network itself and a line concept on this network.

Of course it would be possible to design the algorithms in a way that a call of

```
[R] make ean
```

automatically generates a line concept if none is existent so far but for different reasons we refrained from this.

Therefore before calling

```
[R] make ean
```

you will always have to provide a line concept. Calling it before calculating a line concept will result in an error.

By calling

```
[R] make ean
```

we calculated the events and activities of our EAN. These are written to

```
[Fi] timetabling/Activities-periodic.giv and [Fi] timetabling/Events-periodic.giv.
```

For instance `[Fi] timetabling/Events-periodic.giv` should look something like this:

```
# event_id; type; stop-id; line-id; passengers; line-direction;
  line-freq-repetition
1; "departure"; 1; 1; 20; >; 1
2; "arrival"; 3; 1; 20; >; 1
3; "departure"; 3; 1; 20; >; 1
...
```

The first line again tells us how to read the file, i.e. e.g. event 2 is an arrival of line 1 at stop 3 carrying 20 passengers.

In order to calculate a timetable from this data we just call

```
[R] make tim-timetable
```

and LINTIM will write a timetable to `[Fo] timetabling/timetable-periodic.tim` in which you can look up the event given by its index and the time it is scheduled to take place.

Given this timetable we can now concentrate on the delay-management or the vehicle-scheduling.

We will try out the DM step first. This is a little bit more complex because there are some prerequisites we have to provide.

First of all we need an aperiodic timetable since the DM-algorithms only work for these.

But we do not really need a new aperiodic timetable. We just need our periodic timetable expanded that is we have to adhere the periods.

For LINTIM we call this "Rollout" and calculate it by calling

```
[R] make ro-rollout
```

The needed “aperiodic” timetable will be written to

`[Fi] delay-management/Timetable-expanded.tim` and will also be included in

`[Fi] delay-management/Events-expanded.tim`.

After calculating this timetable we can create some delays by calling

`[R] make dm-delays`

This will call the delay-generator which generates source delays for our given network. More on how the delay-generator works and how to control it can be found in Section 4.10.

After creating some delays we finally want to calculate a disposition timetable and do that by calling

`[R] make dm-disposition-timetable`

The timetable will be calculated and written to

`[Fi] delay-management/Timetable-disposition.tim`.

For concluding our first LINTIM-cycle we now want to calculate a vehicle scheduling.

For this we first have to consider, that all the trips that have to be completed by some vehicle have to be known. In a periodic timetable this might not be the case. Because of this we have to rollout the whole trips and we can do so by setting `[C] rollout_whole_trips` to true. Changing a config-parameter is done in the following way:

Change to

`[Fo] basis`

and write

```
rollout_whole_trips; true
```

into `[Fi] basis/Private-Config.cnf`.

Now for calculating the vehicle-schedules we first have to repeat the steps from and including

`[R] make ro-rollout`

We then have to calculate the trips, the vehicles have to do. We can do so by typing

`[R] make ro-trips`

and the trips will be written to `[Fi] delay-management/Trips.giv`.

Now calling

`[R] make vs-vehicle-schedules`

calculates the vehicle schedule and it is written to

`[Fi] vehicle-scheduling/Vehicle_Schedules.vs`.

In the end, we want to evaluate the created vehicle schedule. By running

`[R] make vs-vehicle-schedules-evaluate`

we evaluate the current vehicle schedule and the computed properties will be written to `[Fi] statistic/statistic.sta`, e.g. `[SK] vs_cost`, the cost of the vehicle schedule and `[SK] vs_feasible`, whether the computed schedule is feasible.

Beside this few make-targets we introduced there are a lot more in LINTIM. Have a look into the makefiles to see which possible targets exist. Which algorithm will be called exactly is defined by the configuration file. For a description of which parameter setting will call which algorithm, see Chapter 2.

Chapter 2

Overview on the Planning Steps

The different public transport optimization problems can be summarized in the following figure:

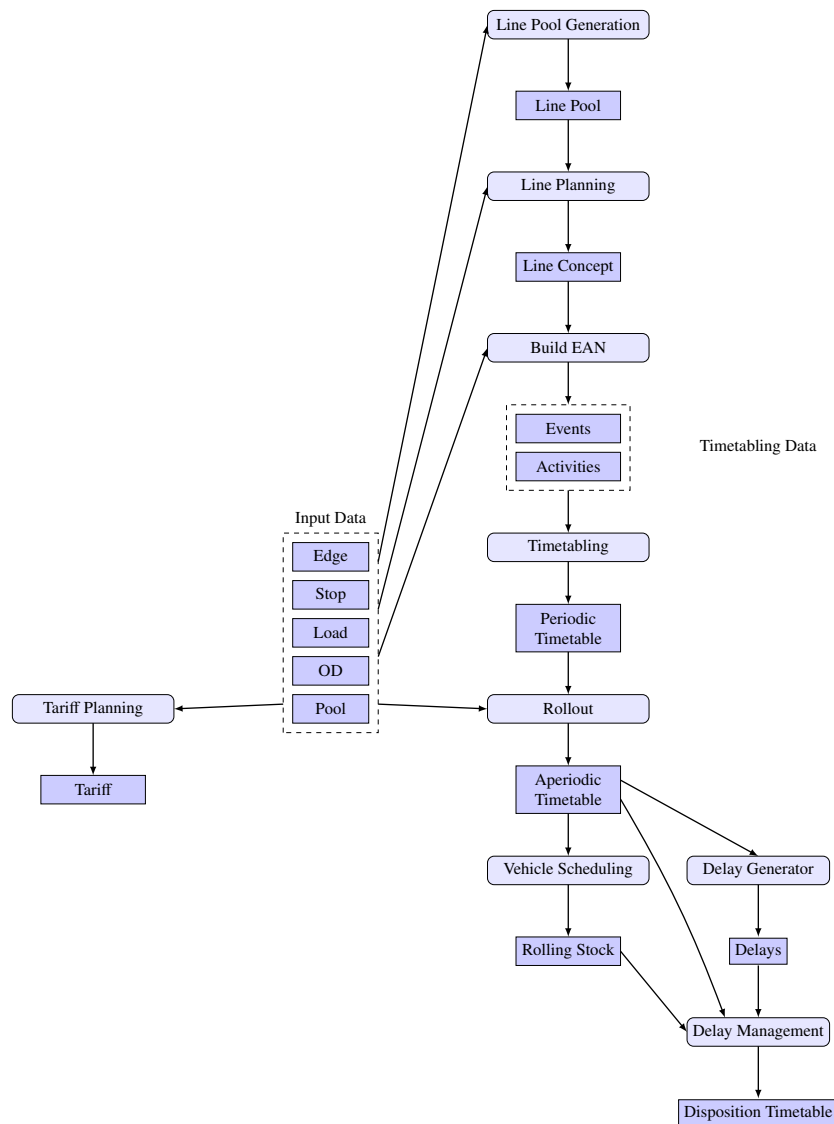


Figure 2.1: Different planning steps considered in LIN-TIM

2.1 Stop Location

In the the stop location step a new PTN is computed according to a given demand and a given infrastructure of stations and tracks.

2.1.1 Input

The following files are needed as input for the classical stop location problems:

- CK `default_existing_stop_file` (Fi `basis/Existing-Stop.giv`) stops of the existing infrastructure network
- CK `default_existing_edge_file` (Fi `basis/Existing-Edge.giv`) edges of the existing infrastructure network
- CK `default_demand_file` (Fi `basis/Demand.giv`) demand at geographical positions

Additionally, there are models for a given infrastructure network. For this, the following files are needed as input:

- CK `filename_node_file` (Fi `basis/Node.giv`) the nodes of the network, including possible stops
- CK `filename_infrastructure_edge_file` (Fi `basis/Edge-Infrastructure.giv`) direct connections between the nodes suitable for public transport
- CK `filename_walking_edge_file` (Fi `basis/Edge-Walking.giv`) possible walking edges between infrastructure nodes
- CK `filename_od_nodes_file` (Fi `basis/OD-Node.giv`) od data based on infrastructure nodes

2.1.2 Output

The following files are produced as output.

- CK `default_stops_file` (Fi `basis/Stop.giv`) stops of the new PTN
- CK `default_edges_file` (Fi `basis/Edge.giv`) edges of the new PTN

2.1.3 Algorithms

Running

R `make sl-stop-location`

will create a new PTN with respect to the given demand points. The following algorithms are available:

- CK `sl_model` CV `dsl` finds an optimal solution for the stop location problem with fixed travel times on PTN edges.
- CK `sl_model` CV `greedy` finds a feasible solution for the stop location problem with fixed travel times on PTN edges with a greedy approach.
- CK `sl_model` CV `dsl-tt` solves CV `dsl` while considering the travel time, including acceleration and deceleration.
- CK `sl_model` CV `dsl-tt-2` solves CV `dsl` while considering the travel time, including acceleration and deceleration.
- CK `sl_model` CV `tt` finds a travel time optimal solution for a given infrastructure network with walking times for the passengers
- CK `sl_model` CV `all` adds every possible stop in a given infrastructure network to the new PTN.

2.2 Line Pool Generation

In the line pool generation step a possible set of lines is computed to use during the line planning step.

2.2.1 Preparation

Run

```
 R make ptn-regenerate-load
```

to compute a new load.

2.2.2 Input

The following files are needed as input:

- CK default_stops_file (Fi basis/Stop.giv) stops of the PTN
- CK default_edges_file (Fi basis/Edge.giv) edges of the PTN
- CK default_loads_file (Fi basis/Load.giv) expected distribution of passengers to PTN edges (depending on CK lpool_model)
- CK default_od_file (Fi basis/OD.giv) OD matrix (depending on CK lpool_model)

2.2.3 Output

The following files are produced as output.

- CK default_pool_file (Fi basis/Pool.giv) line pool, set of possible lines
- CK default_pool_cost_file (Fi basis/Pool-Cost.giv) costs of lines in line pool

2.2.4 Algorithms

To compute a line pool run

```
 R make lpool-line-pool
```

The following algorithms are available:

- CK lpool_model CV tree_based a heuristic based on MST which computes a line pool that at least allows for a feasible line concept for a given load (see 3.2.1)
- CK lpool_model CV restricted_line_duration same as CV tree_based but with additional constraints on the duration of a line (see 3.2.2)
- CK lpool_model CV k_shortest_paths a heuristic which computes the k shortest path for all OD pairs as line pool (see 3.2.3)
- CK lpool_model CV terminal-to-terminal enumerates the complete line pool, starting and ending each line at a terminal (see 3.2.4).
- CK lpool_model CV center-periphery identifies some nodes as centers and some nodes as periphery and then constructs lines between these different pairs of nodes (see 3.2.5).

2.3 Line Planning

In the line planning step a feasible line concept is determined by assigning frequencies to all lines in the line pool.

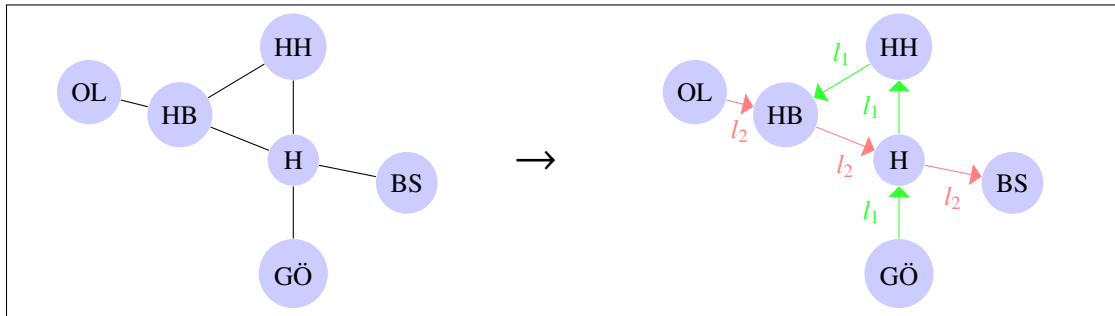


Figure 2.2: Line Planning Step

2.3.1 Preparation

Run

`make ptn-regenerate-load`

to compute a new load.

2.3.2 Input

The following files are needed as input:

- `default_stops_file` (`basis/Stop.giv`) stops of the PTN
- `default_edges_file` (`basis/Edge.giv`) edges of the PTN
- `default_pool_file` (`basis/Pool.giv`) line pool
- `default_pool_cost_file` (`basis/Pool-Cost.giv`) costs of line pool
- `default_loads_file` (`basis/Load.giv`) expected distribution of passengers to PTN edges (depending on `lc_model`)
- `default_od_file` (`basis/OD.giv`) OD matrix (depending on `lc_model`)

2.3.3 Output

The following file is produced as output.

- `default_lines_file` (`line-planning/Line-Concept.lin`) line pool, set of possible lines

2.3.4 Algorithms

To compute a line concept run

`make lc-line-concept`

The following algorithms are available:

- `lc_model` `cost_optimization_model` minimizing the total costs of a line concept (see 3.3.1)
- `lc_model` `cost_restricting_frequencies` the `cost_model`, but with a restriction on the number of frequencies (see 3.3.1)

- CK lc_model CV `direct` optimization model maximizing the number of passengers who can travel on a shortest path from their origin to their destination without having to transfer (see 3.3.2)
- CK lc_model CV `direct_restricting_frequencies` the CV `direct`-model, but with a restriction on the number of frequencies (see 3.3.2)
- CK lc_model CV `direct_relaxation` relaxation of CK lc_model CV `direct`
- CK lc_model CV `cost_greedy_1` greedy heuristic trying to minimize the costs
- CK lc_model CV `cost_greedy_2` another greedy heuristic trying to minimize the costs
- CK lc_model CV `mult-cost-direct` an IP minimizing the weighted sum of costs and direct travelers
- CK lc_model CV `mult-cost-direct-relax` an IP minimizing the weighted sum of costs and direct travelers. Capacity restrictions are aggregated for each edge.
- CK lc_model CV `traveling-time-cg` a column generation procedure minimizing the estimated travel time of passengers. (see 3.3.4)
- CK lc_model CV `traveling-time-mip` (M)IP model for choosing line frequencies and passenger routes minimizing the estimated traveling times. (see 3.3.5)
- CK lc_model CV `minchanges_ip` integer program trying to minimize the weighted number of transfers (see 3.3.7)
- CK lc_model CV `minchanges_cg` column generation procedure trying to minimize the weighted number of transfers (see 3.3.7)
- CK lc_model CV `game` a game-theoretic approach which distributes lines equally among the edges in order to avoid congestion and delays

2.4 Ridepool Generation

In addition to the generation of a line pool and a line concept, it is also possible to generate a ride pool and a ride concept. In ridepooling, we generate several areas of the PTN, i.e. strongly connected subgraphs represented as edge lists. For each area a number of vehicles can be assigned that are allowed to travel freely in the area and operate as an on-demand ridepooling service.

By running

```
 R make rpool-ride-pool
```

we generate a ride pool, to select some of the areas in the next step.

2.4.1 Input

The following files are needed as input:

- CK `default_stops_file` (Fi `basis/Stop.giv`) stops of the PTN
- CK `default_edges_file` (Fi `basis/Edge.giv`) edges of the PTN
- CK `default_pool_file` (Fi `basis/Pool.giv`) line pool
- CK `default_loads_file` (Fi `basis/Load.giv`) expected distribution of passengers to PTN edges (depending on CK `lc_model`)

2.4.2 Output

The following file is produced as output.

- filename_rpool_file (basis/Ridepool.giv) ride pool, set of possible areas

2.4.3 Algorithms

The following models are available:

- rpool_model all Only one ridepooling area is constructed consisting of all edges of the PTN.
- rpool_model demand_heuristic Ridepooling areas are generated from edges with low demand.
- rpool_model tree_based Ridepooling areas are generated using a maximal spanning tree with respect to the edge loads.

2.5 Ride Concept

By running

```
 make rc-ride-concept
```

a line concept and a ride concept are generated simultaneously.

2.5.1 Input

The following files are needed as input:

- default_stops_file (basis/Stop.giv) stops of the PTN
- default_edges_file (basis/Edge.giv) edges of the PTN
- default_pool_file (basis/Pool.giv) line pool
- filename_rpool_file (basis/Ridepool.giv) ride pool, set of possible areas
- default_loads_file (basis/Load.giv) expected distribution of passengers to PTN edges (depending on lc_model)

2.5.2 Output

The following file is produced as output.

- filename_rc_file (basis/Ride-Concept.lin) ride concept, set of areas with assigned number of vehicles
- default_lines_file (line-planning/Line-Concept.lin) line concept, set of lines with assigned frequencies

2.5.3 Algorithms

The following models are available:

- rc_model LPRP_ALPHA An IP is solved, where the edge-specific demand factors $\alpha_{r,e}$ are part of the input.
- rc_model LPRP_BETA An IP is solved, where the edge-specific demand factors $\alpha_{r,e}$ are determined as continuous variables.

2.6 Periodic Timetabling

In periodic timetabling for each Event of a previously created Event-Activity-Network is assigned a time, resulting in a timetable.

2.6.1 Preparation

Run

```
[R] make ean
```

to create an Event-Activity-Network from an existing line concept.

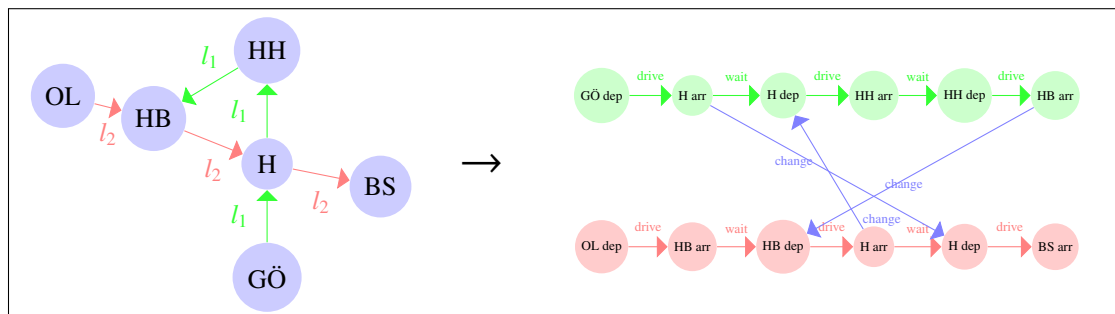


Figure 2.3: Creation of an EAN

2.6.2 Input

The following files are needed as input:

- `[CK] default_activities_periodic_file ([Fi] timetabling/Activities-periodic.giv)` Activities generated by the line concept.
- `[CK] default_events_periodic_file ([Fi] timetabling/Events-periodic.giv)` Events generated by the line concept.

For some timetabling procedures also the following files are necessary:

- `[CK] default_stops_file ([Fi] basis/Stop.giv)` stops of the PTN
- `[CK] default_edges_file ([Fi] basis/Edge.giv)` edges of the PTN
- `[CK] default_lines_file ([Fi] line-planning/Line-Concept.lin)` line concept calculated in the previous planning step
- `[CK] filename_tim_fixed_times ([Fi] timetabling/Fixed-timetable-periodic.tim)` fixed time intervals for some events

2.6.3 Output

The following files are produced as output.

- `[CK] default_timetable_periodic_file ([Fi] timetabling/Timetable-periodic.tim)`

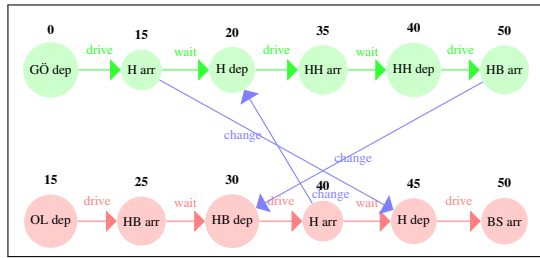


Figure 2.4: Periodic Timetabling Step

2.6.4 Algorithms

To compute a line concept run

`R` `make tim-timetable`

The following algorithms are available by setting the config parameter `CK tim_model` to one of the following:

- `CV MATCH` (default value) Heuristic that sets the times of driving and waiting activities to their lower bounds and then tries to minimize change durations.
- `CV con_prop` Heuristic that fixes events and propagates the implied constraints to the whole network.
- `CV csp` Heuristic that transforms the problem to a Constraint Satisfaction Problem and finds a feasible solution for it. *Currently not included in the release version of LinTim.*
- `CV ns_improve` Improvement procedure (known as Network-Simplex or Modulo-Simplex) that requires a feasible timetable.
- `CV csp_ns` Runs `csp` and `ns_improve` afterwards. *Currently not included in the release version of LinTim.*
- `CV con_ns` Runs `con_prop` and `ns_improve` afterwards.
- `CK ip` Models the Periodic Timetabling Problem as an IP and solves it.
- `CK cb_ip` Models the Periodic Timetabling Problem as a cycle based IP and solves it.
- `CK ns_cb` First improve a given feasible solution using the network simplex and afterwards optimize it using a cycle based IP
- `CK phase-one` Uses a phase 1 simplex method for finding a feasible timetable

2.7 Tariff Planning

Tariff planning computes a new tariff minimizing the deviation from given reference prices.

2.7.1 Input

The following files are needed as input:

- `CK default_stops_file` (`Fi` `basis/Stop.giv`), stops of the PTN
- `CK default_edges_file` (`Fi` `basis/Edge.giv`), edges of the PTN
- `CK default_od_file` (`Fi` `basis/OD.giv`), OD matrix

- filename_tariff_reference_price_matrix_file (basis/Reference-Price-Matrix.giv), matrix of reference prices

Depending on the parameter values a routing in the PTN is also needed:

- filename_routing_ptn_input (basis/Routing-ptn.giv), routing in the ptn

2.7.2 Output

The following files are produced as output independent of taf_model:

- filename_tariff_price_matrix_file (tariff/Price-Matrix.taf), prices for each OD pair
- filename_tariff_properties_file (statistic/tariff-properties.sta), statistic file containing information whether the no-elongation property and the no-stopover property (see Section 3.7.4) are fulfilled for the computed tariff. For a zone tariff it is also checked whether the zones are connected. Additionally, the tariff model and in case of a zone tariff the counting type are specified.

If taf_model is zone, then the two following files are produced as output as well:

- filename_tariff_zone_file (tariff/Zones.taf), assignment of stops to zones
- filename_tariff_zone_price_file (tariff/Zone-Prices.taf), prices per number of traversed zones

If taf_model is network_distance or zone, the routing that is computed is also written as output file:

- filename_routing_ptn_output (basis/Routing-ptn.giv), routing in the ptn

2.7.3 Algorithms

To compute a tariff, run

```
 make taf-tariff
```

The following algorithms are available:

- taf_model flat, optimization model determining a flat tariff
- taf_model beeline_distance, optimization model determining an affine beeline distance tariff,
- taf_model network_distance, optimization model determining an affine network distance tariff,
- taf_model zone, optimization model determining a zone tariff.

2.8 Vehicle Scheduling

In the vehicle scheduling problem a set of routes for service vehicles is calculated to serve the given public transportation system. There are two base models, one based on an aperiodic timetable, the other only on a line concept. The following information is based on the classic formulations, based on the aperiodic timetable.

2.8.1 Preparation

Run

`R` make ro-rollout

and

`R` make ro-trips

with `CK` rollout_whole_trips set to `CV` true to create all input files needed for the vehicle scheduling problem.

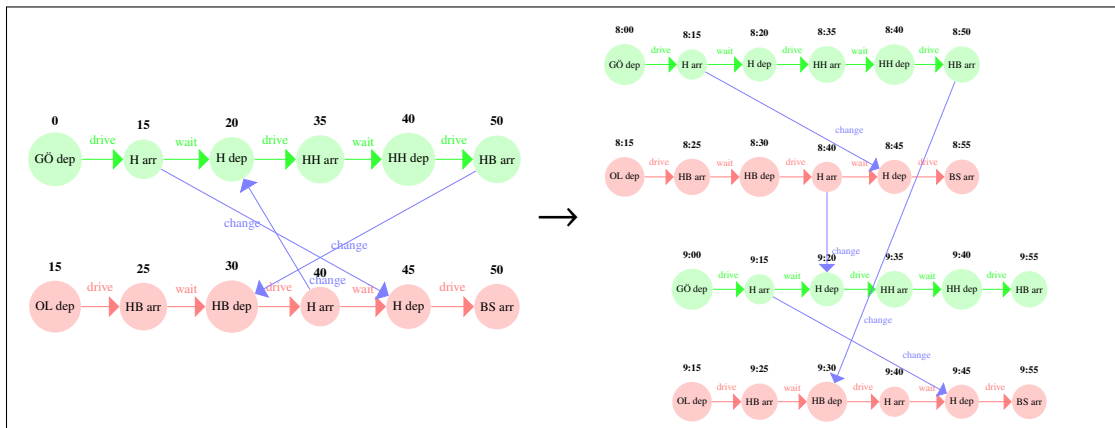


Figure 2.5: Rollout Step

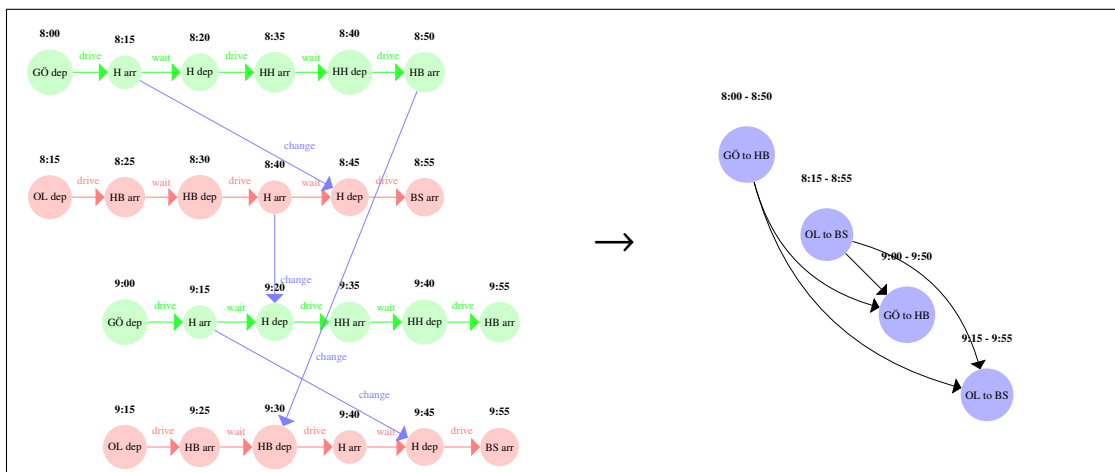


Figure 2.6: Rollout to Trips Step

2.8.2 Input

For the rollout

The following files are needed as an input for the rollout-step:

- `CK` default_edges_file (`Fi` basis/Edge.giv) edges of the PTN
- `CK` default_headways_file (`Fi` basis/Headway.giv) headways of the PTN

- default_lines_file (line-planning/Line-Concept.lin) frequencies of the lines
- default_events_periodic_file (timetabling/Events-periodic.giv) periodic events
- default_activities_periodic_file (timetabling/Activities-periodic.giv) periodic activities
- default_timetable_periodic_file (timetabling/Timetable-periodic.tim) periodic timetable

Only for the model

The following files are needed as an input for the vehicle scheduling step:

- default_stops_file (basis/Stop.giv) stops of the PTN
- default_edges_file (basis/Edge.giv) edges of the PTN
- default_trips_file (delay-management/Trips.giv) trips for the vehicle schedule
- default_events_expanded_file (delay-management/Events-expanded.giv) aperiodic events

2.8.3 Output

The following files will be produced:

- default_vehicle_schedule_file (vehicle-scheduling/Vehicle_Schedules.vs) the vehicle schedule

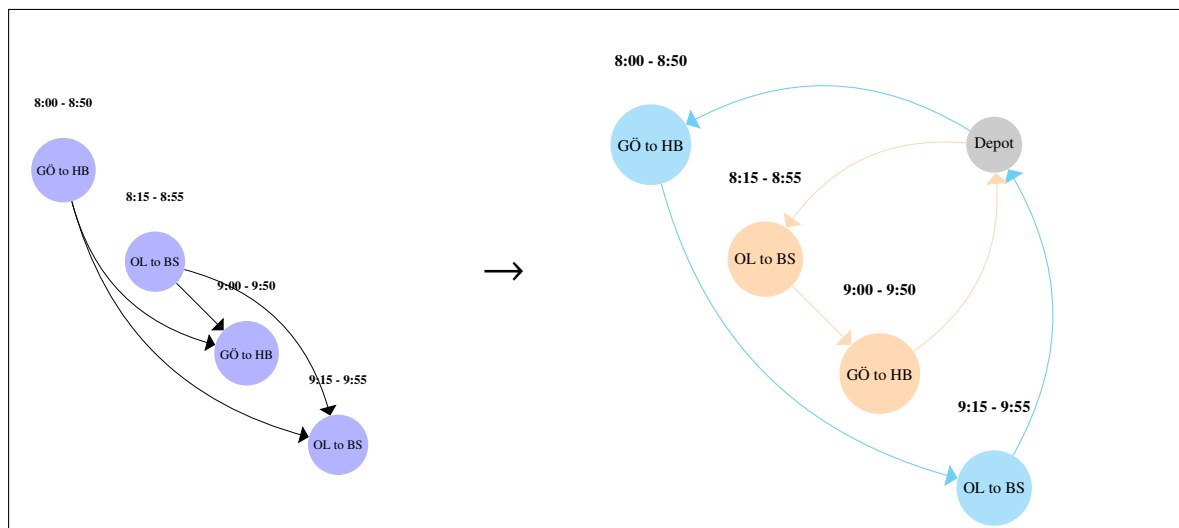


Figure 2.7: Vehicle Scheduling

2.8.4 Algorithms

To compute a vehicle schedule run

make vs-vehicle-schedules

. The following models are available

- CK vs_model CV MDM1 Minimizing the number of vehicles (see 3.8.1)
- CK vs_model CV MDM2 Minimizing the number of vehicles (see 3.8.2)
- CK vs_model CV ASSIGNMENT_MODEL Minimizing the overall costs (see 3.8.3)
- CK vs_model CV TRANSPORTATION_MODEL Minimizing the overall costs (see 3.8.4)
- CK vs_model CV NETWORK_FLOW_MODEL Minimizing the overall costs (see 3.8.5) (see 3.8.1)
- CK vs_model CV CANAL_MODEL More detailed version of CV ASSIGNMENT_MODEL (see 3.8.6)
- CK vs_model CV LINE_BASED vehicle scheduling only based on line planning (see 3.8.7)
- CK vs_model CV SIMPLE will create a vehicle schedule driving the lines back and forth (see 3.8.8)
- CK vs_model CV IP solve a simple ip model (see 3.8.9)

2.9 Delay Management

Delay management computes a new (disposition) timetable based on an existing timetable and unforeseen delays that make the original timetable infeasible.

2.9.1 Preparation

If you have not already done so for the vehicle scheduling part, run

```
 R make ro-rollout
```

to expand a previously computed periodic timetable on a periodic Event-Activity Network into an aperiodic timetable on an aperiodic Event-Activity Network.

2.9.2 Input

For the rollout

The following files are needed as an input for the rollout-step:

- CK default_edges_file (Fi basis/Edge.giv) edges of the PTN
- CK default_headways_file (Fi basis/Headway.giv) headways of the PTN
- CK default_lines_file (Fi line-planning/Line-Concept.lin) frequencies of the lines
- CK default_events_periodic_file (Fi timetabling/Events-periodic.giv) periodic events
- CK default_activities_periodic_file (Fi timetabling/Activities-periodic.giv) periodic activities
- CK default_timetable_periodic_file (Fi timetabling/Timetable-periodic.tim) periodic timetable

Aperiodic Event-Activity Network

These files, generated by the rollout step, are actually used for delay management:

- default_events_expanded_file (delay-management/Events-expanded.giv) for the events
- default_activities_expanded_file (delay-management/Activities-expanded.giv) for the activities
- default_timetable_expanded_file (delay-management/Timetable-expanded.tim) for the initial timetable

Delays

There are two types of delays, which are both optional, and which go into separate files:

- default_event_delays_file (delay-management/Delays-Events.giv)
- default_activity_delays_file (delay-management/Delays-Activities.giv)

You can either manually enter delays on events and/or activities through these files, or use an automatic (random) delay generator by running

```
 make dm-delays
```

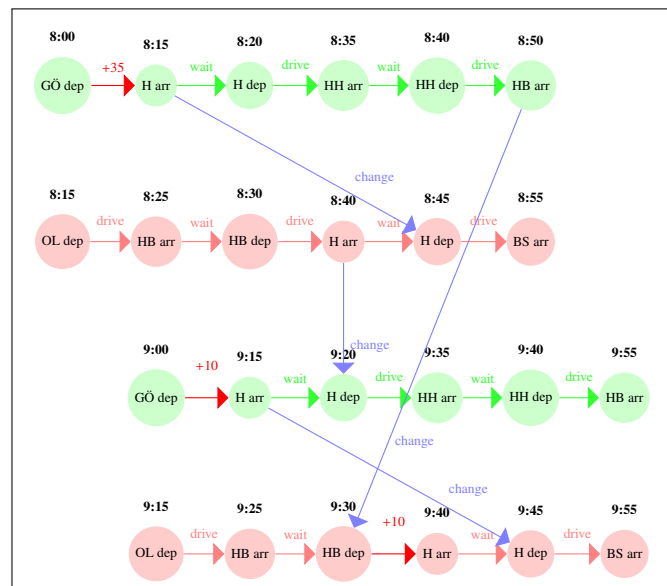


Figure 2.8: Generation of Delays

2.9.3 Output

The result of the delay management step is a new disposition timetable with no departure earlier than in the original timetable, and with all the delays respected: default_disposition_timetable_file (delay-management/Timetable-disposition.tim)

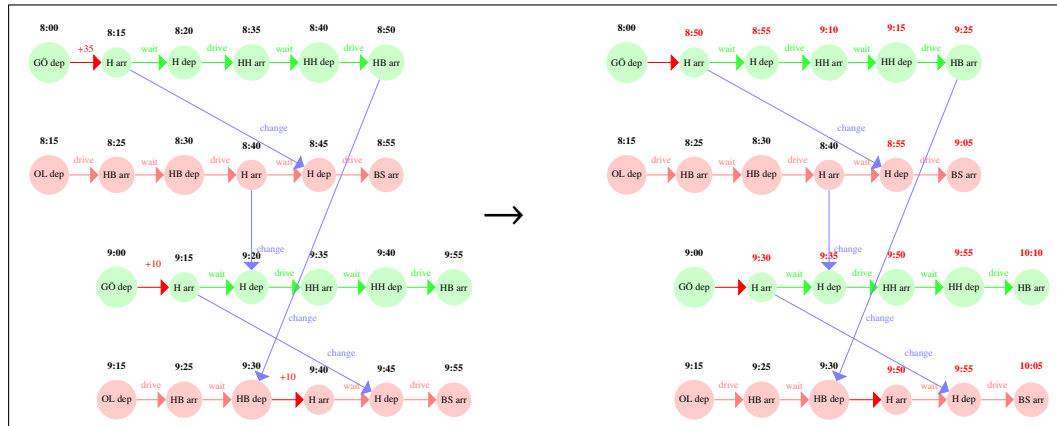


Figure 2.9: Delay Management Step

2.9.4 Algorithms

The delay management step is invoked via

`R` `make dm-disposition-timetable`

The main algorithms implemented in `LINTIM` are the IP-based algorithms

- `CK` `DM_method` `CV` `DM1`
- `CK` `DM_method` `CV` `FSFS`
- `CK` `DM_method` `CV` `FRFS`
- `CK` `DM_method` `CV` `EARLYFIX`
- `CK` `DM_method` `CV` `PRIORITY`
- `CK` `DM_method` `CV` `PRIOREPAIR`
- `CK` `DM_method` `CV` `best-of-all` which computes all of the above and then chooses the best solution
- `CK` `DM_method` `CV` `DM2`
- `CK` `DM_method` `CV` `DM2-pre`

These need a solver configured via `CK` `DM_solver` (like `CV` `Xpress` or `CV` `Gurobi`, see Section 1.2 for details). In contrast, the most basic method without any optimization is just delaying all the events according to the delays, `CK` `DM_method` `CV` `propagate`, where a maximum waiting time for change activities can be configured in seconds by `CK` `DM_propagate_maxwait`, and headway activities can be turned around automatically whenever this would not result in additional delay for the train that was originally scheduled to go first, by setting `CK` `DM_propagate_swapHeadways` to `CV` `true` (the default).

2.10 Integrated Planning

`LINTIM` also contains algorithms to compute multiple planning stages at once or in non-ordinary order.

2.10.1 Algorithms

Timetabling and Passenger Routing: Run

`make int-tim-pass`

to solve the integrated timetabling and passenger routing problem. More information can be found in Section 3.10.1.

Timetabling and Vehicle Scheduling: Run

`make int-tim-veh`

to solve the integrated timetabling and aperiodic vehicle scheduling problem. The passenger routes are fixed in this model. More information can be found in Section 3.10.2.

Line Planning and Timetabling: Run

`make int-lin-tim-pass`

to solve the integrated line planning and timetabling problem. This also includes passenger routing in the timetabling stage. More information can be found in Section 3.10.3.

Timetabling and Vehicle Scheduling: Run

`make int-lin-tim-pass-veh`

to solve the integrated line planning, timetabling and aperiodic vehicle scheduling problem. This also includes passenger routing in the timetabling stage. More information can be found in Section 3.10.4.

Robust Timetabling and Vehicle Scheduling using Machine Learning Run

`make int-rob-ml-algo`

to solve the problem of finding a robust timetable and vehicle schedule based on the current solution. More information can be found in Section 3.10.5.

2.10.2 The Eigenmodel

The eigenmodel is an iterative approach to integrated public transport planning, re-organizing the sequential planning approach to allow new optimization models, solving the original problem in different orderings. For more details, see Section 3.10.6.

Chapter 3

Detailed Description of Algorithms

3.1 Stop Location

3.1.1 Without a given infrastructure network

Running

`R` `make sl-stop-location`

will create a new PTN with respect to the given demand points. Here, all demand points have to be *covered* by at least one station, i.e., the distance between the demand point and the nearest station has to be less than a given radius.

The parameters used for adjusting the model are the following:

- `CK` `sl_distance` norm used for measuring the distance between demand points, stations etc. Currently the only option is `euclidean_norm`.
- `CK` `sl_radius` maximal distance a demand point may have from a station to be covered.
- `CK` `sl_destruction_allowed` whether it is allowed to remove station that are not covering any demand points.
- `CK` `sl_new_stop_default_name` name prefix to be given to new stops.

Fixed travel time on edges

The first step of the classical stop location problem which uses fixed travel times on the edges is to compute a finite dominating set of candidates for new stations. When using the euclidean norm for measuring distance this finite dominating set can easily be computed as the intersection of the tracks and circles around the demand point with the given radius and the already existing stops.

Optimization model For the optimization model define the constants

$$a_{ps} = \begin{cases} 1 & \text{if demand point } p \text{ is covered by candidate } s \\ 0 & \text{otherwise} \end{cases}$$

and the variables

$$x_s = \begin{cases} 1 & \text{if candidate } s \text{ is established as station} \\ 0 & \text{otherwise} \end{cases}$$

The objective is to minimize the number of established stations such that all demand points are covered. The following optimization model is solved to find an optimal solution for the stop location problem.

$$\begin{aligned}
(DSL) \quad & \min \sum_{s \in \mathcal{S}} x_s \\
\text{s.t.} \quad & \sum_{s \in \mathcal{S}} a_{ps} x_s \geq 1 \quad \forall p \in \mathcal{P} \\
& x_s \in \{0, 1\} \quad \forall s
\end{aligned}$$

For more information, see [34, 38].

Greedy heuristic The greedy heuristic find a feasible solution to the stop location problem by successively adding the candidate which covered most uncovered demand points at this point in time.

For more information, see [34, 38].

Travel time considering acceleration/deceleration

When considering the acceleration and deceleration phases of vehicles, the following parameters have to be set:

- `sl_acceleration`
- `sl_deceleration`
- `sl_waiting_time`

For more information, see [4].

3.1.2 For a given infrastructure network

If a complete infrastructure network, i.e., an infrastructure network with walking and node-based od-information, is given, the stop location models `sl_model` `tt` and `all` can be used. For `tt`, a selection of stops is chosen such that the overall travel time of the passengers (containing public transport use as well as walking) is minimized. Additionally, creating stops is penalized by `sl_cost_of_stop`. For `all`, all possible stop points are converted to stops in the PTN.

Given forbidden edges in the infrastructure (`sl_forbidden_edges`) and given restricted turns in the infrastructure (`sl_restricted_turns`) can be converted into the resulting ptn information as well when their corresponding config parameter is set to `true`.

3.2 Line Pool Generation

A new line pool is created by running

```
 make lpool-line-pool
```

There are four main approaches implemented in `LINTIM`.

1. Different variants of a tree-based heuristic. These are invoked by setting `lpool_model` to either `tree_heuristic` or `restricted_line_duration` and are described in Sections 3.2.1 and 3.2.2, respectively.
2. A center-periphery heuristic, which is called with the parameter setting `lpool_model` `center-periphery` and described in Section 3.2.5.
3. k -shortest paths. This is called by setting `lpool_model` to `k_shortest_paths` and is described in Section 3.2.3.
4. All paths between pairs of a given set of terminal stations. This is executed when setting `lpool_model` to `terminal-to-terminal` and described in Section 3.2.4.

3.2.1 Creating a new line pool with the tree based heuristic

For an undirected PTN a line pool \mathcal{L} may be created from an existing PTN (default_edges_file (basis/Edge.giv), default_stops_file (basis/Stop.giv)), a given default_loads_file (basis/Load.giv) (see Chapter 8), and a given default_od_file (basis/OD.giv) by running

```
 make lpool-line-pool
```

with lpool_model tree_based, which creates a line pool default_pool_file (basis/Pool.giv) and a corresponding default_pool_cost_file (basis/Pool-Cost.giv). How the line costs are computed can be seen in Section 3.2.6.

The algorithm iteratively creates minimum spanning trees, on which lines are created in three different possible ways:

- as a path from a leaf of the MST to another leaf,
- as a path from a leaf of the MST to a *terminal* or
- as a path from a terminal to another terminal.

Here *terminals* are nodes of a high node degree. Each of the three classes of lines has to fulfill different requirements, which can be seen in the discussion of the configuration parameters. Lines are created until a feasible line concept can be found within the line pool or until the maximal number of iterations is reached. One iteration consists of the following steps:

1. Determine a set of preferred edges.
2. Compute minimum spanning trees and create lines until all preferred edges are covered sufficiently often or no other line can be added.
3. Test whether a feasible line concept can be found in the constructed pool.

In the first iteration preferred edges are chosen from the usage rate in the shortest paths of the OD pairs. Later, the lower frequencies given in the loads file are lowered until a feasible line concept can be found for the new frequencies, and the edges for which the original frequencies are not met are chosen as preferred edges.

The edge weight used to compute the minimum spanning tree is zero if the edge is preferred and the physical length of the edge otherwise.

The configuration parameters are:

- lpool_max_iterations: the maximal number of iterations.
- lpool_ratio_od: the ratio of the most frequently used edges in shortest paths of the passengers, which are preferred in the first iteration.
- lpool_node_degree_ratio: the percentage of the maximal node degree, which has to be attained to qualify a node as a terminal. In the first iteration the node degree depends on the incident edges in the PTN, later it depends on the lines passing the node.
- lpool_min_cover_factor: a preferred edge has to be covered $\lceil \frac{f_e^{\min}}{\text{lpool_min_cover_factor}} \rceil$ times in order to be sufficiently covered.
- lpool_max_cover_factor: if a new line covers an edge more than $f_e^{\max} \cdot \text{lpool_max_cover_factor}$ it cannot be used in the line pool.
- lpool_min_edges: the minimal number of edges in a line from a leaf to a terminal or from a terminal to another terminal.

- `lpool_min_distance_leaves`: the minimal euclidean distance between two leaves to allow for a line between them.
- `lpool_add_shortest_paths`: determines whether shortest paths are to be added as additional lines to the line pool.
- `lpool_ratio_shortest_paths`: the percentage of the maximal number of passengers in an OD pair which has to be attained in order to add the shortest path for an OD pair as a line. This parameter is only relevant if `lpool_add_shortest_paths` is set to true.
- `lpool_append_single_edges`: Add all links as separate lines to the line pool.

Note that all lines which are created here are cycle-free, as they are either a path in a minimal spanning tree or a shortest path in a network with non-negative edge-lengths.

Possible additional restrictions on the created lines are

- `lpool_restrict_terminals` Only allow lines that start or end at terminals given in `filename_terminals_file` (`basis/Terminals.giv`)
- `lpool_restrict_turns` Only allow lines that do not contain a restricted turn given in `filename_turn_restrictions` (`basis/Restricted-Turns.giv`)
- `lpool_restrict_forbidden_edges` Do not allow the forbidden links in `filename_forbidden_links_file` (`basis/Edge-forbidden.giv`) to be contained in lines

For more information, see [10].

3.2.2 Creating a line pool while restricting the duration of the lines

When running

```
 make lpool-line-pool
```

with the parameter `lpool_model` set to `restricted_line_duration` the tree based heuristic (see 3.2.1) is performed with additional constraints on the duration of lines. This is influenced by the following parameters:

- `ean_model_weight_drive` to decide how the duration of a line is computed
- `ean_model_weight_wait` to decide how much waiting time is added in each station
- `period_length` used to determine the feasible duration interval
- `vs_turn_over_time` used to determine the feasible duration interval
- `lpool_restricted_maximum_buffer_time` used to determine the feasible duration interval
- `lpool_restricted_allow_half_period` determines if lines which fit into the interval at exactly half a period minus the corresponding buffer times are allowed to be added

The feasible interval for the duration of a line `mod` `period_length` is defined as

$$[\text{input_type="checkbox"/}\code{period_length} - \text{input_type="checkbox"/}\code{vs_turn_over_time} \\ - \text{input_type="checkbox"/}\code{lpool_restricted_maximum_buffer_time}, \\ \text{input_type="checkbox"/}\code{period_length} - \text{input_type="checkbox"/}\code{vs_turn_over_time}].$$

Note: There will be no shortest paths added to line pools created by this heuristic, i.e.,

`lpool_add_shortest_paths` has no influence.

For more information, see [24].

3.2.3 Creating a line pool by k shortest paths

Another possibility is to create a line pool with corresponding line costs by using the k shortest paths for each OD pair as lines and then deleting lines which are nested in other lines. To do so run

```
 R make lpool-line-pool
```

with the parameters

- CK `lpool_model` CV `k_shortest_paths`
- CK `lpool_number_shortest_paths`, which gives the number of shortest paths which are to be computed for each OD pair.

3.2.4 Terminal-to-terminal

When terminals are given, i.e., CK `filename_terminals_file` (Fi `basis/Terminals.giv`), running

```
 R make lpool-line-pool
```

with the parameters

- CK `lpool_model` CV `terminal-to-terminal`

will result in the enumeration of all possible lines starting and ending at a terminal and therefore finding all possible lines respecting the terminal restrictions. Note that this may result in large computation times and a large number of lines in the linepool, depending on your PTN.

3.2.5 Center-Periphery

Another method to create a line pool is running

```
 R make lpool-line-pool
```

with the parameter

- CK `lpool_model` CV `center-periphery`

This algorithm identifies some nodes as centers and some nodes as periphery to construct lines between those different pairs of nodes. It is a heuristic that tries to identify natural patterns in the PTN and the OD data. The following parameters have to be specified:

- CK `lpool_centers_fraction` Fraction of nodes that can become centers
- CK `lpool_periphery_radius_factor` Factor for the mean distance of two nodes in the PTN to choose periphery nodes
- CK `lpool_direct_periphery_lines_factor`
- CK `lpool_center_radius_factor` Percentage of the mean distance in the PTN determining the radius of the centers
- CK `lpool_concatenate_lines_factor` Factor for the mean OD value to choose node pairs for which direct lines are created by concatenating existing lines
- CK `lpool_min_degree_center` Minimal node degree that a center node must have
- CK `lpool_min_times_edge_covered` Factor for the minimal frequency of each edge to determine how many times it should be covered by a line
- CK `lpool_max_iter_postprocessing` Maximal number of iterations for each postprocessing step

- `lpool_opt_cost` Determines, if lines are created along shortest paths w.r.t. edge costs or edge lengths
- `lpool_plot_centers` Optional parameter. If set to True, plots of the PTN are created where the centers are highlighted in red and the periphery nodes are highlighted in blue. The images are written to the `/graphics` directory.

Choice of centers and periphery

Let V be the set of nodes in the PTN and $n = |V|$. For each node v in the PTN we compute the number of *interactions* for this node as a measure of its importance as

$$\sum_{u \in V: u \neq v} OD_{u,v} + OD_{v,u}.$$

Only nodes with a degree of at least `lpool_min_degree_center` in the PTN are candidates for centers. We order the set of those nodes non-increasingly by the interactions computed above. The parameter `lpool_centers_fraction` states which portion of the candidates with the highest interaction values should remain center candidates. We look in the interval

$$n \cdot \text{code} + [-0.2 \cdot n, 0.2 \cdot n]$$

for the greatest difference in the interactions between two neighboring nodes in the sorted list to identify the biggest jump in the interactions around the desired portion. Among those candidates we choose the centers in such a way that no centers have a distance less than the mean distance between any two nodes in the PTN multiplied with `lpool_center_radius_factor` and the sum of the interactions of the chosen centers is maximized. This is modelled by an IP which is solved using Gurobi.

Now we determine the periphery nodes. All Endstations, i. e. all nodes with degree 1, are defined to be periphery nodes. Furthermore, all non-center nodes with a distance greater than

$$\text{code} \cdot \sum_{c \text{ center}} \sum_{u \in V: u \neq c} \frac{\text{distance}(u, c)}{(\text{number of centers})(n - 1)}$$

(the mean distance from a node to a center) to it's closest center become periphery nodes.

Line generation

If the parameter `lpool_opt_cost` is set to True then all shortest paths are computed with respect to the costs of the edges, otherwise with respect to their lengths.

For each pair of centers lines are generated along all shortest paths between them.

For each pair of a center and a periphery node lines are generated along all shortest paths between them, if this line is not contained in another line which was already generated.

For each pair of periphery nodes lines are generated along all shortest paths, if the corresponding OD-value is greater than the mean OD-value multiplied with `lpool_direct_periphery_lines_factor`.

As a next step we concatenate for each node pair with an OD-value greater than the mean OD-value multiplied with `lpool_concatenate_lines_factor` all yet generated lines from the start node to it's closest center, the lines between the closest centers of the nodes and the lines from the closest center of the end node to itself. This gives direct connections between the most important OD pairs, but they fit with the center-periphery pattern we want to establish.

If after this procedure there is still a node not covered by a line, we create lines from this node to its closest center along all shortest paths.

In a last step lines along small detours are created. For this we look at all edges that are covered by the smallest number of lines. For those edges, the closest peripheries and centers of both endpoints are determined. Then lines containing the specified edge are created, starting from one of the closest periphery or center nodes of the left node and ending in one of the closest nodes of the right node. This is done along all shortest paths and for all such pairs of closest peripheries or centers. If this procedure creates cycles it is aborted and the single edge is added as a line to the pool.

Postprocessing

The idea of the postprocessing step is to generate lines, that contain edges which are not yet covered by enough lines proportionally to their minimal frequency. In a while loop, we consider a residual network which is initially the same as the PTN. In every iteration the edge weights are updated and lines are created along the shortest paths found in the residual network. We use all line generating algorithms from the previous section. If no new line was found or the maximal number of iterations was reached the while loop is aborted. Let $n(e)$ be the number of lines that contain the edge e and let e_{res} denote the copy of the edge e in the residual network. Then we compute for each residual edge it's lower frequency bound as

$$f_{min}(e_{res}) = \max\{0, f_{min}(e) - n(e) \cdot \min_{L:e \in L} \min_{e' \in L} f_{min}(e')\}$$

where $e \in L$ means that the edge e is contained in the line L . The length of the residual edge is then

$$l(e_{res}) = \exp(f_{min}(e) - f_{min}(e_{res})) \cdot l(e)$$

with $l(e)$ being the original length of the edge. The shortest paths are now computed with respect to this new length. If we set `CK lpool_opt_cost` to `True` the costs of the edges are used instead of their lengths. After this frequency based postprocessing step we do a covered based postprocessing step. The only difference is that the new lengths are set to

$$l(e_{res}) = n(e) \cdot l(e).$$

3.2.6 Line costs

The costs of the lines created by

```
R make lpool-line-pool-cost
```

are of the following form

$$\begin{aligned} \text{cost}_l = & \text{CK lpool_costs_fixed} \\ & + \sum_{e \in l} (\text{CK lpool_costs_length} \cdot \text{length}_e + \text{CK lpool_costs_edges}) \\ & + \text{CK lpool_costs_vehicles} \cdot \left[x \cdot \frac{\text{duration}_l + \text{CK vs_turn_over_time}}{\text{CK period_length}} \right], \end{aligned}$$

where x is 1 for directed and 2 for undirected lines (since undirected lines need to be traversed in both directions). The duration of a line is computed as described in Section 3.2.2.

For a given line pool `CK default_pool_file` (`Fi basis/Pool.giv`) a corresponding cost file `CK default_pool_cost_file` (`Fi basis/Pool-Cost.giv`) can be created by running

```
R make lpool-line-pool-cost.
```

3.3 Line Planning

The line planning problem can be solved by running

```
R make lc-line-concept.
```

The following subsection describe the corresponding algorithms.

3.3.1 Cost

Running

`R` `make lc-line-concept`

with `CK` `lc_model` `CV` `cost`, `CV` `cost_greedy_1` or `CV` `cost_greedy_2` results in solving the line planning model such that the operational costs are minimized. Operational costs in line planning are defined as line based costs cost_l for all line $l \in \mathcal{L}$ and are calculated once per frequency. This means the operation costs of a line concept with line frequencies f_l for line $l \in \mathcal{L}$ is

$$\sum_{l \in \mathcal{L}} \text{cost}_l \cdot f_l.$$

Optimal solution

Running

`R` `make lc-line-concept`

with `CK` `lc_model` `CV` `cost` results in solving the classic costs minimizing line planning problem, described in [36], to optimality. The corresponding integer program is

$$\begin{aligned} \text{(LP-Cost) } \min & \sum_{l \in \mathcal{L}} \text{cost}_l \cdot f_l \\ \text{s.t. } & f_e^{\min} \leq \sum_{l \in \mathcal{L}: e \in l} f_l \leq f_e^{\max} \quad \forall e \in E \\ & f_l \in \mathbb{Z} \quad \forall l \in \mathcal{L}. \end{aligned}$$

which is solved either by the solver Gurobi or by the solver Xpress, depending on whether `CK` `lc_solver` is set to `CV` `GUROBI` or `CV` `XPRESS`.

System frequency

Running

`R` `make lc-line-concept`

with `CK` `lc_model` `CV` `cost` and `CK` `lc_common_frequency_divisor` set to a value unequal to 1, will result in solving the problem with a system frequency, i.e., a frequency is only allowed in a solution, if it is the multiple of the system frequency `CK` `lc_common_frequency_divisor`. A value ≤ 0 will test any system frequency (except for 1) and output the best solution. For more information, see [9].

Heuristic solutions

Running

`R` `make lc-line-concept`

with `CK` `lc_model` `CV` `cost_greedy_1` or `CV` `cost_greedy_2` results in solving a heuristic for the cost model described in this section. Lines are added to the line concept in a greedy way (w.r.t. the costs of the lines) until the lower frequency bounds on the edges are fulfilled. Note that these algorithms ignore the upper frequency bounds and are therefore not guaranteed to find a feasible solution w.r.t. these bounds. The algorithms are described in [33].

Restricting the number of frequencies

Running

```
[R] make lc-line-concept
```

with `[CK] lc_model [CV] cost_restricting_frequencies` results in solving the cost model, while restricting the number of possible frequencies. The resulting model has more variables than the original problem, which may result in much longer running times. Even if the number of possible frequencies is unrestricted (-1) this is still not the same model as cost due to `[CK] lc_maximal_frequency`.

- `[CK] lc_solver` either `[CV] GUROBI` or `[CV] XPRESS`, the solver to use to solve the model
- `[CK] lc_number_of_possible_frequencies` restrict the number of possible frequencies (-1=infinity)
- `[CK] lc_timelimit` the time limit for the solver (-1=infinity)
- `[CK] lc_maximal_frequency` the maximal allowed frequency

Fixed Lines

Running

```
[R] make lc-line-concept
```

with `[CK] lc_model [CV] cost` and `[CK] lc_respect_fixed_lines` set to `[CV] true`, will result in solving the cost model while fixing the line frequencies given by `[CK] filename_lc_fixed_lines` (`[Fi] line-planning/Fixed-Lines.lin`). Fixed lines will count towards fulfilling the lower frequency bounds for feasibility and need to be included in the line pool, i.e., `[CK] default_pool_file` (`[Fi] basis/Pool.giv`) and `[CK] default_pool_cost_file` (`[Fi] basis/Pool-Cost.giv`). The capacities for fixed lines need to be given in `[CK] filename_lc_fixed_line_capacities` (`[Fi] line-planning/Line-Capacities.lin`).

Forbidding Links

It is possible to forbid the usage of certain links in the PTN by setting `[CK] lc_respect_forbidden_edges` to `[CV] true` and giving the forbidden links in `[CK] filename_forbidden_links_file` (`[Fi] basis/Edge-forbidden.giv`). Then, the upper bounds for all the corresponding links will be set to 0 in the optimization problem, guaranteeing that lines using these links will not be used in a feasible line concept. This may be useful when considering a PTN with multiple public transport modes, i.e., having tracks and streets and optimizing a bus network that may not use tracks. Can be combined with setting fixed lines for the forbidden edges.

3.3.2 Direct

Running

```
[R] make lc-line-concept
```

with `[CK] lc_model [CV] direct` results in solving an optimization model which aims to maximize the number of passengers which can travel on a shortest path from their origin to their destination without having to transfer between lines. The shortest path is determined w.r.t. `[CK] ean_model_weight_drive`. Upper and lower frequency bounds have to be fulfilled similar to the cost model and additionally capacity constraints on all edges have to be satisfied. Fixing lines and forbidding links is possible here as well, see the documentation for the cost model in Section 3.3.1.

The following parameters control the behavior of the algorithm:

- `ean_model_weight_drive`
- `gen_passengers_per_vehicle`
- `lc_budget`
- `lc_common_frequency_divisor`
- `lc_direct_optimize_costs`
- `lc_mip_gap`
- `lc_mult_relation`
- `lc_respect_fixed_lines`
- `lc_respect_forbidden_edges`
- `lc_timelimit`
- `period_length`

For more information on the model, see [3].

Restricting the number of frequencies

Running

```
 make lc-line-concept
```

with `lc_model` `direct_restricting_frequencies` results in solving the direct model, while restricting the number of possible frequencies. The resulting model has more variables than the original problem, which may result in much longer running times. Even if the number of possible frequencies is unrestricted (-1) this is still not the same model as direct due to `lc_maximal_frequency`.

- `gen_passengers_per_vehicle`
- `lc_budget`
- `ean_model_weight_drive`
- `lc_common_frequency_divisor`
- `lc_timelimit`
- `lc_maximal_frequency`

System frequency

Running

```
 make lc-line-concept
```

with `lc_model` `direct` and `lc_common_frequency_divisor` set to a value unequal to 1, will result in solving the problem with a system frequency, i.e., a frequency is only allowed in a solution, if it is the multiple of the system frequency `lc_common_frequency_divisor`. A value ≤ 0 will test any system frequency (except for 1) and output the best solution. For more information, see [9].

Aggregating the passengers per OD pair

Running

```
 R make lc-line-concept
```

with CK `lc_model` CV `direct_relaxation` results in solving the direct model, while aggregating the passengers per OD pair. This is a relaxation of the original model, see [3].

Multicriteria optimization

Setting CK `lc_direct_optimize_costs` to CV `true` will result in solving the direct model with a weighted sum, accounting for the line costs of the resulting line concept as well. As a weight factor, CK `lc_mult_relation` will be used.

3.3.3 Cost direct weighted sum

Executing

```
 R make lc-line-concept
```

with CK `lc_model` set to CV `mult_cost_direct` or CV `mult_cost_direct_relax` solve programs which are weighted sums between the cost model (Section 3.3.1) and the direct travelers model (Section 3.3.2). In the relaxed version (i.e.,

CV `mult_cost_direct_relax`) the vehicle capacity is not considered for each vehicle but only the aggregated capacity for each edge is considered. The capacity consideration can be turned off by setting CK `lc_mult_cap_restrict`. The weight can be set by CK `lc_mult_relation` where CV `0` refers to the direct travellers model and CV `1` to the cost model. The tolerance of feasibility, integrality and optimality can be set by CK `lc_mult_tolerance`. A time limit in seconds can be set by CK `lc_time_limit`, but it will only stop the computation if a feasible solution was already found. Otherwise the computation will continue until a feasible solution is found and stop then.

Additionally, there is the possibility to consider system frequencies, i.e., a common integer divisor for all frequencies. For this, set CK `lc_common_frequency_divisor` to something different than CV `1`. When setting it to a value smaller or equal to CV `0`, different prime values are tested as a system frequency and the best in terms of objective value is used as output. Note that testing prime numbers is enough for finding an optimal solution.

3.3.4 Traveling time without frequencies

Executing

```
 R make lc-line-concept
```

with CK `lc_model` CV `traveling_time_cg` solves the traveltime model as stated in [39, (LPMT1)] under the name (LPMT1). This model does not include line frequencies but only decides which lines are established. It routes all passengers over established lines and minimizes their resulting total travel time. Each established line incurs some cost and the total cost is bounded by a budget. This model is solved by a column generation procedure in which the passenger paths are generated throughout the column generation iterations. It is implemented as part of [17]. Various different methods exist in order to compute a feasible starting tableau. That is

- CK `lc_traveling_time_cg_cover` can be set to true or false and is a method to include passenger paths based on the idea that every edge is covered by at least one line.
- CK `lc_traveling_time_cg_k_shortest_paths` can be set to an integer value. This adds a number of shortest paths.

- `lc_traveling_time_cg_add_sol_1` can be set to true or false. The passenger paths which are based on the line concept (a file) given in `lc_traveling_time_cg_add_sol_1_name` are added.
- `lc_traveling_time_cg_add_sol_2` can be set to true or false. The passenger paths which are based on the line concept (a file) given in `lc_traveling_time_cg_add_sol_2_name` are added.
- `lc_traveling_time_cg_add_sol_3` can be set to true or false. The passenger paths which are based on the line concept (a file) given in `lc_traveling_time_cg_add_sol_3_name` are added.

Then the actual column generation procedure is started. Four different versions of constraints (corresponding to 1, 2, 3, 4) can be used which are set by `lc_traveling_time_cg_constraint_type`. Finally the following parameters are important for execution.

- `lc_traveling_time_cg_max_iterations`: This many column generation iterations are executed at most.
- `lc_traveling_time_cg_termination_value`: This is the gap in percent between lower and upper bound below which the best solution is returned.
- `lc_traveling_time_cg_weight_change_edge`: The weights of the transfer (change) edges in the Change&Go-Graph are determined by this value.
- `lc_traveling_time_cg_weight_od_edge`: The weights of the OD edges in the Change&Go-Graph are determined by this value.
- `lc_traveling_time_cg_relaxation_constraint`: boolean for additional relaxation constraint $y_i \forall i \in \mathcal{L}$
- `lc_traveling_time_cg_solve_ip`: if set to true the integer program corresponding to the final linear program should be solved in the last step to approximate an integer solution.

3.3.5 Traveling time with frequencies

Executing

```
 make lc-line-concept
```

with `lc_model` `traveling_time_mip` and `lc_traveling_time_mip_minimize` `"time"` solves the traveling time model with line frequencies as stated in [39, (LPMTF)]. In contrast to the formulation presented in this paper, flow variables belonging to the same origin are aggregated, as in [1]. It uses the solver chosen with `lc_solver` to solve the model. The following additional options can be given:

- `lc_traveling_time_mip_use_loads`: If this is set to true, then the upper and lower bounds on the frequency of service on each edge in the PTN given in the `default_loads_file` are respected. This corresponds to constraint (13) from [39, (LPMTF)] and a symmetric constraint for the lower bound. Otherwise, no bounds on the frequency are respected, i.e., the model only incorporates constraints (10)–(12) and (14) from the referenced model.
- `lc_traveling_time_mip_integer_flow`: Boolean to specify whether the computed passenger flows have to be integral.
- `lc_traveling_time_mip_integer_frequencies`: Boolean to specify whether the computed line frequencies have to be integral.
- `ean_model_weight_drive`: Determines the method used to estimate the driving time of a vehicle on an edge, based on the bounds given in the edge file, see Section 7.9.

- `CK` `ean_model_weight_wait`, `CK` `ean_default_minimal_waiting_time`, and `CK` `ean_default_maximal_waiting_time`: Determine the method used to estimate the waiting time of a vehicle at a station, see Section 7.9.
- `CK` `ean_change_penalty`, `CK` `ean_default_minimal_change_time`: Each transfer is charged with the sum of these two parameters.
- `CK` `lc_budget`: Allowed total cost of the chosen line concept. It is assumed that running a line with frequency f incurs a cost of f times the value specified in the `CK` `default_pool_cost_file`.
- `CK` `gen_passengers_per_vehicle`: Used to determine the total frequency needed to serve all passengers using a line along an edge.

3.3.6 Cost with traveling time bound

Executing

`R` `make lc-line-concept`

with `CK` `lc_model` `CV` `traveling_time_mip` and `CK` `lc_traveling_time_mip_minimize` `CV` `"cost"` solves a variant of the traveling time model with frequencies, in which the traveling time is bounded by `CK` `lc_traveling_time_mip_time_budget` and the cost of the line concept is minimized. Apart from that, it uses the same configuration parameters as the traveling time model with frequencies (except `CK` `lc_budget`).

3.3.7 Minchanges

Running

`R` `make lc-line-concept`

with `CK` `lc_model` `CV` `minchanges_ip` or `CV` `minchanges_cg` results in solving a program to minimize the number of passenger weighted transfers. For further reference see [15].

Integer program

The integer program corresponding to method `CV` `minchanges_ip` is

$$\text{(IP-LPT)} \quad \min \sum_{i,j \in V} \sum_{p \in \mathcal{P}_{CG}^{ij}} d_p c_p \quad (3.1)$$

$$\sum_{p \in \mathcal{P}_{CG}^{ij}} d_p \geq C_{ij} \quad \forall i, j \in V \quad (3.2)$$

$$\sum_{i,j \in V} \sum_{\substack{p \in \mathcal{P}_{CG}^{ij} \\ (e,l) \in p}} d_p \leq A f_l \quad \forall l \in \mathcal{L}, \forall e \in l \quad (3.3)$$

$$\sum_{\substack{l \in \mathcal{L} \\ e \in l}} f_l \leq f_e^{\max} \quad \forall e \in E \quad (3.4)$$

$$d_p \in \mathbb{N}_0 \quad \forall p \in \mathcal{P}_{CG} \quad (3.5)$$

$$f_l \in \mathbb{N}_0 \quad \forall l \in \mathcal{L} \quad (3.6)$$

Since paths of passengers have to be tracked in order to obtain their transfers, the model is based on the Change&Go-Graph CG proposed in [39]. Paths in the Change&Go-Graph are referred to as \mathcal{P}_{CG} . The number c_p then gives the number of transfers on a path $p \in \mathcal{P}_{CG}$. The variables d_p and f_l specify the number of passengers on path p and the frequency of line $l \in \mathcal{L}$, respectively.

The following parameters are used to execute the computation:

- `CK` `lc_minchanges_nr_ptn_paths` determines the maximum number of paths in the PTN on which passengers from each OD pair are allowed to travel. This ensures that also $|\mathcal{P}_{CG}|$ is bounded.
- `CK` `lc_minchanges_xpress_miprelstop`. This parameter is passed to the execution of Xpress and determines the gap (in percent) between lower and upper bound which has to be reached such that the best solution is returned.
- `CK` `lc_minchanges_nr_max_changes`. Since the number of paths in the Change&Go-Graph could become very large this parameter is used to bound them. Only paths which have less or equal transfers (changes) are considered. A value of 0 means that all paths are considered.
- `CK` `gen_passengers_per_vehicle`. This parameter corresponds to the A in constraint (3.3) and determines the vehicle capacity.

Column Generation procedure

In the column generation procedure the integer program (IP-LPT) is relaxed and initially only solved for a subset of all possible paths \mathcal{P}_{CG} . Throughout the column generation procedure paths which are likely to improve the current solution are determined and added to the program. The column generation procedure ends if no such paths can be found anymore. The problem which is solved in order to determine paths which are likely to improve the current solution is an all pairs shortest path problem. Since the correspondence of the solution of this problem to the primarily determined paths in the PTN, \mathcal{P}_G has to be checked, two different implementations can be used via `CK` `lc_minchanges_pricing_method`.

- `CV` `exact`: For each path $p \in \mathcal{P}_G$ the corresponding subgraph of CG is constructed and herein the all-pairs shortest path problem is solved.
- `CV` `heuristic`: The all-pairs shortest path problem is solved in the entire Change&Go-Graph CG for all pairs of nodes. It may happen that for a pair of nodes the shortest path does not correspond to a path in \mathcal{P}_G . In this case a warning is returned because the computation could be wrong. Still, this procedure is much faster since the Change&Go-Graph does not need to be constructed in every iteration.

Additional to the parameters in Section 3.3.7 the following parameters are of relevance.

- `CK` `lc_minchanges_nr_cg_paths_per_ptn_path`: For the starting tableau of the column generation procedure a set of initial paths has to be computed. This parameter determines how many paths in the Change&Go-Graph are computed for each path in the PTN.
- `CK` `lc_minchanges_cg_var_per_it`: Only at most this many variables are added in each column generation iteration.
- `CK` `lc_minchanges_max_reduced_costs_included_IP`: After the column generation only variables which have reduced costs less than or equal to this value are included in the final IP.

For more information on the model, see [15].

3.3.8 Game

Running

`R` `make lc-line-concept`

with `CK` `lc_model` set to `CV` `game` results in solving a game theoretic model where each line acts as a player and aims to minimize its own (expected delay). The delay is dependent on the traffic loads along its edges, i.e, a line tries to choose less-frequent edges. The algorithm uses a potential function to find a line plan at an equilibrium which is a system optimum. This line plan is computed by an integer program. For more information, see [40].

3.4 Ridepool Generation

A new ridepool is generated by running

```
R make rpool-ride-pool
```

There are three models to create a ridepool. The following parameters are needed for all models:

- **CK** rpool_min_edges Minimal number of edges in a ridepooling area. The default value is 1.
- **CK** rpool_max_edges Maximal number of edges in a ridepooling area. The default value is 4.
- **CK** rpool_costs_fixed Costs for one vehicle to operate in one ridepooling area.

A ridepooling area r is represented as a set of edges of the PTN. Those edges must form a strongly connected subgraph of the PTN. For each edge $e \in r$ there is a constant $\alpha_{r,e} \in \mathbb{R}_{\geq 0}$ specifying the fraction of the period length for which one single vehicle in r will probably be used on the edge e to serve the demand on this edge. By default, those values are computed by the following formula:

$$\alpha_{r,e} = N_e \cdot \frac{\text{CK period_length}}{\sum_{e \in r} \delta_e \cdot N_e},$$

where δ_e is the minimal duration of each edge and N_e is defined by

$$N_e = \begin{cases} L_e + 1 & \text{if } L_e \text{ is odd,} \\ L_e & \text{else,} \end{cases}$$

and L_e denotes $\lceil \frac{\text{load}(e)}{\text{CK rc_passengers_per_vehicle}} \rceil$.

Alternatively, one can also set them to an arbitrary value. Similar to lines, to each area will be assigned a number of vehicles operating in the area. The capacity of one ridepooling vehicle is controlled by **CK** rc_passengers_per_vehicle.

3.4.1 All

By setting **CK** rpool_model **CV** all, only one ridepooling area is generated consisting of all edges of the PTN.

3.4.2 Demand heuristic

For the usage of **CK** rpool_model **CV** demand_heuristic, the parameter **CK** rpool_load_factor is needed. The algorithm constructs recursively areas consisting of edges e with

$$\text{load}(e) \leq \text{CK rpool_load_factor} \cdot \text{CK gen_passengers_per_vehicle}, \quad (3.7)$$

i.e. the load of the edge is at most a fraction of **CK** rpool_load_factor of the capacity of a line vehicle. Ridepooling can be used to cover edges with small loads, otherwise one had to establish a line for those edges which would come with big costs.

The algorithm iterates over all edges e in the PTN. Condition (3.7) is checked and if it is fulfilled, a new area r is constructed consisting only of edge e . Now we check recursively for all edges that are incident to an edge in r if (3.7) holds and add them to r until the maximal number of edges is reached or no incident edge with (3.7) is found. The area r is added to the ride pool, if it contains at least **CK** rpool_min_edges edges.

3.4.3 Tree based

The tree bases heuristic can be used by setting `[CK] rpool_model [CV] tree_based`. It computes a maximal spanning tree with respect to the loads on the edges. This tree may represent a possible line concept, covering mainly edges with high demand. The algorithm tries to cover all non-tree edges by ride pool areas.

We iterate over all leaf nodes of the spanning tree. For a leaf node v , we create a new area r and add all incident non-tree edges to r . In a next step, we look for non-tree edges connecting the edges in r and add them to r . If the number of edges of r is between the lower and upper bound, r is added to the ride pool.

3.4.4 Adapt Line Pool

For the ridepooling problem it can be reasonable to adapt a given line pool. By running

`[R] rpool_adapt_lpool`

we add additional lines to the existing linepool. There are two methods how to trim existing lines in order to create the additional ones. By setting `[CK] lpool_adapt_method` to `[CV] trim_all`, we add for each existing line the three lines where `[CK] rpool_nb_edges_to_trim` are trimmed, respectively, from the left side, from the right side, and from both sides. By setting `[CK] lpool_adapt_method` to `[CV] trim_low_demand`, we use a heuristic that deletes only deletes edges with low demand from existing lines. `[CK] lpool_trim_capacity_factor` defines how many times the vehicle capacity is considered as low demand. Going along a line, we add the edges to the new line until we reach an edge of low demand. Then, the new line is added to the line pool.

3.5 Ride Concept

By running

`[R] make_rc-ride-concept`

a new ride concept and a new line concept are generated simultaneously. The following parameter values are needed as input:

- `[CK] rc_solver` specifies the solver to use the IP formulations.
- `[CK] rc_timelimit` Timelimit in seconds for the solver. Value of -1 means no restriction.
- `[CK] rc_mip_gap` MIP Gap for the solver. Value of -1 means no restriction.
- `[CK] rc_threads` Number of threads to use for the solver. Value of -1 means no restriction.
- `[CK] rc_write_lp_file` Whether to write the IP formulation as a file or not.
- `[CK] rc_passengers_per_vehicle` Capacity of a single ridepooling vehicle.

A ride concept consists of ridepooling areas with an assigned number of vehicles operating in this area.

3.5.1 Lineplanning and ridepooling with fixed demand factors

By setting `[CK] rc_model [CV] LPRP_ALPHA`, the following IP formulation is solved:

$$\text{minimize } \sum_{l \in \mathcal{L}^{\text{pool}}} c_l \cdot f_l + \sum_{r \in \mathcal{R}^{\text{pool}}} c_r \cdot v_r \quad (3.8a)$$

$$\text{subject to } t_e^{\min} \leq \sum_{l \in \mathcal{L}^{\text{pool}}: e \in l} A \cdot f_l + \sum_{r \in \mathcal{R}^{\text{pool}}: e \in r} B \cdot \alpha_{r,e} \cdot v_r \leq t_e^{\max} \forall e \in E, \quad (3.8b)$$

$$f_l, v_r \in \mathbb{N} \quad \forall l \in \mathcal{L}^{\text{pool}}, r \in \mathcal{R}^{\text{pool}} \quad (3.8c)$$

c_l and c_r are the costs of establishing the line l , or operating one vehicle in the area r , respectively. The variable f_l is the frequency of the line l and v_r is the number of vehicles in the area r . The first constraint ensures that the demand which is covered on each edge is at least the load of the edge (t_e^{\min}) and at most the upper (line) frequency bound of the edge multiplied by the (line) vehicle capacity (t_e^{\max}). A is the capacity of a line vehicle and B is the capacity of a ridepooling vehicle.

3.5.2 Lineplanning and ridepooling including demand factors as variables

By setting `CK rc_model CV LPRP_BETA`, the following IP formulation is solved:

$$\text{minimize } \sum_{l \in \mathcal{L}^{\text{pool}}} c_l \cdot f_l + \sum_{r \in \mathcal{R}^{\text{pool}}} c_r \cdot v_r \quad (3.9a)$$

$$\text{subject to } t_e^{\min} \leq \sum_{l \in \mathcal{L}^{\text{pool}}: e \in l} A \cdot f_l + \sum_{r \in \mathcal{R}^{\text{pool}}: e \in r} B \cdot \beta_{r,e} \leq t_e^{\max} \quad \forall e \in E, \quad (3.9b)$$

$$\sum_{e \in \mathcal{R}} d_e \cdot \beta_{r,e} \leq T \cdot v_r \quad \forall r \in \mathcal{R}^{\text{pool}}, \quad (3.9c)$$

$$\beta_{r,e} \in \mathbb{R}_{\geq 0} \quad \forall r \in \mathcal{R}^{\text{pool}}, e \in r, \quad (3.9d)$$

$$f_l, v_r \in \mathbb{N} \quad \forall l \in \mathcal{L}^{\text{pool}}, r \in \mathcal{R}^{\text{pool}} \quad (3.9e)$$

In this model, the edge-specific factors $\alpha_{r,e}$ from the first model are now determined as continuous variables $\beta_{r,e}$. The additional constraint links the number of vehicles to the factor $\beta_{r,e}$.

3.6 Timetabling

3.6.1 Modulo network simplex algorithms

There are different ways to use the Modulo Network Simplex Algorithm, depending on how to provide a starting solution:

- `CK tim_model CV ns_improve` It is assumed that `Timetable-periodic.tim` already contains a feasible starting solution; only improvement steps are taken.
- `CK tim_model CV csp_ns` A starting solution is found using Abscon; high reliability, small running times, but the starting solution quality is usually bad – see Section 3.6.2.
- `CK tim_model CV con_ns` A starting solution is found using constraint propagation; may take too long for some networks, but has good quality when it succeeds – see Section 3.6.3.
- `CK tim_model CV ns_cb` It is assumed that `Timetable-periodic.tim` already contains a feasible starting solution. It is improved with the network simplex. Afterwards, a cycle based IP is called. `CK tim_use_old_solution` needs to be set to `CV true` such that the network simplex solution is used as a starting solution for the IP.

There are two search procedures that may be further specified, one for local search and one for fundamental search for cuts, see [12]. The first is represented by the parameter `CK tim_nws_loc_search`, the second by `CK tim_nws_tab_search`.

The possible local search algorithms are:

- `CV SINGLE_NODE_CUT`.
The first improving single node cut that is found will be used. No further parameters have to be specified.

- **RANDOM_CUT.**
Single node cuts are chosen at random, ignoring whether they are improving or not. This will be repeated 10 times. This procedure is likely to give better results than `SINGLE_NODE_CUT`, but will take longer. No further parameters have to be specified.
- **WAITING_CUT.**
Cuts are chosen along each waiting edge cut. This will only improve `SINGLE_NODE_CUT` if the interval $[l_e, u_e]$ is especially small for waiting activities. No further parameters have to be specified.
- **CONNECTED_CUT.**
Cuts are found using a local search technique. This will be repeated up to 3 times. Usually yields the best results.

These are the possible fundamental search algorithms. Their setting will have the largest impact on the quality and time consumption of the solution.

- **TAB_FULL.**
All possible base exchanges are considered and the best one is chosen. This is usually quite time consuming but gives high quality results. No further parameters have to be specified. This may be considered as the default setting.
- **TAB_SIMPLE_TABU_SEARCH.**
As in `TAB_FULL`, all base exchanges are considered, but a tabu list gives the possibility to leave local optima again. Parameters are:
 - `tim_nws_ts_memory`. The length of the tabu list.
 - `tim_nws_ts_max_iterations`. The number of iterations that are allowed before searching for a local cut.

Because of the tabu list this algorithm is even slower than `TAB_FULL` but will seldom give better results because of the large number of neighbors in every step.

- **TAB_SIMULATED_ANNEALING.**
Base exchanges are chosen at random and used despite of being non-improving considering a steadily cooling temperature. Parameters are:
 - `tim_nws_sa_init`. The starting temperature.
 - `tim_nws_sa_cooldown`. The cooling factor < 1 .

This algorithm may improve `TAB_FULL` significantly. The time consumption is about the same.

- **TAB_STEEPEST_SA_HYBRID.**
A mix of `TAB_FULL` and `TAB_SIMULATED_ANNEALING`. This will usually yield the best results but takes longer than `TAB_FULL`. The same parameters are used as in `TAB_SIMULATED_ANNEALING`.
- **TAB_PERCENTAGE.**
A fast algorithm that decreases the quality of the solution only slightly. Parameters are:
 - `tim_nws_percentage`. An integer < 100 that gives the size of the search space.
- **TAB_FASTEST.**
Similar to `TAB_PERCENTAGE`. Parameters are
 - `tim_nws_min_pivot`. The minimum relative improvement a base exchange has to give.
 - `tim_nws_dyn_pivot`. The value by which the first parameter is multiplied if no cut fulfilling the criteria is being found.

For more information, see [13].

3.6.2 Constraint propagation

This is a way to find a feasible solution. The corresponding parameter is:

- `tim_model; "con_prop"`

A solution is found by fixing any event time, and propagating this information through the network, thus removing infeasible solutions. A backtracking procedure is used to fix times differently, if there is no feasible solution anymore.

Parameters are:

- `tim_cp_sortmode; "UP", "DOWN", "RANDOM"` Determines how event times are fixed. "UP" tries to tighten them as far as possible, while "DOWN" tries to relax them as far as possible. "RANDOM" chooses randomly from the set of locally feasible times, and often succeeds where the other two settings don't.
- `tim_cp_check_feasibility; true/false` If set to true, a heuristic check for feasibility is performed before the actual constraint propagation. This takes some time, but may help to determine infeasibility.

3.6.3 Abscon

Currently not included in the release version of LinTim.

To use Abscon, set

- `tim_model; "csp"`

The problem of finding a feasible timetable is then translated to a generic constraint satisfaction problem, and the third-party solver Abscon is started to find a feasible solution. If the problem is feasible, a feasible solution can be found relatively fast; however, its objective value tend to be worse than the one generated by constraint propagation. No parameters.

3.6.4 MATCH

To use MATCH, set

- `tim_model; "MATCH"`

A feasible timetable is found by a matching-merge heuristic. The details of this method can be looked up in [25].

3.6.5 PESP-IP

To use the pesp ip formulation, set

- `tim_model; "ip"`

This will try to solve an integer programming model of the periodic timetabling problem, see [33]. The IP model is implemented in Xpress and Gurobi.

For Gurobi, a solution limit, a best bound stop value, starting solution procedure and a MIPFocus are implemented (see Gurobi documentation):

- `tim_pesp_ip_solution_limit`
- `tim_pesp_ip_best_bound_stop`
- `tim_pesp_ip_mip_focus`
- `tim_use_old_solution`

For all parameters the default value of 0 will disable the respective option.

For more information on the model, see [41].

3.6.6 Cycle-based IP

To use the cycle based mip formulation, set

- `tim_model; "cb_ip"`

This will try to solve a cycle based integer programming model of the periodic timetabling problem, see [33]. You can set a time limit, a thread limit and a desired gap by setting

- `tim_mip_gap`
- `tim_timelimit`
- `tim_threads.`

The following parameter is for a (heuristic) preprocessing step where edges with few passengers are removed:

- `tim_pesp_cb_passenger_cut.`

Additionally for Gurobi, a solution limit, a best bound stop value, and a MIPFocus are implemented (see Gurobi documentation):

- `tim_pesp_cb_solution_limit`
- `tim_pesp_cb_best_bound_stop`
- `tim_pesp_cb_mip_focus_stop.`

For all parameters the default value of 0 will disable the respective option.

For more information on the model, see [41].

3.6.7 Phase 1 simplex

To use the phase 1 simplex method, set `tim_model` to `phase-one`. The idea of this model is to construct an auxiliary PESP instance that is easy to solve and a feasible solution can be converted into a feasible solution for the original problem or prove the infeasibility of the original problem. For more information on this procedure, see [14].

3.6.8 Adaptions

Fixed times

Some timetabling models are able to handle additional restrictions on the events, namely an additional interval for each one. Note that this interval may only include one value, fixing some events to a specific time.

To use this feature, set `tim_respect_fixed_times` to `true` and add `filename_tim_fixed_times` (`timetabling/Fixed-timetable-periodic.tim`) for the additional information.

3.7 Tariff Planning

Running

```
 make taf-tariff
```

will determine a new tariff minimizing the deviation from reference prices in a prespecified model. The available models are the following:

- `taf_model` `flat`, optimization model determining a flat tariff

- taf_model CV beeline_distance, optimization model determining an affine beeline distance tariff
- taf_model CV network_distance, optimization model determining an affine network distance tariff
- taf_model CV zone, optimization model determining a zone tariff

All models optimize prices such that the new tariff is close to the reference prices, which can e.g. be obtained from a former tariff, given in filename_tariff_reference_price_matrix_file (basis/Reference-Price-Matrix.giv).

3.7.1 General Remarks

In tariff planning we only consider node pairs with different nodes and call them non-trivial OD pairs:

$$D := (V \times V) \setminus \{(v, v) : v \in V\}.$$

Tariff planning always produces a price matrix file filename_tariff_price_matrix_file (tariff/Price-Matrix.taf) as output. Prices for trivial OD pairs, i.e. pairs with the same origin and destination node, are set to zero.

Objective function and weight options

In each of the available tariff models there are two options for the objective function. For each option one can choose one out of three possible weight options. The objective function is determined by taf_objective and taf_weights_objective.

If taf_objective has the value

- sum_absolute_deviation, the objective function is the weighted sum of absolute deviations between the new prices and the reference prices (see equation (3.10)),
- max_absolute_deviation, the objective function is the weighted maximum absolute deviation between the new prices and the reference prices (see equation (3.11)).

If taf_weights_objective has the value

- od, the price deviations are weighted by the OD data,
- unit, the price deviations have weight 1,
- reference_inverse, the price deviations are weighted by the inverse of the given reference prices.

This results in one of the two objectives

$$\sum_{d \in D} C_d |r_d - \pi_d|, \quad (3.10)$$

$$\max_{d \in D} C_d |r_d - \pi_d| \quad (3.11)$$

with the above defined set D of all non-trivial OD pairs and reference prices r_d for all non-trivial OD pairs $d \in D$. The new prices that have to be computed are denoted by π_d . The weights C_d refer either to the OD data if taf_weights_objective od, or is $C_d = 1$ for all $d \in D$ if unit, or is $C_d = \frac{1}{r_d}$ for all $d \in D$ if reference_inverse. Both objective functions are applied in a linearized form in the programming formulations.

Routing options

In the tariff models `CV` distance and `CV` zone the prices are optimized with respect to given passenger paths in the PTN. Which paths are used is determined by `CK` `taf_routing_generation`:

- `CV` `fastest-paths`, a new routing using fastest paths with respect to the lower time bounds on the edges is created,
- `CV` `read-all`, a routing for all non-trivial OD pairs given in `CK` `filename_routing_ptn_input` (`Fi` `basis/Routing-ptn.giv`) is read and used,
- `CV` `read-partial-fill`, a partial routing given in `CK` `filename_routing_ptn_input` (`Fi` `basis/Routing-ptn.giv`) is read. Unspecified paths for non-trivial OD pairs are filled with fastest paths with respect to the lower time bounds on the edges.

In all cases the used routing is stored to `CK` `filename_routing_ptn_output` (`Fi` `basis/Routing-ptn.giv`).

Solver options

The following parameters control the behavior of the solver in all models.

- `CK` `taf_solver` determines the solver to be used. Note that currently only Gurobi is supported.
- `CK` `taf_threads` determines the maximum number of threads to use for the solver (-1=use default value, i.e., no restriction). Note that this will only be used for a possible solver integration of the chosen model, not for the rest of the algorithm.
- `CK` `taf_timelimit` sets a time limit for the solver in seconds (-1=use default value).
- `CK` `taf_write_lp_file` determines whether to write the lp file of the model to solve.
- `CK` `taf_mip_gap` sets the MIP optimization gap for the solver. The solver will terminate with an optimal solution if the gap between lower and upper objective bound is less than this value times the absolute value of the incumbent objective value.

3.7.2 Flat Tariff

Running

```
R make taf-tariff
```

with `CK` `taf_model` `CV` `flat` determines a new flat tariff, i.e. a fixed price f for all paths W in the given PTN.

The flat tariff model solves the following linear program to optimality:

$$\begin{aligned} \min \quad & g(f) \\ \text{s.t.} \quad & f \geq 0, \end{aligned}$$

where the objective function g is determined by `CK` `taf_objective` and `CK` `taf_weights_objective` as described in Section 3.7.1, i.e.

$$g(f) = \begin{cases} \sum_{d \in D} C_d |r_d - f| & \text{for } \text{CV} \text{ sum_absolute_deviation,} \\ \max_{d \in D} C_d |r_d - f| & \text{for } \text{CV} \text{ max_absolute_deviation.} \end{cases}$$

The computed flat price is written as `CK` `taf_flat_price` to the State-Config-file `CK` `filename_state_config` (`Fi` `basis/State-Config.cnf`). The following file is produced as output:

- `CK` `filename_tariff_price_matrix_file` (`Fi` `tariff/Price-Matrix.taf`), price matrix containing the price f for each OD pair.

3.7.3 Distance Tariffs

Running

`R` `make taf-tariff`

with `CK` `taf_model` `CV` `beeline_distance` or `CV` `network_distance` results in determining a new distance-based tariff, i.e. the price $p(W)$ for travelling along the path W in the given PTN is determined by $p(W) = f + l(W) \cdot p$ where $f \geq 0$ are fixed costs and $p \geq 0$ is a price factor that is multiplied with the distance $l(W)$, which is either the Euclidean distance between the start and the end station of the path or the sum of all edge lengths of the path.

The distance-based model solves the following program to optimality:

$$\begin{aligned} \min \quad & g(f, p) \\ \text{s. t.} \quad & f, p \geq 0, \end{aligned}$$

where the objective function g is determined by `CK` `taf_model`, `CK` `taf_objective` and `CK` `taf_weights_objective`. This results in one of the two objectives:

$$g(f) = \begin{cases} \sum_{d \in D} C_d |r_d - (f + l_d \cdot p)| & \text{for } \text{CV} \text{ sum_absolute_deviation,} \\ \max_{d \in D} C_d |r_d - (f + l_d \cdot p)| & \text{for } \text{CV} \text{ max_absolute_deviation.} \end{cases}$$

as described in Section 3.7.1. The value of l_d is determined by `CK` `taf_model`:

- `CV` `beeline_distance`, the distance l_d of a non-trivial OD pair $d \in D$ is calculated as the Euclidean distance in km in the plane,
- `CV` `network_distance`, the distance l_d of a non-trivial OD pair $d \in D$ is the length (in terms of edge length) of the path of this node pair in the routing. The parameter `CK` `taf_routing_generation` specifies how the routing is determined as described in Section 3.7.1.

The computed fixed price f and the price factor p are written as `CK` `taf_fixed_costs` and `CK` `taf_factor_costs` to the State-Config-file `CK` `filename_state_config` (`Fi` `basis/State-Config.cnf`). The following file is produced as output:

- `CK` `filename_tariff_price_matrix_file` (`Fi` `tariff/Price-Matrix.taf`), price matrix containing the prices for each OD pair.

3.7.4 Zones

Running

`R` `make taf-tariff`

with `CK` `taf_model` `CV` `zone` determines a new zone tariff by determining zones and a price list. The set of zones is a partition of the set of nodes of the PTN. Prices are given as a price list that assigns a price to the number of traversed zones. The price of a path depends on the number of traversed zones on that path. We say that a zone is traversed if a node of this zone is part of the path, in particular the zones of the start node and end node of a path are traversed. The price for traversing more zones than the maximal specified number in the price list is just the price for traversing the maximal specified number of zones.

It is also possible to determine only a price list for given zones or only zones for a given price list.

The parameters are

- `CK` `taf_zone_counting`: Specifies how the number of traversed zones is counted. If `CV` `single`, then each different zone is only counted once. If `CV` `multiple`, then a zone is counted each time that it is entered. For example consider the path from station 1 to station 3 in the PTNs with zones given in Figure 3.1. In the multiple counting case, the number of traversed zones is 3 in both PTNs. In the single counting case, the number of traversed zones is 3 in the left PTN and 2 in the right PTN because there are only two different zones on the path and the reentry is not counted.



Figure 3.1: PTNs with zones

- $\boxed{\text{CK}}$ `taf_zone_n_zones` integer number specifying the maximum number of zones,
- $\boxed{\text{CK}}$ `taf_zone_enforce_all_zones` boolean, determines whether exactly $\boxed{\text{CK}}$ `taf_zone_n_zones`-many zones ($\boxed{\text{CV}}$ `true`) or at most that many zones ($\boxed{\text{CV}}$ `false`) must be determined,
- $\boxed{\text{CK}}$ `taf_zone_connected` boolean, specifies whether the subgraph of a zone, induced by the nodes assigned to the zone, needs to be connected (in case of a directed graph it is weakly connected),
- $\boxed{\text{CK}}$ `taf_zone_enforce_no_elongation` boolean, determines whether the no-elongation property must be satisfied. This property ensures, that it is never cheaper for passengers to buy a ticket for more zones than they actually want to travel through. Let p_k be the price of a path that traverses k zones. The no-elongation property is satisfied if it holds that

$$p_k \leq p_{k+1} \quad \text{for all } k \in \{1, \dots, (\boxed{\text{CK}} \text{ taf_zone_n_zones}) - 1\}.$$

- $\boxed{\text{CK}}$ `taf_zone_enforce_no_stopover` boolean, determines whether the no-stopover property must be satisfied. This property ensures that it is never cheaper for a passenger to buy two separate tickets for one journey and combine them instead of buying one ticket for the whole journey. Let again p_k be the price of a path that traverses k zones. The no-stopover property in the case of single counting is

$$p_k \leq p_i + p_j \quad \text{for all } k \in \mathbb{N}_{\geq 1}, i, j \in \{1, \dots, k\} \text{ with } i + j \geq k + 1.$$

In the case of multiple counting the property holds if

$$p_k \leq p_i + p_j \quad \text{for all } k \in \mathbb{N}_{\geq 1}, i, j \in \{1, \dots, k\} \text{ with } i + j = k + 1.$$

- $\boxed{\text{CK}}$ `taf_zone_symmetry_breaking`, determines which symmetry breaking model (see below) should be used. Possible values are $\boxed{\text{CV}}$ `A`, $\boxed{\text{CV}}$ `B` and $\boxed{\text{CV}}$ `NONE`.
- $\boxed{\text{CK}}$ `taf_routing_generation`, determines which routing should be used, see Section 3.7.1.
- $\boxed{\text{CK}}$ `taf_zone_only_zones` boolean, specifies whether only zones based on given prices must be computed,
- $\boxed{\text{CK}}$ `taf_zone_only_prices` boolean, specifies whether only prices based on given zones must be computed.

The objective in the zone model is to minimize the objective function $g(\pi)$ such that the above mentioned constraints are satisfied. The objective function is determined by $\boxed{\text{CK}}$ `taf_objective` and $\boxed{\text{CK}}$ `taf_weights_objective` as described in Section 3.7.1, i.e.

$$g(\pi) = \begin{cases} \sum_{d \in D} C_d |r_d - \pi_d| & \text{for } \boxed{\text{CV}} \text{ sum_absolute_deviation,} \\ \max_{d \in D} C_d |r_d - \pi_d| & \text{for } \boxed{\text{CV}} \text{ max_absolute_deviation.} \end{cases}$$

Here π_d refers to the price of OD pair d for travelling along the path given in the routing determined by $\boxed{\text{CK}}$ `taf_routing_generation` (see Section 3.7.1).

The results are written to:

- `filename_tariff_price_matrix_file` (`tariff/Price-Matrix.taf`), containing the prices,
- `filename_tariff_zone_file` (`tariff/Zones.taf`), containing the assignment of stops to zones and
- `filename_tariff_zone_price_file` (`tariff/Zone-Prices.taf`), price list of the zone tariff.

Symmetry Breaking

When determining a zone tariff some feasible solutions may only vary in name. Therefore symmetry breaking constraints can be introduced to the MILP solving the problem.

Let $x_{vz} = 1$ if and only if the stop with stop-ID v is allocated to the zone with zone-ID z and 0 else. Then the following constraints can be introduced:

$$x_{vz} = 0 \quad \text{for all } v \in \{1, \dots, \min\{n, N\}\}, z \in \{v+1, \dots, N\} \quad (3.12)$$

$$x_{vz} \leq \sum_{k=1}^{v-1} x_{k,z-1} \quad \text{for all } v \in \{3, \dots, n\}, z \in \{3, \dots, N\} \quad (3.13)$$

$$\sum_{v \in V} x_{vz} \geq \sum_{v \in V} x_{v,z+1} \quad \text{for all } z \in \{1, \dots, N-1\} \quad (3.14)$$

Here, N denotes the maximum number of zones (`taf_zone_n_zones`) and n is the number of nodes in the PTN.

The first constraint (3.12) ensures that the i -th stop can only be in the first i zones. The second one (3.13) ensures that a stop is only allowed in a zone if a node with a smaller stop-ID is in the zone with the next smaller zone-ID. The third one (3.14) orders the zones descending by size such that the one with the smallest zone-ID is the biggest in terms of number of nodes.

The parameter `taf_zone_symmetry_breaking` determines which one of them will be used. Three options are available:

- A, implementing (3.12) and (3.13). This seems to be the fastest.
- B, implementing (3.14). This seems to be slower.
- NONE, no symmetry breaking constraints are applied.

Only Prices

If only prices should be optimized for given zones, `taf_zone_only_prices` (boolean) must be set to `true`. By default it is `false`.

If `true`, the same MILP as described above is solved, but the zones are fixed. Therefore a zone file `filename_tariff_zone_file` (`tariff/Zones.taf`) must be given, otherwise the algorithm fails.

Only Zones

If only zones should be optimized for given prices, `taf_zone_only_zones` (boolean) must be set to `true`. By default it is `false`.

If `true`, the same MILP as described above is solved, but the prices for travelling through a certain number of zones are fixed. Therefore a zone-price file `filename_tariff_zone_price_file` (`tariff/Zone-Prices.taf`) must be given, otherwise the algorithm fails.

3.8 Vehicle Scheduling

The vehicle scheduling step can be invoked via

```
[R] make vs-vehicle-schedules
```

It assumes that there is an aperiodic Event-Activity Network with a given timetable for the aperiodic events and a set of trips to cover, which can be generated from a periodic timetable by the auxiliary rollout algorithm (see Section 4.9).

3.8.1 Mdm1

Running

```
[R] make vs-vehicle-schedules
```

with the `[CK] vs_model` set to `[CV] MDM1` will result in running a model minimizing the number of vehicles used in the vehicle schedule. For two consecutive trips the last station of the first trip has to be equal to the first station of the second trip. A depot, given by `[CK] vs_depot_index`, is considered. For more information on the model, see [2].

3.8.2 Mdm2

Running

```
[R] make vs-vehicle-schedules
```

with the `[CK] vs_model` set to `[CV] MDM2` will result in running a model that is equivalent to `[CV] MDM1`, except that no depot is considered. For more information on the model, see [2].

3.8.3 Assignment model

Running

```
[R] make vs-vehicle-schedules
```

with the `[CK] vs_model` set to `[CV] ASSIGNMENT_MODEL` will result in running a model minimizing the overall costs, considering vehicle costs (`[CK] vs_vehicle_costs`) and empty meters costs (given by the respective distance in time). A depot, given by `[CK] vs_depot_index`, can be considered.

Two consecutive trips can have different end and start stations respectively. Whether they can be connected relies on the end time of trip one, the start time of trip two, the distance between the two respective stations (in terms of minimal running times on shortest path) and a minimal turnover time (`[CK] vs_turn_over_time`). Note that the turnover time is not a simple restriction on the time between two connected consecutive trips, but includes the time needed to travel to the later station, i.e., it is the designated time the vehicle needs to be available at the later station before departing again.

For more information on the model, see [2].

3.8.4 Transportation model

Running

```
[R] make vs-vehicle-schedules
```

with the `[CK] vs_model` set to `[CV] TRANSPORTATION_MODEL` will result in running a model minimizing the overall costs, considering vehicle costs by driving to/from the depot, given by `[CK] vs_depot_index`, and (fixed) penalty costs `[CK] vs_penalty_costs` for not giving service on a trip. For more information on the model, see [2].

3.8.5 Network flow model

Running

```
R make vs-vehicle-schedules
```

with the `CK vs_model` set to `CV NETWORK_FLOW_MODEL` will result in running a model minimizing the overall costs considering both vehicle and empty meters costs. A depot, given by `CK vs_depot_index`, is considered. The number of vehicles can be bounded. For more information on the model, see [2].

3.8.6 Canal model

Running

```
R make vs-vehicle-schedules
```

with the `CK vs_model` set to `CV CANAL_MODEL` will result in running a more detailed version of `CV ASSIGNMENT_MODEL` incorporating the actual waiting times at every node and furthermore the considered period can be extended. Thus, each station can be regarded as a depot if trains from one day wait at the station for a service from that station the next day. Also, depot and maintenance decisions for locations which are farther away from the actual station can be taken. The minimal turnover time (`CK vs_turn_over_time`) will be respected. For more information on the model, see [43].

3.8.7 Line-based

Running

```
R make vs-vehicle-schedules
```

with the `CK vs_model` set to `CV LINE-BASED` will result in running a model based on line planning only. This model runs with the `CK vs_line_based_method` set to `CV 4`, `CV 3` or `CV 2` and `CK vs_line_based_alpha` set to `CV 0.3`. Here the `CK vs_line_based_method` describes the program type and the `CK vs_line_based_alpha` describes the value of α . For more information on the model, see [18].

3.8.8 Simple

Running

```
R make vs-vehicle-schedules
```

with the `CK vs_model` set to `CV SIMPLE` will result in a homogeneous vehicle schedule, i.e., all vehicles will serve only one line, back and forth.

3.8.9 IP model

Running

```
R make vs-vehicle-schedules
```

with the `CK vs_model` set to `CV IP` will result in a simple ip model to determine a cost efficient vehicle schedule. Trips are determined compatible, if the shortest path w.r.t. the lower bounds is sufficient to serve the trips after each other. A depot, given by `CK vs_depot_index` can be considered. Currently, only `CV GUROBI` is allowed as `CK vs_solver`. A time limit for the ip model can be set via `CK vs_timelimit`, where `CV -1` disables this option. The cost of a vehicle is determined using `CK vs_vehicle_costs` and the cost of an empty trip by `CK vs_eval_cost_factor_empty_trips_length` and `CK vs_eval_cost_factor_empty_trips_duration`. The minimal turnover time (`CK vs_turn_over_time`) will be respected. For more information on the model, see [2].

3.9 Delay Management

The delay-management step can be invoked via

```
[R] make dm-disposition-timetable
```

It assumes that there is an aperiodic Event-Activity Network with a given timetable for the aperiodic events, which can be generated from a periodic timetable by the auxiliary rollout algorithm (see Section 4.9), and some primary delays on events and/or activities (see Section 4.10). The lower bounds on the drive, wait (dwell) and fixed-circulation activities can be automatically reduced to account for a globally applied buffer that is contained in the lower bounds but may be exploited in case of delays. To this end, the parameter `[CK] DM_lower_bound_reduction_factor` can be set to a value below `[CV] 1.0`.

Note that during all these steps – in contrast to preceding planning steps like line planning or periodic timetabling – time intervals are measured in seconds, points in time in seconds since 0:00. E.g., if an activity has a lower bound of 60, this means 60 seconds, and if the time of an event is 28 800, this means 08:00 a.m.

The following parameters are used by all methods:

- `[CK] DM_verbose`: enable verbose output
- `[CK] DM_enable_consistency_checks`: enable (time-consuming) consistency checks of input data and results
- `[CK] DM_debug`: enable debugging output (also enables verbose output and consistency checks)

3.9.1 Propagate

The mere propagation of delays to produce a feasible disposition timetable is done when `[CK] DM_method` is set to `[CV] propagate`. After applying the given delays on events and on the lower bounds on activity durations, the (rolled-out) events are traversed in a topological sorting. Upon visit of each event, its time becomes fixed (since, due to the topological sorting, all events taking place earlier have been fixed before) and its successor events (targets of outgoing activities) are delayed as much as necessary to fulfill the lower bound on the duration of the respective activity.

During this propagation procedure, change activities can be cut off (so that delays will not propagate along them) based on a maximum waiting time: If the target event of a change activity would be delayed by more than `[CK] DM_propagate_maxwait` seconds, then this change activity is not respected at all. If all change activities shall be maintained, this parameter must be set to a very large value (e.g. the duration of the time horizon according to the rollout parameters, in seconds).

Furthermore, the headway constraints can be swapped around in those cases where the train that was originally scheduled first is so late that the train that was originally scheduled to go second can actually go first without affecting the train originally scheduled first. To enable this swapping of headways, `[CK] DM_propagate_swapHeadways` must be set to `[CV] true`.

3.9.2 Integer-Linear-Programming based methods

The aim of delay management is to react to delays in such a way that the effect on the passengers is minimal. To this end, one has to decide for each connection whether it should be maintained or not (i.e., if a connecting train waits for a delayed feeder train or not) and for each pair of trains using the same piece of track which train should go first. The delay management problem is for example described in [27]. The following parameters are used by all delay management algorithms:

- `[CK] DM_solver`: Defines which MIP solver should be used. Possible choices are Gurobi and Xpress. Please note that your environment (e.g. the CLASSPATH variable) has to be set up properly.
- `[CK] DM_solver_time_limit`: Time limit for the MIP solver in seconds – after this time, the solver is interrupted and the best solution found so far is used. If set to 0, no time limit is used.

- **CK** `DM_lower_bound_reduction_factor`: Describes how much buffer time is included in the minimal duration of the activities in the event-activity network. The lower bounds read from the input are multiplied with this number, so setting **CK** `DM_lower_bound_reduction_factor` to 1 does not change the lower bounds, while setting it to a value in $]0, 1[$ reduces the lower bound of all activities.

The variable **CK** `DM_method` defines which algorithm should be used to solve the delay management problem:

- CV** **DM1**: Computes an optimal solution of the MIP formulation (DM1) presented in [34, 35]. This is the slowest algorithm provided. To perform the calculation, the rollout must have been done where the parameter **CK** `rollout_passenger_paths` is set to **CV** `true` since the algorithm minimizes the delays on the passenger paths given in **CK** `default_passenger_paths_file`.
- CV** **DM2**: Computes an optimal solution of the MIP formulation (DM2) presented in [34, 35]. This is an approximation for (DM1) and a bit faster but still far slower than the other algorithms.
- CV** **DM2-pre**: The same as **CV** `DM2`, but with a preprocessing step. Computes an optimal solution of the MIP formulation (DM2) after applying Algorithm 3.2 from [27, p. 38] for reducing the size of the event-activity network. For more information, see [34, 35].
- CV** **FSFS**: “First scheduled, first served” – fixes the forward headways, deletes the backward headways, and solves the resulting uncapacitated delay management problem with fixed headways to optimality using DM1 or DM2, as specified in **CK** `DM_opt_method_for_heuristic`, see Algorithm 4.1 in [27, p. 56]. For more information, see [27, 28]. *Heuristic algorithm – might not find the global optimum.*
- CV** **FRFS**: “First rescheduled, first served” – fixes the headways according to the optimal solution of the corresponding uncapacitated delay management problem, then solves the resulting uncapacitated delay management problem with fixed headways to optimality using **CV** `DM1` or **CV** `DM2`, as specified in **CK** `DM_opt_method_for_heuristic`, see Algorithm 4.2 in [27, p. 57]. For more information, see [27, 28]. *Heuristic algorithm – might not find the global optimum.*
- CV** **EARLYFIX**: Similar to **CV** `FRFS` – but also fixes the changing activities according to the solution of the corresponding uncapacitated delay management problem by using **CV** `DM1` or **CV** `DM2`, as specified in **CK** `DM_opt_method_for_heuristic`, see Algorithm 4.3 in [27, p. 57]. For more information, see [27, 28]. *Heuristic algorithm – might not find the global optimum. Note that **CV** `FRFS` is always at least as good as this method [27, Lemma 4.5], while this method might be faster on instances with many changing activities.*
- CV** **PRIORITY**: Similar to **CV** `FSFS` – but also fixes the “most important” connections (the variable **CK** `DM_method_prio_percentage` defines how many percent of all connections should be maintained), see Algorithm 4.4 in [27, p. 57]. For more information, see [27, 28]. *Heuristic algorithm – might not find the global optimum. Uses **CV** `DM1` or **CV** `DM2`, as specified in **CK** `DM_opt_method_for_heuristic` for optimization. Note that **CV** `FSFS` is always at least as good as this method [27, Lemma 4.6], while this method might be faster on instances with many changing activities.*
- CV** **PRIOREPAIR**: Fixes the connections according to their weights like **CV** `PRIORITY`, relaxes the headway constraints, and solves the resulting problem using **CV** `DM1` or **CV** `DM2`, as specified in **CK** `DM_opt_method_for_heuristic`. Then it uses this solution to fix the headways and solves the problem again (again **CV** `DM1` or **CV** `DM2`) (see Algorithm 4.7 in [27, p. 68]). For more information, see [27, 28]. *Heuristic algorithm – might not find the global optimum.*
- CV** **best-of-all**: Runs **CV** `FSFS`, **CV** `FRFS` and **CV** `PRIOREPAIR` consecutively and takes the best solution. Due to [27, Lemma 4.5] and [27, Lemma 4.6], it’s sufficient to run **CV** `FSFS`, **CV** `FRFS`, and **CV** `PRIOREPAIR` and to ignore **CV** `EARLYFIX` and **CV** `PRIORITY`. Uses **CV** `DM1` or **CV** `DM2`, as specified in **CK** `DM_opt_method_for_heuristic` for optimization. If

`CK DM_best_of_all_write_objectives` is set to `CV true`, this will output all objective values of the different methods into

`CK filename_dm_best_of_all_objectives` (`Fi statistic/dm_objectives.sta`). For more information, see [28]. *Heuristic algorithm – might not find the global optimum.*

`CV PASSENGERFIX`: Uses a IP to fix the headways of passenger paths with the most passenger weight sum possible without contradictions and solves the following smaller problem with `CV DM1`. Note that all headways on a path get fixed if and only if no headway contradicts the earlier decisions. Otherwise no headway gets fixed. Same requirement as `CV DM1`. The IP is very big and slow!

`CV PASSENGERPRIOFIX`: A heuristic for the IP of `CV PASSENGERFIX`, fixes the headways of the first `CK DM_method_prio_percentage` percent of the passenger paths sorted by weight. Fixes any headway for a path only if this is possible without contradiction to the previous paths. After that, it solves the smaller problem with `CV DM1`. Same requirement as this method.

`CV FIXFSFS`: First uses the fixing method of `CV PASSENGERPRIOFIX` on as many paths as possible, again sorted by weight. After that it uses the fixing method of `CV FSFS` to fix the remaining headways. After that, it solves the reduced problem with `CV DM1` with the same requirement.

`CV FIXFRFS`: Like `CV FIXFSFS`, just uses the fixing method of `CV FRFS` instead of `CV FSFS`

3.10 Integrated Planning

The common parameters for all integrated programs are the following. Whether these parameters are used is dependent on the specific problems.

`CK int_max_threads` The maximal number of cpu threads used for optimization

`CK int_factor_travel_time` The objective factor for the travel time

`CK int_factor_drive_time` The objective factor for the drive time of the passengers

`CK int_factor_transfer_time` The objective factor for the transfer time of the passengers

`CK int_factor_wait_time` The objective factor for the waiting time of the passengers

`CK int_factor_penalty_time_slice` The penalty for changing time slices for the passengers. Only applicable on models respecting time slices. Only applicable for models with passenger routing.

`CK int_time_slices` The number of time slices to use. Only applicable for models with passenger routing.

`CK int_number_of_periods` The number of periods to consider the vehicle schedule for. Lines will not be cut off at the end of the planning period. Only applicable for models with vehicle scheduling.

`CK int_restrict_to_system_frequency` Whether to use a system frequency, i.e., a common divisor for all frequency values. Only applicable for models with line planning.

`CK int_system_frequency` The value for the system frequency, i.e., the common divisor for all frequency values. Only applicable for models with line planning.

`CK int_check_lower_frequencies` Whether the model should respect the lower frequency bounds, i.e., the minimal number of times edges in the public transport network need to be covered. Only applicable for models with line planning.

`CK int_check_upper_frequencies` Whether the model should respect the upper frequency bounds, i.e., the maximal number of times edges in the public transport network may be covered. Only applicable for models with line planning.

int_set_starting_timetable Whether to set the starting values for timetabling. Only applicable for models not containing line planning.

int_solver_type The solver to use.

3.10.1 Integrated timetabling and passenger routing

An implementation of the integrated periodic timetabling and passenger routing problem. For details on the model, see [30].

tim_pass_use_preprocessing Whether to use an exact preprocessing method to reduce the problem size before optimization.

tim_pass_use_cycle_base Whether to use a cycle-base formulation. This is normally much faster.

tim_pass_restrict_transfer_stations Whether to use an auxiliary IP to restrict the transfer stations. This method is only exact if all drive- and wait-activities are fixed.

tim_pass_add_fixed_passenger_paths Whether to add the non-routed passengers as fixed weights to the model.

tim_pass_number_of_routed_od_pairs The number of routed od pairs.

tim_pass_choose_routed_od_pairs How to choose the routed od pairs. The following methods are possible:

POTENTIAL Choose the od pairs with the most potential, i.e., compute the shortest path w.r.t. lower bounds on the EAN, evaluate these paths w.r.t. the difference of upper and lower bound on each activity and weight the result by the number of passengers of the od pair.

LARGEST_WEIGHT Choose the od pairs with the largest weight.

SMALLEST_WEIGHT Choose the od pairs with the smallest weight.

LARGEST_WEIGHT_WITH_TRANSFER Choose the od pairs with the largest weight that additionally have at least one transfer in their shortest path w.r.t. the lower bounds on the EAN.

LARGEST_DISTANCE Choose the od pairs with the largest euclidian distance.

DIFF Similar to **POTENTIAL** but without the additional scaling by the number of passengers.

RANDOM Random.

tim_pass_time_limit The time limit for the optimization.

tim_pass_mip_gap The mip gap for the optimization.

tim_pass_write_lp_output Whether to write the lp output. Will additionally compute an IIS for infeasible programs.

3.10.2 Integrated timetabling and aperiodic vehicle scheduling

Solve the integrated periodic timetabling and aperiodic vehicle scheduling problem. Includes passenger routing for the timetabling step. For more information, see [30].

tim_veh_allow_empty_trips Whether to allow empty trips in the vehicle schedule.

tim_veh_use_lower_bound Whether to include an additional lower bound on the objective function.

tim_veh_time_limit The time limit for the optimization.

tim_veh_mip_gap The mip gap for the optimization.

tim_veh_write_lp_output Whether to write the lp output. Will additionally compute an IIS for infeasible programs.

3.10.3 Integrated line planning and timetabling

Solve the integrated line planning and periodic timetabling problem. Includes passenger routing for the timetabling stage. For more information, see [30].

- lin_tim_pass_use_preprocessing** Whether to use an exact preprocessing method to reduce the problem size before optimization.
- lin_tim_pass_add_fixed_passenger_paths** Whether to add the non-routed passengers as fixed weights to the model.
- lin_tim_pass_number_of_routed_od_pairs** The number of routed od pairs.
- lin_tim_pass_factor_line_cost** The factor for the line costs.
- lin_tim_pass_time_limit** The time limit for the optimization.
- lin_tim_pass_mip_gap** The mip gap for the optimization.
- lin_tim_pass_write_lp_output** Whether to write the lp output. Will additionally compute an IIS for infeasible programs.
- lin_tim_pass_choose_routed_od_pairs** How to choose the routed od pairs. The following methods are possible:
 - LARGEST_WEIGHT** Choose the od pairs with the smallest weight.
 - SMALLEST_WEIGHT** Choose the od pairs with the smallest weight.
 - LARGEST_DISTANCE** Choose the od pairs with the largest euclidian distance.
 - RANDOM** Random.

3.10.4 Integrated line planning, timetabling and vehicle scheduling

Solve the integrated line planning, periodic timetabling and aperiodic vehicle scheduling problem. Includes passenger routing for the timetabling stage. For more information, see [30].

- lin_tim_pass_veh_use_preprocessing** Whether to use an exact preprocessing method to reduce the problem size before optimization.
- lin_tim_pass_veh_add_fixed_passenger_paths** Whether to add the non-routed passengers as fixed weights to the model.
- lin_tim_pass_veh_number_of_routed_od_pairs** The number of routed od pairs.
- lin_tim_pass_veh_time_limit** The time limit for the optimization.
- lin_tim_pass_veh_mip_gap** The mip gap for the optimization.
- lin_tim_pass_veh_write_lp_output** Whether to write the lp output. Will additionally compute an IIS for infeasible programs.
- lin_tim_pass_choose_routed_od_pairs** How to choose the routed od pairs. The following methods are possible:
 - LARGEST_WEIGHT** Choose the od pairs with the smallest weight.
 - SMALLEST_WEIGHT** Choose the od pairs with the smallest weight.
 - LARGEST_DISTANCE** Choose the od pairs with the largest euclidian distance.
 - RANDOM** Random.

3.10.5 Robust Timetabling and Vehicle Scheduling Using Machine Learning

This algorithm tries to improve the robustness of the given timetable and vehicle schedule by using a machine-learned oracle and meta-heuristics for robustness prediction and determining possible improvement steps. For more information, see [22].

For this model to work, a machine-learned oracle needs to be trained first. This step is not part of `LINTIM`. For more information on the training process, see [21]. To compute the key features described there and in the publication above, use

```
R make int-rob-ml-key-features
```

This will create CK `filename_robustness_tensor_file_name` (FI `statistic/data.tensor`) which can then be used for training externally.

The following configuration parameters determine the behavior of the algorithm.

CK **ean_change_penalty** the change penalty to respect when routing passengers

CK **ean_default_maximal_change_time** the maximal change time. Will be used when CK `rob_create_missing_changes` is set to CV `true`

CK **ean_default_minimal_change_time** the minimal change time. Will be used when CK `rob_create_missing_changes` is set to CV `true`

CK **filename_robustness_ml_model** the filename of the machine-learned model to consider. Will only be used when CK `rob_use_api_for_prediction` is set to CV `false`.

CK **gen_passengers_per_vehicle** the vehicle capacity

CK **rob_max_changes** the maximal changes allowed in the key feature vector used for robustness prediction

CK **rob_max_group_size** the maximal passenger group size to route. Grouping passengers may improve routing runtime.

CK **rob_max_iteration** the maximal number of iterations the algorithm is allowed to perform before aborting

CK **rob_max_travel_time** the maximal travel time in the key feature vector used for robustness prediction

CK **rob_max_turnaround_time** the maximale turnaround time allowed in the key feature vector used for robustness prediction

CK **rob_output_every_solution** whether every solution should be written to disk. If set to CV `true`, a subfolder CK `rob_debug_output_path` will be used to store the result of every iteration. Note that this may take up a large amount of disk space when used on large datasets with many iterations.

CK **rob_reroute_interval** the interval to reroute, i.e., setting this to CV `5` will result in rerouting taking place every fifth iteration. Increasing this value may improve the runtime but decrease the prediction quality.

CK **rob_routing_end_time** the time when the routing of the passengers should stop. You should allow enough time for your transportation system to settle after ending the routing of passengers. Events outside of the routing window will not be considered for the key features. Note that we will consider at most 4 hours, i.e., setting this higher will have no effect.

CK **rob_routing_start_time** the time when the routing of the passengers should start. You should allow enough “startup” time for your transportation system to settle before starting the routing of passengers.

CK **rob_start_solutions_file** the start solution file to read. Start solutions are read for the genetic algorithm (i.e. **CK** **rob_use_genetic_algorithm** **CV** true) or when a specific start solution should be used for the local search (i.e. **CK** **rob_local_search_start_solution** ≠ -1). The file should be a zip file containing the possible start solutions each in a separate folder, named e.g. **Fi** **A_10** for start solution with index 10. In this folder should be a valid **LNTIM** dataset.

CK **rob_use_api_for_prediction** will not read the model directly but use an api provided on port **CK** **rob_api_port**. The algorithm will send the key feature vector separated with ";" and expect the resulting values as a ";" separated vector as well, followed by "n". The average of the received vector will be used as the prediction value for the given key feature vector.

CK **rob_use_single_ann_models** will not read a single neural network model but one for each of the four robustness objectives. Will insert "_1", ..., "_4" into the filename, i.e., for **CV** **model.h5** in **CK** **filename_robustness_ml_model**, this will try to read **Fi** **model_1.h5**, ... **Fi** **model_4.h5**.

Specific for the local search, i.e., with **CK** **rob_use_genetic_algorithm** set to **CV** false

CK **rob_ls_allowed_travel_time_increase** the allowed travel time increase of the passengers, i.e., when this is set to **CV** 1.1 the algorithm allows an average travel time increase of 10% before aborting

CK **rob_ls_buffer_increase_per_step** the amount of buffer to add in each step, in seconds.

CK **_ls_candidates_per_type** determines how many candidates per activity type should be added in each neighborhood

CK **rob_ls_change_weight** the weight factor for change activities in the neighborhood selection process

CK **rob_ls_drive_weight** the weight factor for drive activities in the neighborhood selection process

CK **rob_ls_propagate_slack_use_percentage** determines the propagation of slack on activities. When set to **CV** true, **CK** **rob_ls_propagate_slack_percentage** gives the ratio of the activity slack to reduce in each step. When set to **CV** false, a minimal slack time of **CK** **rob_ls_propagate_slack_min_time** will be used instead.

CK **rob_ls_select_by_ratio** when set to true, not the absolute robustness improvement but the robustness improvement divided by the lost passenger travel time will be used to determine the best solution in each neighborhood.

CK **rob_ls_turn_weight** the weight factor for turnover activities in the neighborhood selection process

CK **rob_ls_use_periodic_timetabling** whether to maintain a periodic timetable in every step or not

CK **rob_ls_wait_weight** the weight factor for wait activities in the neighborhood selection process

Specific for the genetic algorithm, i.e., **CK** **rob_use_genetic_algorithm** set to **CV** true

CK **rob_ga_breedings_per_iteration** the number of breedings to perform per generation

CK **rob_ga_mutation_amount** the maximal amount of mutation to use in each mutation

CK **rob_ga_number_mutations_at_breeding** the number of vector entries to mutate during the breeding process

CK **rob_ga_number_mutations_at_start** the number of vector entries to mutate in the start solutions

CK **rob_ga_number_start_solutions** the number of start solutions to use.

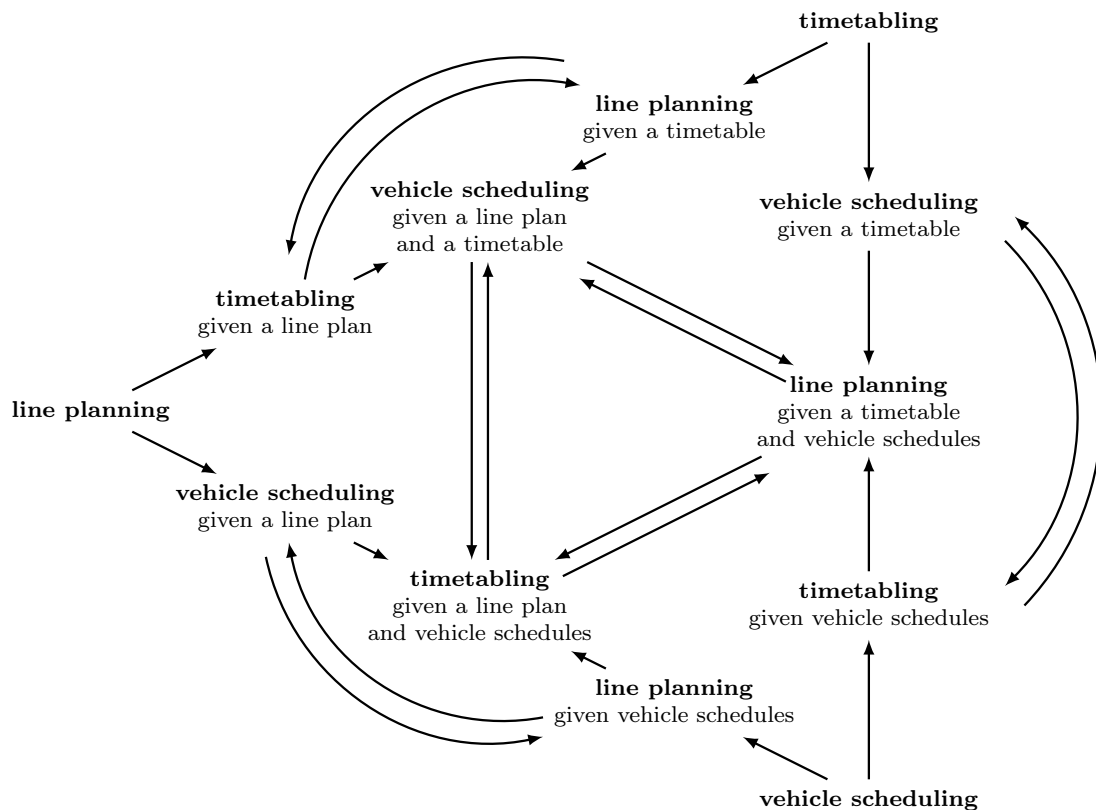


Figure 3.2: Depiction of the eigenmodel described in [37].

- rob_ga_only_best_breeding** whether to use only the best/fittest or all of the population for breeding
- rob_ga_seed** the random seed
- rob_ga_selection** determines how to choose the next generation. While **QUALITY** will only keep the best/fittest solutions, **PARETO** will keep all non-dominated (w.r.t. predicted robustness and travel time) individuals and add the best/fittest solutions if those are not enough (compared to **rob_genetic_solution_pool_size**).
- rob_ga_solution_pool_size** the number of solutions in each generation
- rob_mip_gap** the mip gap for the vehicle scheduling subproblem. Set to -1 to disable.
- rob_threads** the thread limit for the vehicle scheduling subproblem. Set to -1 to disable.
- rob_timelimit** the time limit for the vehicle scheduling subproblem. Set to -1 to not set a time limit.

3.10.6 Eigenmodel

The eigenmodel is a theoretical model for iteratively solving the integrated public transport model. A representation can be seen in Figure 3.2. For more information, see [37].

Tim-Veh-To-Lin

Implementation of one of the steps of the inner circle of the eigenmodel. For a fixed line plan and vehicle schedule, compute a new periodic timetable. For more information, see [29]. Note that this model will only work for line frequencies of 1.

- tim_veh_to_lin_time_limit** The time limit for the optimization.
- tim_veh_to_lin_mip_gap** The mip gap for the optimization.
- tim_veh_to_lin_write_lp_output** Whether to write the lp output. Will additionally compute an IIS for infeasible programs.
- DM_earliest_time_EM** The earliest time for events to consider for this model. Should be large enough that the time between **DM_earliest_time_EM** and **DM_latest_time_EM** is free of any aperiodic side effects.
- DM_latest_time_EM** The latest time for events to consider for this model. Should be small enough that the time between **DM_earliest_time_EM** and **DM_latest_time_EM** is free of any aperiodic side effects.

Chapter 4

Auxiliary Algorithms

4.1 Dataset Generation

With the dataset generator it is possible to create new artificial datasets. To use it, navigate into the `[Fo]` /datasets directory and run

```
[R] make dg-generate-dataset
```

This will create a new subdirectory in `[Fo]` /datasets.

4.1.1 Input

As input, only some parameters in the file `[Fi]` /dataset-generation/basis/Config.cnf are needed.

`[CK]` `ptn_name` The name for the new dataset. As default, this is set to be `new_generic_dataset`.

`[CK]` `dg_model` specifies the method by which the new dataset is created.

Depending on the chosen `[CK]` `dg_model`, some more config parameters are required; see below.

4.1.2 Output

As output a new directory `[Fo]` /datasets/`[CK]` `ptn_name` is created. The config file from the directory `[Fo]` dataset-generation is copied into the new dataset. This is then ready to be used as a dataset with all functionalities of `LINTM`.

4.1.3 Algorithms

Parametrized City

`[CK]` `dg_model` `[CV]` `parametrized_city`

The model is based on a paper by Fielbaum et al. [5]. This model divides a city into various zones. The authors state, that most big cities consist of one Central Business District (CBD) surrounded by some subcenters. As output, the files `[Fi]` Stop.giv, `[Fi]` Edge.giv and `[Fi]` OD.giv are created in the new directory. The PTN is generated by the following procedure:

First, the CBD is represented by a node in the center of the PTN. The CBD is surrounded by n zones, each of which consists of a subcenter-node and a periphery node. All the subcenters are then connected to the CBD and their neighboring subcenters. The periphery nodes are only connected to their own subcenter.

The distance between a subcenter and the geometrical center C of the graph is L . It is not necessary that the CBD is located at C , but it can have an offset to C by ηL along an axis CBD-subcenter. The distance

between a periphery and its subcenter is gL .

Considering the creation of the OD-Matrix, the parameter Y states how many trips are generated in total. They are evenly splitted among the n zones, such that exactly $\frac{Y}{n}$ trips start in each zone. A fraction a of those trips start in the subcenter and a fraction of $b = 1 - a$ depart from the periphery. Usually we have $b < a$. A fraction of α of all trips generated in a periphery goes to the CBD and a fraction of β goes to it's own subcenter. The rest ($\gamma = 1 - \alpha - \beta$) goes evenly splitted to all other (foreign) subcenters.

To use the Parametrized city model, specify the following parameters:

- gen_vehicle_speed Speed of the vehicles in km/h.
- dg_param_city_number_subcenters Number of subcenters surrounding the CBD. The PTN has $2n + 1$ nodes.
- dg_param_city_alpha Trips proportion from periphery that go to the CBD.
- dg_param_city_beta Trips proportion from periphery to own subcenter. From α and β we calculate the value of $\gamma = 1 - \alpha - \beta$ representing the trips proportion from periphery to foreign subcenters.
- dg_param_city_eta Portion of displacement of the CBD from the center of the city in an axis CBD-subcenter.
- dg_param_city_Y Total number of trips generated.
- dg_param_city_L Distance from any subcenter to the geometrical center of the city.
- dg_param_city_g Distance periphery-subcenter / Distance subcenter-CBD.
- dg_param_city_a Trips proportion that depart from the periphery. From this we calculate the value of $b = 1 - a$ representing the trips proportion that depart from a subcenter.

According to [5] the parameters in table 4.1.3 should give a reasonable model of the corresponding cities.

	Santiago	Bordeaux	Los Angeles
n	7	3	7
α	0.25	0.18	0.0033
β	0.22	0.72	0.287
γ	0.53	0.1	0.68
η	0	0	0
Y	2,565,622	250,000	4,500,000
L	10	6.6	11.65
g	0.85	1.2	0.79
a	0.78	0.3	0.91
b	0.22	0.7	0.09

Table 4.1: These parameters should reproduce a reasonable model for the corresponding cities.

Ring

- dg_model ring

This model creates an undirected PTN consisting of some concentric rings and a center node. For each edge the lower bound is set to 1 and the upper bound is set to 20. The following parameters control the layout:

- dg_ring_number_of_rings Number of concentric rings that are generated.
- dg_ring_nodes_per_ring Number of nodes that each ring consists of

`dg_ring_length_1` If this boolean parameter is set to `true`, the lengths of all edges are equal to 1.

`dg_ring_radius` specifies the radius of the inner ring, i.e. the lengths of the edges from the center to the nodes of the inner ring. The lengths of all other edges are set according to the euclidean distance in the plane. Only used, if `dg_ring_length_1` is `false`.

There are different methods for the creation of the OD data, specified by `dg_ring_demand_type`. OD-values are always created symmetrically and they are equal to zero if both nodes are identical. Available options are:

`UNIFORM` All OD-values are set to 1.

`UNIFORM_CENTRE` If one of the nodes is the center node, the OD value is 100, otherwise it is 10.

`RANDOM` All OD-values are set to random integers between 1 and 100.

`RANDOM_NEIGHBOUR_CENTRE` If one of the nodes is the center node, then the OD-value is a random integer between 40 and 150. If there exists an edge between both nodes, the OD-value is a random integer between 20 and 50 and otherwise it is a random integer between 0 and 30.

`SPOKE_RING` For each pair of nodes we compute a shortest path in the PTN with respect to the euclidean distance (even if the edge lengths are set to 1). Along this path we count the number of ring edges and the number of spoke edges. A spoke edge is an edge between two nodes in different rings. The following parameters need to be specified:

`dg_ring_spoke_edge_demand`

`dg_ring_ring_edge_demand`

`dg_ring_demand_scaling_factor`

The OD value is then computed as

$$\text{scaling_factor} \cdot \left(\frac{\text{spoke_edge_demand}}{\#\text{spoke edges} + 1} + \frac{\text{ring_edge_demand}}{\#\text{ring edges} + 1} \right)$$

rounded to the nearest integer.

4.2 OD Matrix Creation

In the OD matrix creation step, an OD matrix is calculated using a given demand and a PTN.

4.2.1 Input

The following files are needed as input:

- `default_stops_file` (`basis/Stop.giv`) stops of the PTN
- `default_edges_file` (`basis/Edge.giv`) edges of the PTN
- `default_demand_file` (`basis/Demand.giv`) demand at geographical positions

4.2.2 Output

The following file is produced as output:

- `default_od_file` (`basis/OD.giv`) OD matrix for one planning period

4.2.3 Algorithms

To compute an OD matrix run

```
R make od-create
```

For all pairs of demand point a shortest path is computed, which includes the path to and from the PTN and might also not use any PTN edges. The demand at one demand point is distributed randomly to all other demand points with probabilities proportional to

$$\frac{\text{demand at other demand point}}{(\text{distance between demand points})^2}$$

The passengers which are computed to travel between to demand points are attributed to the OD pairs consisting of the first and last station on the shortest path. If the shortest path does not contain any stations, the passengers are not counted towards the OD matrix.

The following parameters can be used to influence the OD matrix which is created:

- `CK od_use_network_distance`: if set to `true`, the distance between demand point which is used for distributing passengers to destination demand points is the travel time between the demand points on the shortest paths. Otherwise it is proportional to the geographical distance between the demand point depending on the norm `CK sl_distance`.
- `CK od_remove_uncovered_demand_points`: if set to `true`, demand points which are more than `CK sl_radius` away from the nearest station are not included in the computation.
- `CK od_network_acceleration`: speed up factor for driving in the PTN compared to traveling directly, also used for driving to and from the network.
- `CK ptn_stop_waiting_time`: the time (in minutes) a vehicle has to stop at each station which is considered during the computation of the shortest path.

4.2.4 Distribute from node demand

If an od demand based on an infrastructure is given, i.e., `CK filename_od_nodes_file` (`Fi basis/OD-Node.giv`), an od distribution algorithm can be used to create a stop based od matrix. For this, run

```
R make od-distribute-from-nodes
```

to obtain `CK default_od_file` (`Fi basis/OD.giv`). This will find travel-time-minimal paths for all passengers and create a stop od matrix based on their chosen route, i.e., the first boarding station and the last alighting station will determine the new od matrix. For this, the walking edges provided in `CK filename_walking_edge_file` (`Fi basis/Edge-Walking.giv`) and a penalty factor for walking, i.e., `CK gen_walking_utility`, will be considered. The drive time on infrastructure edges is based on `CK ean_model_weight_drive` and the waiting time at stations is calculated based on `CK ean_model_weight_wait`. Additionally, the obtained assignment from node od pair to stop od pair can be written to `CK filename_od_node_assignment_file` (`Fi basis/OD-Node-Assignment.giv`) by setting `CK od_node_write_assignment` to `CV true`.

4.3 Load distribution

This step takes the OD matrix and distributes the passengers to the PTN. The resulting edge loads are used as an input for following steps, e.g. most line planning algorithms. This section first handles the setting of `CK load_generator_model` to `CV LOAD_FROM_PTIN`. For the case of setting `CK load_generator_model` to `CV LOAD_FROM_EAN`, see 4.3.4, and for the case of setting it to `CV spanners`, see 4.3.5.

4.3.1 Input

The following files are needed as input:

- default_stops_file (basis/Stop.giv)
- default_edges_file (basis/Edge.giv)
- default_od_file (basis/OD.giv)

When parameter load_generator_use_cg is set to true, the line pool is needed as well to build the Change&Go-network, i.e.,

- default_pool_file (basis/Pool.giv)
- default_pool_cost_file (basis/Pool-Cost.giv)

4.3.2 Output

The following file is produced as output:

- default_loads_file (basis/Load.giv)

4.3.3 Algorithms

To compute a new load, run

```
 make ptn-regenerate-load
```

There are different objective functions to distribute the passengers, namely

- load_generator_type SP
- load_generator_type REWARD
- load_generator_type REDUCTION

SP distributes the passengers on shortest paths. For determining the length of a PTN edge, parameter ean_model_weight_drive is used.

The load generators REWARD and REDUCTION are iterative and include an additional term, rewarding in different ways the bundling of passengers. The weight of the additional terms is determined by load_generator_scaling_factor. REDUCTION adds a penalty depending on the usage of the edge in PTN (high penalty for low usage) and REWARD rewards an edge more if less passengers are needed to fill the next vehicle on the edge. For a more detailed description of the models, see [8].

There are two other parameters to determine the behavior of the algorithm:

load_generator_use_cg When this is set to true, a Change&Go-network is used for routing the passengers. This includes the knowledge of the line pool, allowing to consider transfers. The cost of a transfer will be the estimated change time (load_generator_min_change_time_factor times ean_default_minimal_change_time; at most ean_default_maximal_change_time) plus ean_change_penalty. For waiting at a stop, the behavior of ean_model_weight_wait is adopted. For a more detailed description of the Change&Go-network see [39]. Since the network to route in is much larger, this increases the runtime, especially for bigger pools. But the resulting load is often more realistic.

load_generator_number_of_shortest_paths This determines the number of shortest paths the passenger are distributed to, i.e., if this is set to K , the K shortest paths are computed in each step. This increases the runtime! To distribute the passengers on the different paths, a logit model with parameter load_generator_sp_distribution_factor is used.

For an undirected PTN the algorithm does not distinguish the direction in which an edge is traversed, i.e., the load on an edge is the sum of the numbers of passengers traversing it in each direction. To determine the lower and upper frequency values in the default_loads_file (basis/Load.giv), the resulting load is divided by the vehicle capacity gen_passengers_per_vehicle. Overall, the following parameters determine the behavior of the algorithm:

- ean_change_penalty
- ean_default_maximal_change_time
- ean_default_maximal_waiting_time
- ean_default_minimal_change_time
- ean_default_minimal_waiting_time
- ean_model_weight_drive
- ean_model_weight_wait
- gen_passengers_per_vehicle
- load_generator_add_additional_load
- load_generator_fixed_upper_frequency
- load_generator_fix_upper_frequency
- load_generator_lower_frequency_factor
- load_generator_max_iteration
- load_generator_min_change_time_factor
- load_generator_model
- load_generator_number_of_shortest_paths
- load_generator_scaling_factor
- load_generator_sp_distribution_factor
- load_generator_type
- load_generator_use_cg
- load_generator_upper_frequency_factor

4.3.4 Using the EAN

If load_generator_model is set to LOAD_FROM_EAN, the EAN is used to determine the load of the PTN edges. Therefore the EAN is read and has to be present, i.e., the files

- default_events_periodic_file (timetabling/Events-periodic.giv)
- default_activities_periodic_file (timetabling/Activities-periodic.giv)

4.3.5 Using spanner MIPs

If `load_generator_model` is set to `spanners`, a mixed-integer formulation of the problem is used to distribute the load and find an optimal subgraph of the PTN. The three measures of quality considered in this model are:

Building cost The cost of the subgraph (i.e., total length of the edges)

Total travel time The sum of travel times for each OD pair, weighted by the demand

Maximum detour factor The maximum (over all OD pairs) ratio between shortest path in subgraph and shortest path in full graph.

The model is described in detail in Heinrich et al. (2023) [16]. The goal of this model is to find a subgraph of the PTN satisfying the given constraints and minimizing an objective. This method gives flexibility to the user in choosing which measures should be considered as constraints and which one as the objective. The model assumes that all passengers choose the shortest path available to them.

The main parameters determining the behavior of the algorithm are:

`load_generator_building_cost`

`load_generator_travel_time`

`load_generator_detour_factor`

`load_generator_max_building_cost`

`load_generator_max_travel_time`

`load_generator_max_detour`

The first three parameters are allowed to take values `obj`, `cons` and `none`. Only one of them should be set to `obj`, thus setting the objective of the MIP to minimize building costs of the spanner, total travel time of passengers, or the maximum detour factor. The rest of the parameters can be set to `cons`, adding a constraint determining an upper bound for the respective expression (determined by the latter three parameters above), or `none`, if, e.g., total travel time should not be considered as either the objective function or a constraint.

Other parameters affecting the MIP are:

`load_generator_gen_cuts` When this is set to `true`, so-called valid inequalities are added to the MIP. These have no effect on the solution, but tend to greatly improve computational performance.

`load_generator_remove_unused_edges` When this is set to `true`, after solving the model, the unused edges are removed from `default_edges_file` (`basis/Edge.giv`).

Finally, the solver parameters can be set using:

`load_generator_solver` determine the solver to be used, the supported solvers are Gurobi, CPLEX and Xpress.

`load_generator_threads` determine the maximal number of threads to use for the solver (-1=use default value, i.e., no restriction).

`load_generator_timelimit` the time limit for the solver in seconds (-1=use default value).

`load_generator_mip_gap` the mip optimization gap for the solver, 0.1 equals a gap of 10 % (-1=use default value).

`load_generator_write_lp_file` whether to write the lp file of the model to solve.

4.4 Headway creation

This is a small helper script to create a headway file for the current dataset. Some older methods still need a headway file present, even if the content is not used.

4.4.1 Input

The following file is needed as input

- `default_edges_file` (`basis/Edge.giv`) edges of the PTN

4.4.2 Output

The following file is produced as output:

- `default_headways_file` (`basis/Headway.giv`) a file containing a default headway value for each edge

4.4.3 Algorithm

To create the headways, run

```
 make ptn-headways
```

This will create a new headway file, using `ptn_default_headway_value` as a value for each edge.

4.5 PTN to EAN

4.5.1 Input

The following files are required as input

- `default_stops_file` (`basis/Stop.giv`) edges of the PTN
- `default_edges_file` (`basis/Edge.giv`) edges of the PTN
- `default_lines_file` (`line-planning/Line-Concept.lin`) a line concept on the PTN

4.5.2 Output

This procedure gives the following output

- `default_events_periodic_file` (`timetabling/Events-periodic.giv`)
- `default_activities_periodic_file` (`timetabling/Activities-periodic.giv`)

4.5.3 Algorithm

To create the Event-Activity-Network (required as input for Timetabling etc.), run

```
 make ean
```

The event-activity-network is then created. To this end for every line departure and arrival events for every station the line passes (every line is executed in both directions, depending on `ptn_is_undirected`) will be created. These events are then connected either with drive or wait activities (respecting the bounds given by the configuration of `ean_default_minimal_waiting_time` etc.). Furthermore it will assign each arc with some weight, corresponding to the amount of passengers driving on it. To this end, the

passengers are routed along shortest paths in the EAN, where the calculation of the edge lengths assumes that the times for each activity are given by `ean_model_weight_drive` (resp. wait/change/etc.). Note that this computation ignores the loads of the edges in the PTN.

Per default `ean_construction_target_model_frequency` is set to

`FREQUENCY_AS_MULTPLICITY`, which additionally creates synchronisation activities between every repetition of each line. This ensures that in the EAN the frequency of each line is indeed respected. Note, that such synchronisation activities have fixed upper and lower bounds, that are equal. If the frequency of a line does not divide the period length, this routine will distribute the remaining time buffer evenly to the different activities.

If headways exist, they can also be created for the EAN by setting

`ean_construction_target_model_headway` to something different than

`NO_HEADWAYS` (which is the default), e.g. to `SIMPLE`.

Individual station limits can be provided by `filename_station_limit_file` (`basis/Station-Limits.giv`) when `ean_individual_station_limits` is set to `true`. For every station in the station limit file, the given individual limits will be used. For stops not in the limit file or entries of -1 the global default values will be used.

Additionally, it is possible to restrict the set of stations where transfers may take place. For this, set `ean_respect_change_stations` to `true` and provide a list of possible transfer stations in `filename_change_station_file` (`basis/Change-Stations.giv`). Transferring in other stations will be forbidden, i.e., no transfer activities will be created there.

It is also possible to enable walking, i.e., transferring between different stops connected by walking edges. For this, `ean_use_walking` must be set to `true` and an infrastructure network with corresponding walking edges needs to be provided that is consistent with the PTN used, i.e., we assume that the node id of the corresponding node is stored in the long name of the stops. Additionally, a total maximal walking time (`sl_max_walking_time`) can be provided, only allowing walking transfers with the given maximal length.

The following parameters control the behavior of the algorithm:

`debug_paths_in_ptn`

`debug_paths_in_ean`

`ean_algorithm_shortest_paths`

`ean_change_penalty`

`ean_construction_skip_passenger_distribution`

`ean_construction_target_model_frequency`

`ean_construction_target_model_headway`

`ean_default_maximal_change_time`

`ean_default_maximal_waiting_time`

`ean_default_minimal_change_time`

`ean_default_minimal_waiting_time`

`ean_discard_unused_change_activities`

`ean_dump_initial_duration_assumption`

`ean_individual_station_limits`

`ean_initial_duration_assumption_model`

`ean_model_weight_change`

`ean_model_weight_drive`

`ean_model_weight_wait`

`ean_random_shortest_paths`

`ean_use_walking`

`period_length`

`sl_max_walking_time`

4.6 EAN buffer activities

There are several algorithms to add buffer times to the EAN. All methods are called using

`make ean-buffer-activities`

and the implementation used is determined by the config parameter `rob_buffer_generator` with the following choices:

- `exponential`: Exponential distribution
- `reverse-exponential`: Reverse exponential distribution
- `uniform-random`: Uniform random buffer distribution
- `exceed-random`: Uniform random distribution with an additional upper bound
- `proportional`: Add a fixed buffer to all activities
- `proportional-restricted`: Buffer all activities with a fixed term, but restrict the number of events or activities to buffer

For `proportional-restricted`, the following config parameters determine the behavior:

- `rob_buffer_link_list`: A given list of link ids to buffer. All activities belonging to the given links will be buffered
- `rob_buffer_on_wait_activity`: The buffer to add to wait activities, only activities determined by the `rob_buffer_stop_percentage` will be buffered.
- `rob_buffer_on_drive_activity`: The buffer to add to drive activities, only activities determined by the `rob_buffer_link_percentage` or `rob_buffer_link_list`: will be buffered.
- `rob_proportional_drive_activity_buffer`: An additional percentage based buffer for the drive activities, should be between 0 and 1
- `rob_buffer_link_percentage`: The percentage of links to buffer. Will buffer all drive activities on the most used links, i.e., the links with the most drive activities. Should be between 0 and 1.
- `rob_buffer_stop_percentage`: The percentage of stops to buffer. Will buffer all wait activities at the most used stops, i.e., the stops with the most changing passengers. Should be between 0 and 1.

The buffered activities will be written to `default_activity_buffer_file` and

`use_buffered_activities` will be set to `true`. Reading

`default_activities_periodic_file` should always return the value for

`default_activity_buffer_file` when `use_buffered_activities` is set to `true`.

4.7 EAN reroute passengers

R `make ean-reroute-passengers`

This generates a passenger distribution (i.e., new weights on the activities) by rerouting the passengers (i.e., the OD pairs) through the periodic EAN on shortest paths with respect to the timetable derived durations. Note that the passengers of the same OD pair will not be split up, but will all use the same shortest path in the EAN.

4.8 Tariff (Reference) Price Matrix

Running

R `make taf-tariff-price-matrix`

creates a price matrix for a specified tariff with given tariff information. No optimization is done.

Running

R `make taf-tariff-reference-price-matrix`

creates a reference price matrix for a specified tariff with given tariff information.

4.8.1 Input

The following files are needed as input:

- CK `default_stops_file` (Fi `basis/Stop.giv`), stops of the PTN,
- CK `default_edges_file` (Fi `basis/Edge.giv`), edges of the PTN.

If CK `taf_model` CV `flat`, the following configuration value is needed:

- CK `taf_flat_price`, the constant price for all paths in the tariff.

If CK `taf_model` CV `beeline_distance` or CV `network_distance`, the following configuration values are needed:

- CK `taf_fixed_costs`, the fixed costs of an affine beeline distance or network distance tariff,
- CK `taf_factor_costs`, the factor costs of an affine beeline distance or network distance tariff.

If CK `taf_model` CV `zone`, then the two following files are additionally needed as input:

- CK `filename_tariff_zone_file` (Fi `tariff/Zones.taf`), assignment of stops to zones,
- CK `filename_tariff_zone_price_file` (Fi `tariff/Zone-Prices.taf`), price list of the zone tariff.

If CK `taf_model` is CV `network_distance` or CV `zone` and CK `taf_routing_generation` is CV `read-all` or CV `read-partial-fill`, then the following file is additionally needed as input:

- CK `filename_routing_ptn_input` (Fi `basis/Routing-ptn.giv`).

4.8.2 Output

Running

```
[R] make taf-tariff-price-matrix
```

yields as output

- [CK] `filename_tariff_price_matrix_file` ([Fi] `tariff/Price-Matrix.taf`), prices for each OD pair,

and

```
[R] make taf-tariff-reference-price-matrix
```

yields as output

- [CK] `filename_tariff_reference_price_matrix_file` ([Fi] `basis/Reference-Price-Matrix.giv`), reference prices for each OD pair.

If [CK] `taf_model` is [CV] `network_distance` or [CV] `zone`, the routing file is produced as output as well:

- [CK] `filename_routing_ptn_output` ([Fi] `basis/Routing-ptn.giv`), routing in the PTN.

In all cases, the following statistic file is also produced as output:

- [CK] `filename_tariff_properties_file` ([Fi] `statistic/tariff-properties.sta`), statistic file containing information whether the no-elongation property and the no-stopover property (see Section 3.7.4) are fulfilled for the computed price matrix.

4.8.3 Algorithms

Price Matrix

Run

```
[R] make taf-tariff-price-matrix
```

to create a price matrix for all OD pairs for a specified tariff with given tariff information. The following models [CK] `taf_model` are available:

- [CV] `flat`, write price matrix for a flat tariff,
- [CV] `beeline_distance`, write price matrix for an affine beeline distance tariff,
- [CV] `network_distance`, write price matrix for an affine network distance tariff,
- [CV] `zone`, write price matrix for a zone tariff.

If [CK] `taf_model` [CV] `flat`, [CK] `taf_flat_price` is read and for each non-trivial OD pair (i.e. for all $d \in D := (V \times V) \setminus \{(v, v) : v \in V\}$) this flat price is written to [CK] `filename_tariff_price_matrix_file` ([Fi] `tariff/Price-Matrix.taf`). For trivial OD pairs $d = (v, v)$ for $v \in V$ the price is set to 0.

Be aware that entries for [CK] `taf_flat_price` in the private configuration file overwrite those specified in the state configuration file (see Section 8.1).

If [CK] `taf_model` [CV] `beeline_distance`, the Euclidean distances l_d between the start and end station of each OD pair d are determined. Then the fixed costs f ([CK] `taf_fixed_costs`) and factor costs p ([CK] `taf_factor_costs`) are read from the configuration file and the prices are determined for each non-trivial OD pair d by $f + l_d \cdot p$ and written to [CK] `filename_tariff_price_matrix_file` ([Fi] `tariff/Price-Matrix.taf`). For trivial OD pairs $d = (v, v)$ for $v \in V$ the price is set to 0.

Be aware that entries for `CK taf_fixed_costs` and `CK taf_factor_costs` in the private configuration file overwrite those specified in the state configuration file (see Section 8.1).

If `CK taf_model CV network_distance`, for each OD pair the summed edge lengths of the routing specified by `CK taf_routing_generation` with the following possible values are used as distances:

- `CV fastest-paths`, a new routing using fastest paths with respect to the lower time bounds is created,
- `CV read-all`, a routing given in `CK filename_routing_ptn_input (Fi basis/Routing-ptn.giv)` is read and used to determine distances,
- `CV read-partial-fill`, a partial routing given in `CK filename_routing_ptn_input (Fi basis/Routing-ptn.giv)` is read and used to determine distances. Unspecified paths are filled with fastest paths with respect to the lower time bounds.

From this the paths lengths l_d are calculated by summing up the edge lengths on the path. For all non-trivial OD pairs $d \in D$ the prices are calculated by $f + l_d \cdot p$ with fixed costs f (`CK taf_fixed_costs`) and factor costs p (`CK taf_factor_costs`). For trivial OD pairs $d = (v, v)$ for $v \in V$ the price is set to 0. The prices are written to `CK filename_tariff_price_matrix_file (Fi tariff/Price-Matrix.taf)`.

Be aware that entries for `CK taf_fixed_costs` and `CK taf_factor_costs` in the private configuration file overwrite those specified in the state configuration file (see Section 8.1).

If `CK taf_model CV zone`, then `CK filename_tariff_zone_file (Fi tariff/Zones.taf)`, `CK filename_tariff_zone_price_file (Fi tariff/Zone-Prices.taf)` and `CK taf_zone_counting` are read. From this for each non-trivial OD pair a path specified by `CK taf_routing_generation` (as explained above and in Section 3.7.1) is used to determine the price by counting the number of traversed zones respecting `CK taf_zone_counting` (as explained in Section 3.7.4). For trivial OD pairs $d = (v, v)$ for $v \in V$ the price is set to 0. The prices are written to `CK filename_tariff_price_matrix_file (Fi tariff/Price-Matrix.taf)`.

Reference Price Matrix

Run

```
R make taf-tariff-reference-price-matrix
```

to create a reference price matrix for all OD pairs for a specified tariff with given tariff information. This command follows the same logic as

```
R make taf-tariff-price-matrix
```

with the difference being that the prices are written to `CK filename_tariff_reference_price_matrix_file (Fi basis/Reference-Price-Matrix.giv)` instead of `CK filename_tariff_price_matrix_file (Fi tariff/Price-Matrix.taf)`.

4.9 Rollout

The periodic event-activity network and the periodic timetable have to be converted to a nonperiodic event-activity network that can be used in the operational phase of public transport.

4.9.1 Input

The following files are needed as input

- default_edges_file (default: basis/Edge.giv)
- default_headways_file (default: basis/Headway.giv)
- default_events_periodic_file (timetabling/Events-periodic.giv)
- default_activities_periodic_file (timetabling/Activities-periodic.giv)

4.9.2 Output

The following files are produced as output:

- default_events_expanded_file (delay-management/Events-expanded.giv) a file containing the aperiodic events
- default_activities_expanded_file (delay-management/Activities-expanded.giv) a file containing the aperiodic activities
- default_timetable_expanded_file (delay-management/Timetable-expanded.tim) a file containing the aperiodic timetable

4.9.3 Algorithm

To roll out, all (nonperiodic) events that take place in the time interval [DM_earliest_time, DM_latest_time] (given in seconds since 0:00) as well as all (nonperiodic) activities connecting those events are taken into account. If rollout_whole_trips is set to true, all trips whose start event or end event are not contained in [DM_earliest_time, DM_latest_time] are deleted. If rollout_discard_unused_change_edges is set to true, changing activities with weight 0 are ignored (this might significantly reduce the size of the nonperiodic event-activity network, speeding up the delay management step). The parameter rollout_for_nonperiodic_timetabling influences the output: if set to true, only forward headways are contained in the output, and for each activity, the output also contains an upper bound on its duration (note that this parameter always should be set to false unless you really know what you are doing!).

Delay Management and Vehicle Scheduling When rolling out for vehicle scheduling, usually a long time period (e.g. a whole day) is considered and rollout_whole_trips *must* be set to true. When rolling out for delay management, usually a short time period (e.g. two hours) is considered and rollout_whole_trips should be set to false. Typically, the combination of vehicle scheduling and delay management could be like this:

1. Set [DM_earliest_time, DM_latest_time] to a “large” time interval, e.g. one day, and rollout_whole_trips to true.
2. Run
 make ro-rollout && make ro-trips
3. Run
 make vs-vehicle-schedules
to generate the vehicle schedules.

4. Set `[CK] DM_earliest_time`, `[CK] DM_latest_time` to the time interval needed for delay management, e.g. two hours, and `[CK] rollout_whole_trips` to `[CV] false`.

5. Run

```
[R] make ro-rollout && make vs-add-circulations-to-ean
```

to roll out for delay management and to add the circulations to the rolled-out event-activity network.

Generating passenger paths For more precise methods of delay management, OD pairs may be rolled out over the delay management period into distinct paths in the aperiodic EAN. As this takes quite some time in the rollout and in the evaluation of the delay management, this has to be explicitly enabled by setting the `rollout_passenger_paths` parameter to true. A new file determined by `default_passenger_paths_file` will be created containing in each line a departure event, an arrival event, the source and target station id, an integral passenger weight and a comma-separated list of change activities. The weights are distributed from the original OD file, where passengers are equally distributed over the time between `DM_earliest_time` and the departure time of their last connection. Every passenger gets assigned to the next possible departure event. If there exists multiple paths with the same arrival time, among them only those with a minimal number of changes and with the latest possible departure time will be kept and considered. A path for which another path with the same or a later departure time but an earlier arrival time exists will not be considered either. If there still are multiple paths for one departure time, the passengers will be divided between them equally but integrally (such that some of them may have 1 passenger less than others). If passenger paths are rolled out, there will be an additional file according to `default_od_expanded_file` will be created. This file contains a timestamped OD demand according to the path-distribution of the passengers.

4.9.4 Requirements and caveats

- If `[CK] DM_enable_consistency_checks` is set to `[CV] true`, IDs in files are checked to be consecutively numbered beginning from 1.

4.9.5 Generating trips

For vehicle scheduling, it is necessary to additionally create the trips after rolling out, i.e., after

```
[R] make ro-rollout
```

with `[CK] rollout_whole_trips` set to `[CV] true`,

```
[R] make ro-trips
```

should be run as well. This method uses the files

- `[CK] default_activities_expanded_file` (`[Fi] delay-management/Activities-expanded.giv`)
- `[CK] default_events_expanded_file` (`[Fi] delay-management/Events-expanded.giv`)

to create

- `[CK] default_trips_file` (`[Fi] delay-management/Trips.giv`)

The file `[CK] default_trips_file` (`[Fi] delay-management/Trips.giv`) will then contain all information regarding line trips that need to be covered of a feasible vehicle schedule.

4.10 Delay generation

To simulate source delays during the operational phase, different delay generators are included in `LINTM`. The following parameters are used by all delay generators:

- The interval `[CK delays_min_time, CK delays_max_time]` defines which events and/or activities might be delayed (only events taking place in this time interval or activities connecting two such events might be delayed). Note that `[CK delays_min_time, CK delays_max_time] ⊆ [CK DM_earliest_time, CK DM_latest_time]` is required.
- The parameters `CK delays_min_delay` and `CK delays_max_delay` define lower and upper bounds on the amount of a source delay. If `CK delays_absolute_numbers` is set to `CV true`, the bounds are in seconds, otherwise the bounds are in % of the nominal duration of a delayed activity (this is needed for delays on activities only).
- If `CK delays_append` is set to `CV true`, the generated source delays are appended to already existing files containing source delays (to allow a combination of delays, generated by different delay generators); if set to `CV false`, existing files containing source delays are replaced. Please note that several source delays of the same event (activity) are not additive: newly generated source delays are simply appended to the file containing the source delays, and this file is read sequentially – so for each event (activity), only the last source delay contained in the file is taken into account.

Which generator is going to be used is controlled by the `CK delays_generator` parameter.

CV uniform_distribution: Adds random source delays to randomly chosen events and/or activities. Its behavior can be controlled by the following parameters:

- `CK delays_events`: If set to `CV true`, source delays on events are generated (can be combined with `CK delays_activities`).
- `CK delays_activities`: If set to `CV true`, source delays on driving activities are generated (can be combined with `CK delays_events`).
- `CK delays_count`: Number of source delays that will be generated. If `CK delays_count_is_absolute` is set to `CV true`, `CK delays_count` is an absolute number; otherwise it defines how many events of all events taking place in `[CK delays_min_time, CK delays_max_time]` (in %) and/or how many driving activities of all driving activities with start event and end event in `[CK delays_min_time, CK delays_max_time]` (in %) will be delayed.

CV events_in_station: Delays *all* events in the station defined by

`CK delays_station_id_for_delays`. If `CK delays_station_id_for_delays` is `CV -1`, the station is chosen randomly. If you want to delay all events in several different stations, you have to run the delay generator several times with different values of `CK delays_station_id_for_delays` and `CK delays_append` set to `CV true`.

CV activities_on_track: Delays *all* driving activities on the track defined by `CK delays_edge_id_for_delays`. If `CK delays_edge_id_for_delays` is `CV -1`, the track is chosen randomly. If you want to delay all driving activities on several different tracks, you have to run the delay generator several times with different values of `CK delays_edge_id_for_delays` and `CK delays_append` set to `CV true`.

CV uniform_background_noise: Adds random source delays to every event and/or activity. Its behavior can be controlled by the following parameters:

- `CK delays_seed`: For reproducible purpose a seed for generating random delay amount is introduced. If `delays seed` is set to `CV 0`, no seed will be set and thus the experiment in general is not reproducible.

- `CK` `delays_events`: If set to `CV` `true`, source delays on events are generated (can be combined with `CK` `delays_activities`).
- `CK` `delays_activities`: If set to `CV` `true`, source delays on driving activities are generated (can be combined with `CK` `delays_events`).
- `CK` `delays_append`: If this is set to `CV` `true`, the already delayed events and activities are not further manipulated.

4.11 Visualization

LINTIM offers algorithms for drawing several states of the public transportation system. The output files can be found in `Fo` `graphics`.

4.11.1 PTN

To create an illustration of the PTN run

```
R make ptn-draw
```

The result for dataset toy is depicted in Figure 4.1.

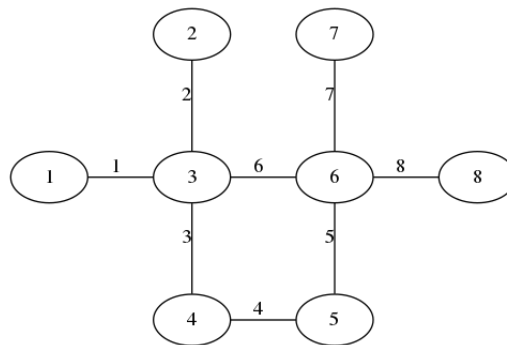


Figure 4.1: The PTN of the toy dataset

The graph can be scaled by adapting `CK` `ptn_draw_conversion_factor`.

Setting `CK` `ptn_draw_use_coordinates` to `CV` `false` results in disregarding the stop-coordinates. Instead, the stops are arranged automatically. The result for dataset toy is depicted in Figure 4.2.

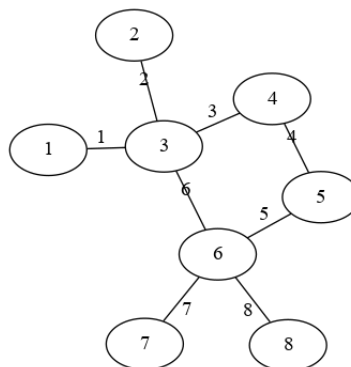


Figure 4.2: The PTN of the toy dataset with automatically arranged stops

To create an illustration of the PTN that is readable even for larger datasets, run

```
R make ptn-draw-interactive
```

The resulting html-script allows for some interaction, like changing node sizes or viewing network information when tracing over the graph. One possible output for dataset `bahn-01` is depicted in Figure 4.3. Edge labels can be enabled with `CK ptn_draw_interactive_graph_edge_labels`.

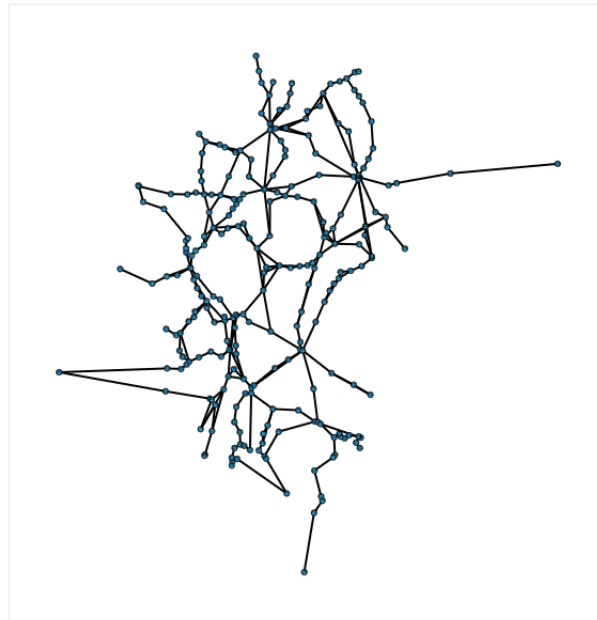


Figure 4.3: The PTN of the bahn-01 dataset

4.11.2 OD

To create an illustration of the OD data run

```
R make od-draw
```

The result for dataset `toy` is depicted in Figure 4.4. The graph displays only those OD pairs whose

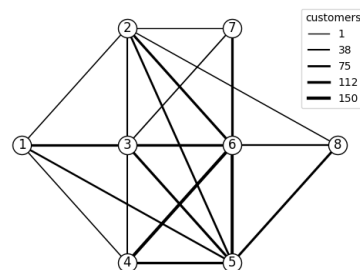


Figure 4.4: The OD data of the toy dataset where the edge width indicates the number of passengers traveling

fractional value in relation to the maximal value of the OD pairs lies within the closed interval given by `CK od_visualization_lower_bound` and `CK od_visualization_upper_bound`. Datasets with

symmetric OD data will be illustrated using undirected graphs. Otherwise a directed graph will be used. The output is saved in `filename_od_visualization_file`. The graph can visualize the logarithm of the number of passengers traveling with `od_visualization_use_log_scale`. The graphs maximal edge width can be adjusted with `od_visualization_max_edge_width`. The number of passengers traveling can be indicated with edge color instead of edge width using the parameter `od_visualization_use_edge_color`. The result for dataset toy is depicted in Figure 4.5. Either graph can be scaled by adapting `od_draw_conversion_factor`.

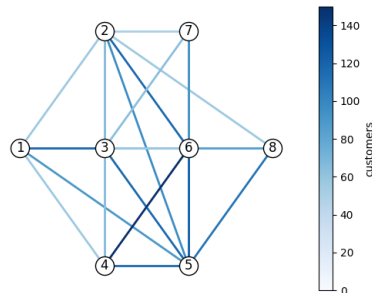


Figure 4.5: The OD data of the toy dataset where the edge color indicates the number of passengers traveling

Alternatively, a heatmap visualization can be used with `od_visualization_use_heatmap`. It can be annotated using `od_visualization_use_annotations`. As with the graph visualization, the heatmap can visualize the logarithm of the number of passengers traveling with `od_visualization_use_log_scale`. The result for dataset toy is depicted in Figure 4.6.

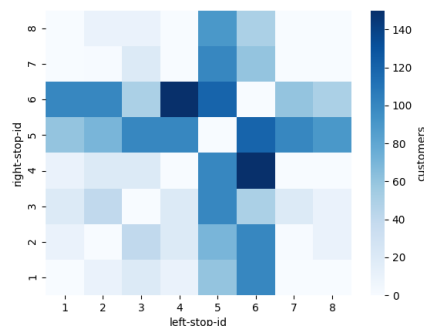


Figure 4.6: The OD data of the toy dataset visualized as a heatmap

4.11.3 Loads

To create an illustration of the traffic loads in the PTN run

```
R make ptn-load-draw
```

Displayed are the links whose traffic load in relation to the maximal traffic load in the network is within the interval given by the fractions `loads_graph_lower_bound` and `loads_graph_upper_bound`. The traffic loads can be illustrated using the edge color or the edge width of the PTN. This can be chosen using `loads_graph_use_edge_color`. The result of the former for dataset toy is depicted on the left hand side of Figure 4.7, whereas the result of the latter is depicted on the right hand side of Figure 4.7. In the

latter case, the maximal edge width can be scaled by adapting `loads_graph_max_edge_width`. The entire figure can be scaled by adapting `loads_draw_conversion_factor`

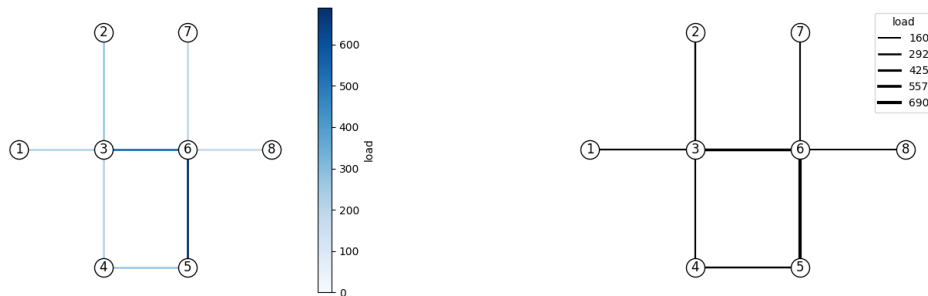


Figure 4.7: The traffic loads of the toy dataset. On the left hand side, the load of an edge is indicated by its width, on the right hand side by its color

4.11.4 Line pool

To create an illustration of the line pool run

```
R> make_lpool_line_pool_draw
```

The result for dataset toy is depicted in Figure 4.8.

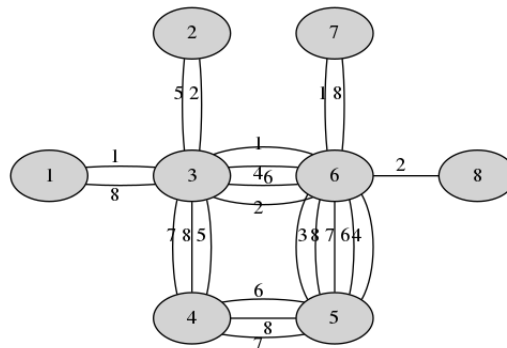


Figure 4.8: The line pool of the toy dataset

The graph can be scaled by adapting `lpool_coordinate_factor`.

4.11.5 Line concept

To create an illustration of the line concept run

```
R> make_lc_line_concept_draw
```

The result for dataset toy is depicted in Figure 4.9.

The graph can be scaled by adapting `lpool_coordinate_factor`.

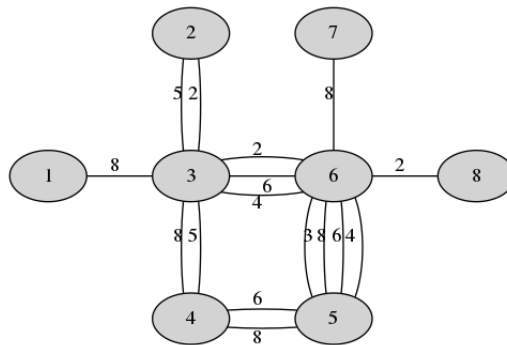


Figure 4.9: One possible line concept of the toy dataset

4.11.6 Timetable

To create an illustration of the timetable, run

```
R make tim-timetable-draw
```

The result for dataset toy is depicted in Figure 4.10. Note, that this command will draw only the ean, if no

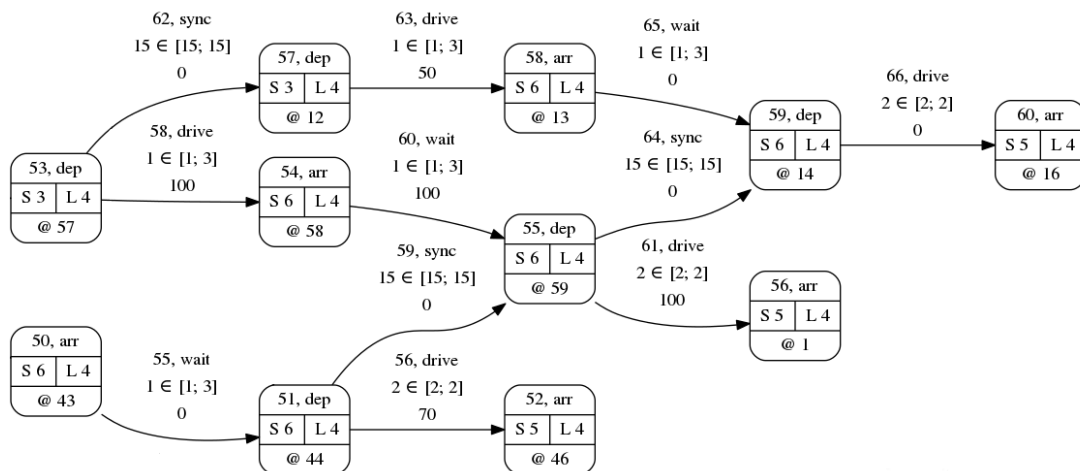


Figure 4.10: Extract of one possible timetable of the toy dataset

timetable is present.

4.11.7 Disposition timetable

To create an illustration of the disposition timetable, run

```
R make dm-disposition-timetable-draw
```

The result for dataset toy is depicted in Figure 4.11. Delayed events will be displayed in red (more delay results in more saturation). Note, that this command will draw only the extended timetable, if no disposition timetable is present.

4.11.8 Tariff

Running

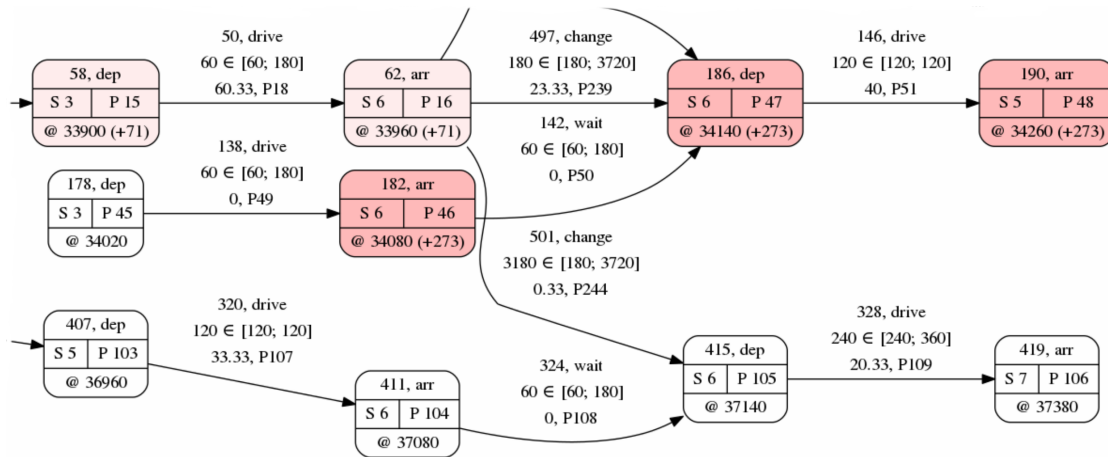


Figure 4.11: Extract of one possible disposition timetable of the toy dataset

`R` `make taf-tariff-draw`

yields a heatmap of prices or of price differences between two specified price matrices and can draw the PTN with nodes assigned to their zones in case of a zone tariff.

Heatmap

If `CK taf_draw_heatmap CV true` (which is the default), executing the above make command generates a heatmap of prices or of price differences for all OD pairs and stores it to `CK filename_tariff_heatmap (Fi graphics/tariff-heatmap.png)`.

Which price matrix or which comparison of price matrices is visualized, is controlled by `CK taf_heatmap_mode` with the following possible values:

- `CV old`, the price matrix specified by `CK taf_evaluate_old_prices` (default: `CV basis/Reference-Price-Matrix`) is visualized,
- `CV new`, the price matrix specified by `CK taf_evaluate_new_prices` (default: `CV tariff/Price-Matrix.taf`) is visualized and
- `CV compare`, the price differences between the price matrices specified by `CK taf_evaluate_new_prices` and `CK taf_evaluate_old_prices` are visualized such that the heatmap shows the change from the old prices to the new prices, i.e. the old prices are subtracted from the new prices.

By default `CK taf_heatmap_mode` is `CV old`.

Furthermore the following features of the heatmap can be controlled:

- `CK taf_heatmap_log_scale` boolean, whether or not the heatmap should use a logarithmic scale. By default it is `CV false`.
- `CK taf_heatmap_use_annotations`, whether or not the heatmap should be annotated with the calculated differences in each square. By default it is `CV false`.

Zones

If `CK taf_draw_zones` is `CV true` (default: `CV false`), then a PTN with stops colored fittingly to their zones is drawn (see Section 4.11.1), which is outputted to `CK filename_tariff_ptn_zone_graph (Fi graphics/ptn-zones.png)`.

An example for the dataset toy with four zones is depicted in Figure 4.12.

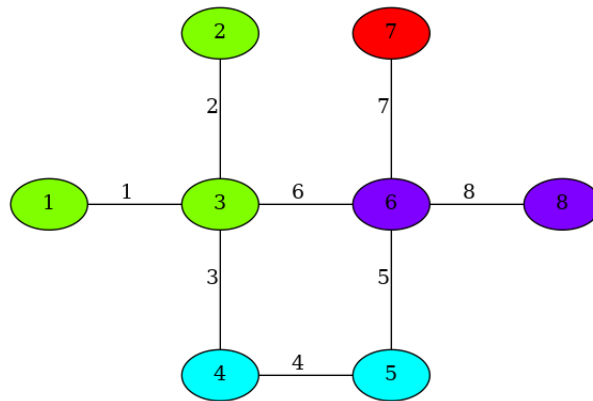


Figure 4.12: The graph of the toy dataset with zones in different colors

4.11.9 mapgui

Additionally, there is an interactive tool for displaying public transportation systems on a map which is used by running

```
[R] make mapgui
```

To decide which step is displayed, set the parameter `[CK] mapgui_show_step` to `[CV] ptn`, `[CV] linepool`, `[CV] lineconcept`, `[CV] timetable` or `[CV] dispotimetable`, respectively. The speed of the visualization is controlled by `[CK] mapgui_visual_speed`.

4.12 Interaction with VISUM

During the work on DFG FOR 2083¹, a basic interface to PTV VISUM [26] was created. For this, `LINTIM` gained the ability to write the periodic timetable in a format that can be easily read by VISUM, as well as reading different infrastructure and solution information from VISUM-net-files. In this section, we will describe the different interfaces and their file requirements. Note that the name of the transport system to read can be set by `[CK] visum_tsyscode`, which defaults to “B”. In this documentation, all attributes will include this default in their name when necessary but the read attributes are dependent on the config key.

4.12.1 Writing files for VISUM

By calling

```
[R] make tim-transform-to-visum
```

`LINTIM` will create a timetable file based on stops (or stop points in VISUM) at

```
[CK] default_timetable_visum_file ([Fi] timetabling/ Timetable-visum-nodes.tim)
```

, that can be read easier by VISUM.

4.12.2 Reading a config file

By calling

```
[R] make config-fill-config-from-visum
```

`LINTIM` will read a visum configuration file provided for `LINTIM` and set the contained config parameters in the `LINTIM` config file `[Fi] basis/Config.cnf`. It will read `[CK] filename_visum_config_file ([Fi] config.net)`. The following parameters will be read

¹<https://for2083.mathematik.uni-kl.de>

LINTIM_BASE_UNIT_FOR_HEADWAY : The system frequency to use, i.e., the common frequency divisor for all line frequencies. Will set `lc_common_frequency_divisor`.

LINTIM_DEFDWELLTIME the default minimal waiting time at each station, will set `ean_default_minimal_waiting_time`.

LINTIM_MIN_TRANSFERTIME the default minimal transfer time at each station, will set `min_change_time`

LINTIM_PERIOD_LENGTH the period length in time units to use. Will set `period_length`.

LINTIM_POSTPREPTIME the turnover time after each line serving. One part of `vs_turn_over_time`, i.e., the values of **LINTIM_POSTPREPTIME** and **LINTIM_PREPREPTIME** will be summed up.

LINTIM_PREPREPTIME the turnover time before each line serving. One part of `vs_turn_over_time`, i.e., the values of **LINTIM_POSTPREPTIME** and **LINTIM_PREPREPTIME** will be summed up.

LINTIM_TIME_UNITS_PER_MINUTE the time units per minute to use. Will set `time_units_per_minute`.

LINTIM_TRANSFER_UTILITY the change penalty to use, i.e., the additional penalty to add for each transfer. Will set `ean_change_penalty`.

LINTIM_TSYS_FOR_ADAPTING the public transport mode to adapt. Will determine, which set of ptn links/infrastructure edges from Visum will be set to usable/forbidden in LinTim. Will set `visum_tsyscode`.

LINTIM_WALKTIME_UTILITY the walk time utility, i.e., the penalty factor for time spend walking. Will set `gen_walking_utility`.

SCENARIO_NAME the name of the dataset, will set `ptn_name`.

4.12.3 Reading the infrastructure

By calling

```
 make ptn-read-infrastructure-from-visum
```

LINTIM will read the infrastructure information on node-level from the provided VISUM-net-file and the corresponding walking information. Note that this will not create a PTN but the underlying infrastructure, i.e., you need to compute the PTN yourself. Whether the walking information is assumed to be symmetric is dependent on `sl_walking_is_directed`. The following files and contents will be read:

filename_net_file (**infrastructure.net**) the infrastructure file with the following objects and attributes

\$ NODE: NO, XCOORD, YCOORD

\$ LINK: FROMNODENO, LENGTH, NO, TONODENO, TSYSSET, T_PUTSYS(B)

filename_visum_walk_file (**walk_times.att**) the walking file with the following objects and attributes

\$ ODDPAIR: FROMZONENO, TOZONENO, WALK_TIME (note that any third attribute will be interpreted as the walk time and only three attributes are allowed here!)

The following files will be written:

- **CK** `filename_node_file` (**Fi** `basis/Node.giv`): The nodes will contain the original visum node number as name.
- **CK** `filename_infrastructure_edge_file` (**Fi** `basis/Edge-Infrastructure.giv`)
- **CK** `filename_walking_edge_file` (**Fi** `basis/Edge-Walking.giv`)

4.12.4 Reading the PTN

By calling

R `make ptn-read-ptn-from-visum`

LINTIM will read the infrastructure information regarding the PTN from the provided VISUM-net-file. Note that the read infrastructure needs to represent a valid LINTIM PTN, i.e., links may only include nodes that are stop points. The following files and contents will be read:

CK `filename_net_file` (**Fi** `infrastructure.net`) the infrastructure file with the following objects and attributes

\$ NODE: NO, XCOORD, YCOORD

\$ STOPPOINT: NO, NODENO

\$ LINK: FROMNODENO, LENGTH, NO, TONODENO, TSYSSET, T_PUTSYS(B)

The following files will be written:

- **CK** `default_stops_file` (**Fi** `basis/Stop.giv`): The stops will contain the original visum node number as short and long name.
- **CK** `default_edges_file` (**Fi** `basis/Edge.giv`)

4.12.5 Reading the demand

Reading stop demand

By calling

R `make od-read-stop-od-from-visum`

LINTIM will read the demand data for the current stops from the provided VISUM-net-file. This step will assume that all zones in the demand matrix are located and named by their corresponding stopping point, which should be present in the short name of the LINTIM stops. Demand from and to the same zone will be ignored and set to 0. The following files and contents will be read:

CK `filename_visum_od_file` (**Fi** `od.att`) the demand file with the following objects and attributes

\$ ODPAIR: FROMZONENO, TOZONENO, DEMAND (note that any third attribute will be interpreted as the demand and only three attributes are allowed here!)

The following file will be written:

- **CK** `default_od_file` (**Fi** `basis/OD.giv`)

Reading node demand

By calling

`make od-read-node-od-from-visum`

LINTIM will read the demand data for the nodes from the provided VISUM-net-file. This step will assume that all zone numbers correspond to the original visum node numbers which should be stored in the names of the LINTIM nodes. The following files and contents will be read:

`filename_visum_od_file` (`od.att`) the demand file with the following objects and attributes

\$ **ODPAIR:** FROMZONENO, TOZONENO, DEMAND (note that any third attribute will be interpreted as the demand and only three attributes are allowed here!)

The following file will be written:

- `filename_od_nodes_file` (`basis/OD-Node.giv`)

4.12.6 Reading stops and lines

For a given infrastructure network and demand, i.e., nodes, infrastructure edges and a node-based demand, given VISUM stops and lines can be read by calling

`make lc-read-stops-and-lines-from-visum`

This step will assume that the original visum node numbers are stored in the names of the LINTIM nodes and that the read lines are undirected. The following files and contents will be read:

`filename_visum_timetable_file` (`vehicle_journeys.att`) the vehicle journey file with the following objects and attributes:

\$ **VEHJOURNEYITEM:** DEP, DIRECTIONCODE, INDEX, LINENAME, TIMEPROFILEITEM\LINEROUTEITEM\STOPPOINT\NO, VEHJOURNEYNO

The following files will be written:

- `default_stops_file` (`basis/Stop.giv`): The stops will contain the original visum node number as short and long name.
- `default_edges_file` (`basis/Edge.giv`)
- `default_pool_file` (`basis/Pool.giv`)
- `default_pool_cost_file` (`basis/Pool-Cost.giv`)
- `default_lines_file` (`lineplanning/Line-Concept.lin`)

4.12.7 Reading a timetable

For a given line concept a timetable for the same lines can be read from provided VISUM-net-files by calling

`make tim-read-timetable-from-visum`

This step will assume that the lines for the VISUM timetable are the same as in the current line concept, including the frequencies but excluding the direction, i.e., LINTIM and VISUM may have the same lines noted in different directions, since lines are assumed to be undirected. The original VISUM node numbers are assumed to be present in the short names of the stops. This method will read the timetable in one specific hour, given by `visum_hour_to_consider`. The corresponding periodic EAN is assumed to be present. The following files will be read:

filename_visum_timetable_file (**vehicle_journeys.att**) the vehicle journey file with the following objects and attributes:

\$ VEHJOURNEYITEM: ARR, DEP, DIRECTIONCODE, INDEX, LINENAME,
TIMEPROFILEITEM\LINEROUTEITEM\STOPPOINT\NO, VEHJOURNEYNO

The following file will be written:

- **default_timetable_periodic_file** (**timetabling/Timetable-periodic.tim**)

4.12.8 Reading fixed lines

By calling

make lc-read-fixed-lines-from-visum

LINTIM will read lines to fix from the provided VISUM-net-file. For this, we assume that there is a transportation system that should be optimized (given by **visum_tsyscode**) and other fixed transportation systems. All fixed lines are read and added to the line pool as well as a fixed line file with their respective frequency and the corresponding capacities. Note that this will change the line pool, i.e., running this multiple times in a row without resetting the pool may lead to unintended consequences. We assume that the short name of the LINTIM stops contains the original VISUM node number.

Afterwards, setting **lc_respect_fixed_lines** to **true** will respect these lines for the line planning problem. This is not supported for all line planning problems, see the corresponding line planning documentation in Section 3.3.

The following file and contents will be read:

filename_net_fixed_lines_file (**visum-fixed-lines.net**) the infrastructure file with the following objects and attributes

\$ LINE: NAME, TSYSCODE
\$ LINEROUTEITEM: DIRECTIONCODE, LINENAME, NODENO
\$ LINK: FROMNODENO, NO, TONODENO
\$ VEHJOURNEY: DEP, LINENAME
\$ VEHUNIT: TOTALCAP, TSYSSET

The following files will be written:

- **filename_lc_fixed_lines** (**line-planning/Fixed-Lines.lin**) the fixed lines
- **filename_lc_fixed_line_capacities** (**line-planning/Line-Capacities.lin**) the capacities of the fixed lines

4.12.9 Reading fixed times

By calling

make tim-read-fixed-times-from-visum

LINTIM will read the timetable of some fixed lines from the provided VISUM-net-file. For this, we assume that there is a transportation system that should be optimized (given by **visum_tsyscode**) and other fixed transportation systems. The fixed lines are assumed to be included in the event-activity-network and the corresponding times will be read.

Afterwards, setting **tim_respect_fixed_times** to **true** will respect these times for the timetabling problem. For more information, see Section 3.6.8.

The following file and contents will be read:

filename_net_fixed_lines_file (**visum-fixed-lines.net**) the infrastructure file with the following objects and attributes

\$ **LINK:** FROMNODENO, NO, TONODENO

\$ **LINEROUTEITEM:** DIRECTIONCODE, LINENAME, NODENO

\$ **TIMEPROFILE:** ARR, DEP, DIRECTIONCODE, LINENAME

The following file will be written:

- **filename_tim_fixed_times** (**timetabling/Fixed-timetable-periodic.tim**) the fixed times

Chapter 5

Evaluation

5.1 Evaluation of the PTN Created by Stop Location

To evaluate the properties of the public transportation network created by stop location, you can use the makefile target

R `make sl-evaluate`

The following parameters will be evaluated and written to

CK `default_statistic_file` (**Fi** `statistic/statistic.sta`):

SK `ptn_feasible_od` For every OD pair exists a path through the PTN. (Only evaluated if an OD matrix exists.)

SK `ptn_feasible_sl` Every demand point that is no more than **CK** `sl_radius` away from the PTN is covered by a stop, i.e., it is no more than **CK** `sl_radius` away from a stop.

SK `ptn_time_average` Average travel-time of all passengers on shortest path through the PTN in seconds. (Only evaluated if an OD matrix exists.)

SK `ptn_obj_stops` Number of stops.

SK `ptn_prop_edges` Number of undirected edges for an undirected PTN, number of directed edges for a directed PTN.

SK `ptn_prop_existing_stops` Number of stops before a stop location algorithm was executed.

SK `ptn_prop_existing_edges` Number of undirected edges for an undirected PTN, number of directed edges for a directed PTN before a stop location algorithm was executed.

SK `ptn_prop_demand_point` Number of demand points.

SK `ptn_prop_relevant_demand_point` Number of demand points that are no more than **CK** `sl_radius` away from the PTN.

SK `ptn_travel_time_realistic` Sum of the realistic travel-time on all edges of the PTN in seconds considering the acceleration (**CK** `sl_acceleration`) and deceleration (**CK** `sl_deceleration`) of the vehicles.

SK `ptn_travel_time_const` Sum of the travel-time on all edges of the PTN in seconds assuming the vehicles would maintain a constant speed of **CK** `gen_vehicle_speed`.

If

C `sl_eval_extended; true`

is set, the following parameters will additionally be evaluated:

SK **ptn_max_distance** Maximal distance any demand point has to the stop nearest to it.

SK **ptn_candidates** Number of candidates considered as new stops during the stop location algorithm.

5.2 Evaluation of the PTN

To evaluate the properties of the public transportation network, you can use the makefile target

R `make ptn-evaluate`

The following parameters will be evaluated and written to

CK **default_statistic_file** (Fi `statistic/statistic.sta`):

SK **ptn_feasible_od** For every OD pair exists a path through the PTN. (Only evaluated if an OD matrix exists.)

SK **ptn_time_average** Average travel-time of all passengers on shortest path through the PTN. (Only evaluated if an OD matrix exists.)

SK **ptn_obj_stops** Number of stops.

SK **ptn_prop_edges** Number of undirected edges for an undirected PTN, number of directed edges for a directed PTN.

5.3 Evaluation of the OD Matrix

To evaluate the properties of the origin destination matrix, you can use the makefile target

R `make od-evaluate`

The following parameters will be evaluated and written to

CK **default_statistic_file** (Fi `statistic/statistic.sta`):

SK **od_prop_entries_greater_zero** Number of entries greater than zero, i.e., of OD pairs (A, B) where more than zero passengers want to travel from A to B .

SK **od_prop_overall_sum** Sum over all entries in the matrix, i.e., all passengers who want to travel in the network.

5.4 Evaluation of the Line Pool

To evaluate the properties of the line pool, you can use the makefile target

R `make lpool-line-pool-evaluate`

The following parameters will be evaluated and written to

CK **default_statistic_file** (Fi `statistic/statistic.sta`):

SK **lpool_cost** $\sum_{l \in \mathcal{L}} cost_l$ - sum over costs per line.

SK **lpool_feasible_circles** No line is containing a circle.

SK **lpool_feasible_od** For every passenger there exists a path through the PTN that is only using edges occurring in the line pool.

SK **lpool_prop_directed_lines** Number of directed lines.

SK **lpool_time_average** Average travel-time of all passengers on shortest path through the PTN where only edges occurring in the line pool are used.

5.5 Evaluation of the Line Concept

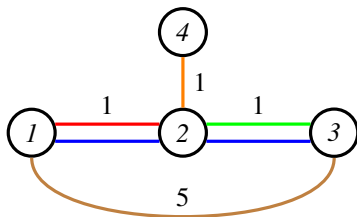
To evaluate the properties of the line concept, you can use the makefile target

R `make lc-line-concept-evaluate`

The following parameters will be evaluated and written to

CK `default_statistic_file` (**Fi** `statistic/statistic.sta`). Some of them (travel times, changes) depend on the routes of the passengers, which can be computed in different ways. Their meaning is illustrated with the following example.

Example.



CK `period_length`: **CV** 60
CK `ean_default_minimal_waiting_time`: **CV** 1
CK `ean_default_maximal_waiting_time`: **CV** 3
CK `ean_default_minimal_change_time`: **CV** 3
CK `ean_model_weight_wait`: **CV** AVERAGE_WAITING_TIME
CK `ean_model_weight_change`: **CV** FORMULA_1
CK `ean_change_penalty`: **CV** 5
CK `gen_passengers_per_vehicle`: **CV** 1

The numbers at the arcs indicate the driving times. Assume that the red line is operated at frequency 2 and all other lines have a frequency of 1. Let there be a demand of 2 between nodes 1 and 3 and of 1 between nodes 1 and 4.

SK `lc_average_distance` Average Euclidean distance between the two endpoints of a line.

SK `lc_average_edges/length` Average number of edges/length of the lines in the line concept.

SK `lc_cost` Sum over costs of line times frequency: $\sum_{l \in \mathcal{L}} cost_l f_l$.

SK `lc_feasible` Lower and upper bounds on frequency on every edge respected: $f_e^{min} \leq \sum_{l \in \mathcal{L}} f_l \leq f_e^{max}$.

SK `lc_min_distance/edges/length` Minimal distance/number of edges/length of the lines in the line concept.

SK `lc_obj_game` Sum of the squared frequencies on all edges: $\sum_{e \in E} f_e^2$.

SK `lc_prop_directed_lines` number of directed lines. If a line is undirected, it is counted twice.

SK `lc_prop_freq_max` The maximal frequency: $\max_{l \in \mathcal{L}} f_l$.

SK `lc_var_distance/edges/length` Variance of the distance/number of edges/length of the lines in the line concept.

For the following two properties, passengers are routed along shortest paths in the network consisting of all edges that are covered with frequency at least one. Every driving edge contributes its driving time (computed according to **CK** `ean_model_weight_drive`) and every intermediate station contributes the waiting time (computed according to **CK** `ean_model_weight_wait`), independently of whether a change is necessary. Hence, the passenger routing respects neither vehicle capacities nor changes.

SK `lc_time_average_without_transfers` Average travel-time of all passengers in the routing described above, where every driving edge contributes its driving time and every intermediate station contributes the waiting time, independently of whether a change is necessary.

SK lc_unicapacitated_direct_travelers Number of travelers that have a shortest path in the routing described above that does not require a transfer. Does not respect the changing times, change penalty, and the capacity of lines, so this is not the same as the objective of the direct travelers model, see Section 3.3.2. For this, please check **SK lc_obj_direct_travelers** in the extended evaluation.

In the example the path 1,2,3 is has cost $1 + 2 + 1 = 4 < 5$. Therefore, both passengers between 1 and 3 are routed along this path and are counted as direct passengers. Moreover, the passenger from 1 to 4 is routed along the path 1,2,4, which also has travel time 4 but cannot be followed without changing. Therefore, **SK lc_time_average_without_transfers** is **SV** 4 and **SK lc_unicapacitated_direct_travelers** is **SV** 2.

When setting the config-parameter **CK lc_eval_extended** to *true*, additionally the following properties will be evaluated and written to **CK default_statistic_file** (**Fi** statistic/statistic.sta). Note that an IP solver is necessary for that. The IP solver used is selected via **CK lc_solver**.

For the following two properties, passengers are routed along shortest paths in the Change&Go graph, where every change contributes the transfer penalty (**CK ean_change_penalty**) plus the transfer time calculated according to **CK ean_model_weight_change**. Hence the vehicle capacities are not respected.

SK lc_perceived_time_average Average travel time including a penalty for each transfer (**CK ean_change_penalty**) in the passenger routing described above.

SK lc_prop_changes Total number of transfers in the passenger routing described above.

In the example, the transfer time between the red and one of the other two lines is computed according to Formula 1 (see Section 7.9) as $\frac{60}{2} + \frac{60}{1} + 5 = 95$. Since capacities are ignored, the passengers between 1 and 3 are routed via the blue line with perceived travel time 4, while the passenger between 1 and 4 has a perceived travel time of 97. Therefore, **SK lc_perceived_time_average** is $\frac{2 \cdot 4 + 97}{3} = \text{SV } 35$ and **SK lc_prop_changes** is **SV** 1.

For the following property, passengers are routed according to a minimum-cost multi-commodity flow in the Change&Go graph, respecting the capacities and the transfer times, taken as **CK ean_default_minimal_change_time** plus **CK ean_change_penalty**.

SK lc_obj_travel_time Average travel time including a transfer penalty (**CK ean_change_penalty**) in the routing described above. This is the objective value of the current solution the travel time model, see Section 3.3.4.

In the example the transfer times are estimated as $3 + 5 = 8$. Only one passenger between 1 and 3 can be routed via the blue line. The other will take the brown line with travel time $5 < 10 = 1 + 8 + 1$. The passenger between 1 and 4 has a travel time of 10. Therefore, the **SK lc_obj_travel_time** is $\frac{4 + 5 + 10}{3} = \text{SV } 6.33333$.

For the following two properties, passengers are routed according to a minimum-cost multi-commodity flow in the Change&Go graph, respecting the capacities and the transfer times, computed according to **CK ean_model_weight_change** plus **CK ean_change_penalty**.

SK lc_capacitated_perceived_time_average Average travel time including a transfer penalty (**CK ean_change_penalty**) in the routing described above.

SK lc_capacitated_prop_changes Total number of transfers in the passenger routing described above.

In the example the transfer time between the red line and another line is again estimated as 95. Therefore, **SK lc_capacitated_perceived_time_average** is $\frac{4 + 5 + 97}{3} = \text{SV } 35.33333$ and **SK lc_capacitated_prop_changes** is **SV** 1.

For the following property, the passengers are routed according to a minimum-cost path-based multi-commodity flow in the Change&Go graph, where every origin-destination pair (commodity) is restricted to paths that correspond to a shortest path in the PTN (with the length of a path including the waiting time at every intermediate station).

SK lc_obj_direct_travelers Total number of passengers travelling without transfer in the passenger routing described above. This is the objective of the direct travelers model for the current solution, see Section 3.3.2.

In the example, the passengers between 1 and 3 can only be routed along the path 1,2,3, which has length $4 < 5$ in the PTN. Therefore, one passenger uses the blue line and the other changes at 2. Also the passenger between 1 and 4 changes at 2. Therefore, **SK lc_obj_direct_travelers** is only **SV** 1.

5.6 Evaluation of the EAN

To evaluate the properties of the event activity network, you can use the makefile target

R make ean-evaluate

The following parameters will be evaluated and written to

CK default_statistic_file (**Fi** statistic/statistic.sta):

SK ean_prop_events $|\mathcal{E}|$ - number of events.

SK ean_prop_events_arrival $|\{e \in \mathcal{E} : e \text{ is arrival}\}|$ - number of arrival events.

SK ean_prop_events_departure $|\{e \in \mathcal{E} : e \text{ is departure}\}|$ - number of departure events.

SK ean_prop_activities $|\mathcal{A}|$ - number of activities.

SK ean_prop_activities_change $|\mathcal{A}_{change}|$ - number of change activities.

SK ean_prop_activities_drive $|\mathcal{A}_{drive}|$ - number of drive activities.

SK ean_prop_activities_wait $|\mathcal{A}_{wait}|$ - number of wait activities.

SK ean_prop_activities_headway $|\mathcal{A}_{headway}|$ - number of headway activities.

SK ean_prop_activities_od $|\{a \in \mathcal{A} : c_a > 0\}|$ - number of activities with more than 0 passengers.

SK ean_prop_activities_od_change $|\{a \in \mathcal{A}_{change} : c_a > 0\}|$ - number of change activities with more than 0 passengers.

SK ean_prop_activities_od_drive $|\{a \in \mathcal{A}_{drive} : c_a > 0\}|$ - number of drive activities with more than 0 passengers.

SK ean_prop_activities_od_wait $|\{a \in \mathcal{A}_{wait} : c_a > 0\}|$ - number of wait activities with more than 0 passengers.

SK ean_time_average $\frac{1}{\sum_{a \in \mathcal{A}} c_a} \sum_{a \in \mathcal{A}} c_a \cdot$ "duration assumption" - estimated average travel time. For duration assumption see 4.5.

Furthermore by setting config-parameter **CK** ean_eval_extended to *true* additionally the following parameter will be evaluated and written to **CK** default_statistic_file (**Fi** statistic/statistic.sta):

SK ean_prop_activities_feas $|\{a \in \mathcal{A} : U_a - L_a < T - 1\}|$ - number of activities that impose constraints.

SK ean_prop_activities_objective $|\{a \in \mathcal{A} : c_a > 0 \text{ or } U_a - L_a < T - 1\}|$ - number of activities that have an influence on the objective value.

SK ean_prop_changes_od_max $\max_{\substack{a \in \mathcal{A}_{change} \\ c_a > 0}}$ "duration assumption of a" - maximal used change duration.

SK ean_prop_changes_od_min $\min_{\substack{a \in \mathcal{A}_{change} \\ c_a > 0}}$ "duration assumption of a" - minimal used change duration.

SK ean_prop_headways_dep Are headways between departures only.

SK ean_prop_headways_interstation Do headways exist between different stations.

Additionally, the loads on the ean will be evaluated and compared to the maximal feasible load on the ptn edges given by the line concept. If the load on the ptn is invalid, i.e., too high, the respective ptn edges and their load will be written to **CK filename_invalid_loads** (**FI** statistic/Invalid-Loads.sta).

Additionally, the maximal load factor will be written as **SK ean_max_load_factor** to

CK default_statistic_file (**FI** statistic/statistic.sta).

5.7 Evaluation of the Timetable

To evaluate the properties of the timetable, you can use the makefile target

R make tim-timetable-evaluate

The following parameters will be evaluated and written to

CK default_statistic_file (**FI** statistic/statistic.sta):

SK tim_feasible $L_a \leq ((\pi_j - \pi_i - L_a) \bmod T) + L_a \leq U_a$ for all $(i, j) = a \in \mathcal{A}$ - Are lower and upper bounds on travel time on each activity respected.

SK tim_obj_ptt1 $\sum_{(i,j)=a \in \mathcal{A}} c_a \left(((\pi_j - \pi_i - L_a) \bmod T) + L_a \right)$ - Sum of weighted travel time. Weights correspond to the number of passengers specified in activity file.

SK tim_obj_slack_average $\frac{1}{|\mathcal{A}|} \sum_{(i,j)=a \in \mathcal{A}} ((\pi_j - \pi_i - L_a) \bmod T)$ - Average of slacks.

SK tim_time_average - Average travel time per passenger. The travel time for every OD pair is calculated according to its shortest path in the EAN.

SK tim_perceived_time_average - Average travel time per passenger. The travel time for every OD pair is calculated according to its shortest path in the EAN with additionally **CK ean_change_penalty** on change activities.

Furthermore by setting config-parameter **CK tim_eval_extended** to *true* additionally the following parameter will be evaluated and written to **CK default_statistic_file** (**FI** statistic/statistic.sta):

SK tim_obj_slack_drive_average $\frac{1}{|\mathcal{A}_{drive}|} \sum_{(i,j)=a \in \mathcal{A}_{drive}} ((\pi_j - \pi_i - L_a) \bmod T)$ - average slack on drive activities.

SK tim_obj_slack_wait_average $\frac{1}{|\mathcal{A}_{wait}|} \sum_{(i,j)=a \in \mathcal{A}_{wait}} ((\pi_j - \pi_i - L_a) \bmod T)$ - average slack on wait activities.

SK tim_obj_slack_change_average $\frac{1}{|\mathcal{A}_{change}|} \sum_{(i,j)=a \in \mathcal{A}_{change}} ((\pi_j - \pi_i - L_a) \bmod T)$ - average slack on change activities.

SK tim_obj_slack_headway_average $\frac{1}{|\mathcal{A}_{headway}|} \sum_{(i,j)=a \in \mathcal{A}_{headway}} ((\pi_j - \pi_i - L_a) \bmod T)$ - average slack on headway activities.

SK tim_overcrowded_time_average the average time that passengers are overcrowded in the vehicles. Does not take any rerouting into account, i.e., will use the passenger weights currently stored in the EAN. A drive or wait activity is overcrowded, if the number of passengers using the activity is larger than **CK gen_passengers_per_vehicle**.

tim_prop_changes_od_max $\max_{\substack{(i,j)=a \in \mathcal{A}_{change} \\ c_a > 0}} (\pi_j - \pi_i) \bmod T$ - maximal used change duration.

tim_prop_changes_od_min $\min_{\substack{(i,j)=a \in \mathcal{A}_{change} \\ c_a > 0}} (\pi_j - \pi_i) \bmod T$ - minimal used change duration.

tim_number_of_transfers Weighted number of transfers.

5.7.1 Capacitated Routing

There is also the possibility to do a capacitated routing of the passengers in the EAN by solving a multicommodity flow problem. To do so, set the config-parameter

tim_cap_eval **true**.

First, the standard timetable evaluation is executed and written to the statistics-file. The algorithm then creates a routing network from the EAN. For each Stop $s \in \mathcal{S}$ we add two additional nodes to the EAN, one origin event and one destination event. We add arcs from the origin event to all departure events corresponding to stop s and from all arrival events of s to the destination event, those activities are added to \mathcal{A} . Let $\text{origin}(s)$ be the origin event to the stop s and $\text{dest}(s)$ the corresponding destination event. The set of all origin events is denoted by \mathcal{E}_{origin} and the set of all destination events by \mathcal{E}_{dest} .

Let $\text{duration}(a)$ be the duration of activity a , which comes from the existing timetable. If the config-parameter **tim_cap_eval_tt** is false, we don't use the timetable but the lower bound of the activity. If a is a change activity, we add a penalty time **ean_change_penalty**. Furthermore we need the capacity C of each vehicle, given by **gen_passengers_per_vehicle**.

For each activity a and each stop s there is a variable $x_{s,a}$ that states how many passengers starting from stop s use activity a . Depending on **tim_cap_eval_integer_flow** they are integer or continuous variables. We set up the following optimization problem and solve it with Gurobi.

$$\min \sum_{s \in \mathcal{S}} \sum_{a \in \mathcal{A}} \text{duration}(a) \cdot x_{s,a} \quad (5.1a)$$

$$\text{s.t.} \quad \sum_{\substack{a \in \mathcal{A} \\ a \in \delta^+(e)}} x_{s,a} - \sum_{\substack{a \in \mathcal{A} \\ a \in \delta^-(e)}} x_{s,a} = \sum_{u \in \mathcal{S}} C_{s,u} \quad \forall e \in \mathcal{E}_{origin}, \forall s \in \mathcal{S}, \quad (5.1b)$$

$$\sum_{\substack{a \in \mathcal{A} \\ a \in \delta^-(e)}} x_{s,a} - \sum_{\substack{a \in \mathcal{A} \\ a \in \delta^+(e)}} x_{s,a} = C_{s,t} \quad \forall e \in \mathcal{E}_{dest}, \forall s \in \mathcal{S} : e = \text{dest}(t), \quad (5.1c)$$

$$\sum_{\substack{a \in \mathcal{A} \\ a \in \delta^+(e)}} x_{s,a} - \sum_{\substack{a \in \mathcal{A} \\ a \in \delta^-(e)}} x_{s,a} = 0 \quad \forall e \in \mathcal{E}_{arr} \cup \mathcal{E}_{dep}, \forall s \in \mathcal{S}, \quad (5.1d)$$

$$\sum_{s \in \mathcal{S}} x_{s,a} \leq C \quad \forall a \in \mathcal{A}_{drive}, \quad (5.1e)$$

$$x_{s,a} \in \{0, \dots, C\} \subseteq \mathbb{N} \quad \forall s \in \mathcal{S}, a \in \mathcal{A} \quad (5.1f)$$

Here, $\delta^+(e)$ is the set of all outgoing arcs of the node e and $\delta^-(e)$ is the set of all ingoing arcs. $C_{u,v}$ denotes the OD-value from stop u to stop v , i.e. the number of passengers that start from the stop u and end at stop v according to the given OD-Matrix. If **tim_cap_eval_integer_flow** is false, constraint (5.1f) is replaced by $x_{s,a} \in [0, C] \subseteq \mathbb{R}$.

If **tim_cap_eval_accumulate_on_links** is true, constraint (5.1e) is replaced by

$$\sum_{\substack{s \in \mathcal{S} \\ a \in \text{activities}(l)}} x_{s,a} \leq C \cdot |\text{activities}(l)| \quad \forall l \text{ Link in PTN}$$

where $\text{activities}(l)$ is the set of all activities a that belong to the link l of the PTN. This gives a weaker version of the IP.

The following parameters will be evaluated and written to

default_statistic_file (statistic/statistic.sta):

tim_capacitated_travel_time $\sum_{a \in \mathcal{A}} \text{passengers}(a) \cdot \text{duration}(a)$ - sum of all travel times, where $\text{passengers}(a)$ denotes the number of passengers using activity a according to the optimal solution of the IP.

tim_capacitated_travel_time_average travel time divided by the total number of passengers - average travel time per passenger.

tim_capacitated_percieved_travel_time $\sum_{a \in \mathcal{A}} \text{passengers}(a) \cdot \text{percieved_duration}(a)$ - sum of all percieved travel times. $\text{percieved_duration}(a) = \text{duration}(a) + \text{ean_change_penalty}$ if $a \in \mathcal{A}_{\text{wait}}$ and $\text{percieved_duration}(a) = \text{duration}(a)$ for all other activities a .

tim_capacitated_percieved_travel_time_average percieved travel time divided by the total number of passengers - average percieved travel time per passenger.

tim_capacitated_max_load $\max_{a \in \mathcal{A}} \frac{\text{passengers}(a)}{C}$ - maximal percentage load. Could be greater than 1, if tim_cap_eval_accumulate_on_links is true.

5.8 Evaluation of the Tariff created by Tariff Planning

To evaluate the properties of a tariff created by

make taf-tariff

or the price matrices created by

make taf-tariff-price-matrix

run

make taf-tariff-evaluate

The following configuration parameters control the evaluation:

- taf_evaluate_old_prices, points towards a price matrix. By default this is the reference price matrix basis/Reference-Price-Matrix.giv.
- taf_evaluate_new_prices, points towards a price matrix. By default this is the tariff price matrix tariff/Price-Matrix.taf.

The following parameters will be evaluated and written to

default_statistic_file (statistic/statistic.sta), where they are sorted alphabetically:

taf_revenue_old The revenue generated by the tariff if all passengers pay the prices in taf_evaluate_old_prices.

taf_revenue_new The revenue generated by the tariff if all passengers pay the prices in taf_evaluate_new_prices.

taf_od_pairs_increased_prices The absolute number of OD pairs for which the price increases comparing the prices in taf_evaluate_old_prices to the prices in taf_evaluate_new_prices.

taf_od_pairs_decreased_prices The absolute number of OD pairs for which the price decreases comparing the prices in taf_evaluate_old_prices to the prices in taf_evaluate_new_prices.

taf_passengers_increased_prices The absolute number of passengers for which the price increases comparing the prices in taf_evaluate_old_prices to the prices in taf_evaluate_new_prices.

- SK taf_passengers_decreased_prices** The absolute number of passengers for which the price decreases comparing the prices in **CK taf_evaluate_old_prices** to the prices in **CK taf_evaluate_new_prices**.
- SK taf_objective_sum_unit** The sum of absolute deviations between the new prices (**CK taf_evaluate_new_prices**) and the old prices (**CK taf_evaluate_old_prices**) all weighted with one.
- SK taf_objective_sum_od** The sum of absolute deviations between the new prices (**CK taf_evaluate_new_prices**) and the old prices (**CK taf_evaluate_old_prices**) weighted with the od values.
- SK taf_objective_sum_reference_inverse** The sum of absolute deviations between the new prices (**CK taf_evaluate_new_prices**) and the old prices (**CK taf_evaluate_old_prices**) weighted with the inverse of the old prices.
- SK taf_objective_max_unit** The maximum of absolute deviations between the new prices (**CK taf_evaluate_new_prices**) and the old prices (**CK taf_evaluate_old_prices**) all weighted with one.
- SK taf_objective_max_od** The maximum of absolute deviations between the new prices (**CK taf_evaluate_new_prices**) and the old prices (**CK taf_evaluate_old_prices**) weighted with the od values.
- SK taf_objective_max_reference_inverse** The maximum of absolute deviations between the new prices (**CK taf_evaluate_new_prices**) and the old prices (**CK taf_evaluate_old_prices**) weighted with the inverse of the old prices.

For the calculation of the objective values only non-trivial OD pairs are considered as described in section 7.12. If there are zero prices for non-trivial OD pairs, **SK taf_objective_sum_reference_inverse** and **SK taf_objective_max_reference_inverse** are None.

5.9 Evaluation of the Trips

To evaluate the properties of the trips, you can use the makefile target

R `make ro-trips-evaluate`

The following parameters will be evaluated and written to

CK default_statistic_file (**Fi** `statistic/statistic.sta`):

- SK ro_trips_feasible** whether the trips are feasible. The trips are considered feasible if they cover every event in the aperiodic event activity network and no event is used in multiple trips.
- SK ro_prop_trips** $|\mathcal{T}|$ - number of trips.
- SK ro_prop_stops_at_begin_or_end** Number of stations that are start or end station of a trip.

5.10 Evaluation of the Vehicle Schedules

To evaluate the properties of the vehicle scheduling, you can use the makefile target

R `make vs-vehicle-schedules-evaluate`

This evaluation will read the following parameters from the config-files:

- CK vs_vehicle_costs** The cost of a vehicle, needed to determine the costs
- CK vs_eval_cost_factor_empty_length** the cost of a kilometer on an empty trip
- CK vs_eval_cost_factor_empty_duration** the cost for the vehicle driving on an empty trip for an hour

CK **vs_eval_cost_factor_full_length** the cost of a kilometer serving a line

CK **vs_eval_cost_factor_full_duration** the cost for the vehicle driving for an hour while serving a line

The following parameters will be evaluated and written to

CK **default_statistic_file** (**Fi** `statistic/statistic.sta`):

SK **vs_cost** The cost of the vehicle schedule, weighted according to the parameters above.

SK **vs_feasible** Whether the current vehicle schedule is feasible. This only checks, whether the time for the empty trips is sufficient, not the viability of the covered lines.

SK **vs_circulations** The number of circulations in the vehicle schedule.

SK **vs_vehicles** The number of used vehicles in the vehicle schedule.

SK **vs_empty_distance** The distance a vehicle drives without passengers in the current vehicle schedule, given in kilometers.

SK **vs_empty_distance_with_depot** The distance a vehicle drives without passengers in the current vehicle schedule including driving from and to the depot, given in kilometers. Will be the same as above if the depot index is not set.

SK **vs_empty_duration** The time needed for empty trips in the current vehicle schedule, given in minutes. Does not include waiting in stations.

SK **vs_empty_duration_with_depot** The time needed for empty trips in the current vehicle schedule including driving from and to the depot, given in minutes. Does not include waiting in stations. Will be the same as above if the depot index is not set.

SK **vs_empty_trips** The number of empty trips in the current vehicle schedule. Does not include waiting in stations.

SK **vs_empty_trips_depot** The number of empty trips to and from the depot.

SK **vs_minimal_waiting_time** The minimal waiting time in a station between two consecutive trips, served by the same vehicle. Only if the station is not changed in the empty trip.

SK **vs_maximal_waiting_time** The maximal waiting time in a station between two consecutive trips, served by the same vehicle. Only if the station is not changed in the empty trip.

SK **vs_average_waiting_time** The average waiting time in a station between two consecutive trips, served by the same vehicle. Only if the station is not changed in the empty trip.

SK **vs_full_distance** The distance a vehicle drives with passengers in the current vehicle schedule, given in kilometers.

SK **vs_full_duration** The time needed for serving trips in the current vehicle schedule, given in minutes.

5.11 Evaluation of the Disposition Timetable

To evaluate the properties of the delay management, you can use the makefile target

R `make dm-disposition-timetable-evaluate`

The following parameters will be evaluated and written to

CK **default_statistic_file** (**Fi** `statistic/statistic.sta`):

dm_feasible Whether the disposition timetable is feasible according to the lower bounds of the activities.

dm_obj_changes_missed_od The number of missed used connections in the disposition timetable.

dm_obj_delay_events_average The average delay of the events in the disposition timetable.

dm_obj_dm2 The objective value of the DM_method DM2.

dm_obj_dm2_average The objective value of DM_method DM2, divided by the number of passengers.

dm_prop_events_delayed The number of delayed events in the disposition timetable.

dm_prop_headways_swapped The number of headways swapped in the disposition timetable, compared to the original timetable.

dm_time_average The average travel time of the passengers according to the disposition timetable.

Furthermore by setting config-parameter **DM_eval_extended** to *true* additionally the following parameters will be written to **default_statistic_file** (**statistic/statistic.sta**). Note, that the rollout must have been done with the parameter **ro_rollout_passenger_paths** set to *true*.

dm_obj_dm1 The objective value of DM_method DM1.

dm_obj_dm1_average The objective value of DM_method DM1, divided by the number of passengers.

dm_passenger_delay The delay of the passenger after rerouting, given the distribution of **DM_passenger_routing_arrival_on_time**.

dm_passenger_delay_average The average delay of the passenger after rerouting, given the distribution of **DM_passenger_routing_arrival_on_time**.

Additionally, when the config-parameter **DM_eval_extended** is set to *true*, the following distributions will be written to **./statistic/statistic_dist.sta**:

dm_dist_delays_events For each possible delay (in seconds) there is one entry giving the number of events with this delay in the disposition timetable.

dm_dist_delays_od For each possible delay (in seconds) there is one entry giving the number of passengers with this delay in the disposition timetable.

Chapter 6

Overview of Supported Integer Programming Solvers

Different algorithms in `LINTIM` use integer programm solvers. Altogether, the following solvers are currently used in `LINTIM`

- Gurobi
- Xpress
- CPLEX
- SCIP
- COIN
- CBC
- GLPK

For an overview, which algorithms support which solver choice, see Table 6.1. For information on how to combine `LINTIM` with one of the solvers above, see Section 1.2.1.

Algorithm	Supported Solvers	Config-Key	Reference
Stop Location dsl, dsl-tt, dsl-tt-2 tt	Xpress Gurobi, Xpress, CPLEX, SCIP, GLPK	<input type="checkbox"/> sl_solver	Section 3.1
Line Pool Generation all	Gurobi, Xpress, CPLEX, SCIP, GLPK	<input type="checkbox"/> lc_solver	Section 3.2
Line Planning Cost CostRestricted CostExtended Direct DirectRelaxation DirectRestricted CostDirect Game Min-Changes Travelling-Time-CG Traveling Time IP	Gurobi, Xpress, CPLEX, SCIP, GLPK Gurobi, Xpress, CPLEX, SCIP, GLPK Gurobi, Xpress, CPLEX, SCIP, GLPK Gurobi, Xpress, CPLEX, SCIP, GLPK Xpress Gurobi, Xpress, CPLEX, SCIP, GLPK Xpress Xpress Xpress Xpress Gurobi	<input type="checkbox"/> lc_solver	Section 3.3
Timetabling IP Aperiodic-robust Cycle-base	Gurobi, Xpress, CPLEX, SCIP, GLPK Xpress Gurobi, Xpress, CPLEX, SCIP, COIN, CBC, GLPK	<input type="checkbox"/> tim_solver	Section 3.6
Tariff Planning Distance Beeline Zone	Gurobi Gurobi Gurobi	<input type="checkbox"/> taf_solver	Section 7.12
Vehicle Scheduling Canal-based IP Line-Based	Xpress Gurobi, Xpress, CPLEX, SCIP, GLPK Xpress	<input type="checkbox"/> vs_solver	Section 3.8
Delay Management IP	Gurobi, Xpress	<input type="checkbox"/> DM_solver	Section 3.9
Integrated Models all	Gurobi, Xpress, CPLEX, SCIP, COIN, CBC, GLPK	<input type="checkbox"/> int_solver	Section 3.10
Tools Line-Rearrange Regenerate-load (spanner-based)	Xpress Gurobi, CPLEX, Xpress	<input type="checkbox"/> load_generator_solver	Section 4.3.5

Table 6.1: Table of all algorithms using an integer programming solver

Chapter 7

Configuration Parameters

This section describes the configuration parameter available in `LINTIM`. For a detailed description of the different algorithms, see Section 3. There, you find a list of corresponding parameters for the different algorithms.

7.1 General

console_log_level the log level to use, determines the amount of output on the console. The possible log levels are:

ERROR: Only write error messages

WARN: Additionally write warnings

INFO: The default. Will give general information about the current step of the algorithm used.

DEBUG: This includes many information to better understand the behavior of the algorithm, e.g., information about substeps of the algorithm, the read configuration values, the read input files, solver output, ...

gen_passengers_per_vehicle the capacity of the vehicles.

gen_walking_utility the penalty factor for walking.

period_length the length of the periodic planning period.

7.2 Stop Location

sl_destruction_allowed whether it is allowed to destroy existing stops

sl_distance the distance function to use

sl_eval_extended activate the extended evaluation

sl_max_walking_time the maximal walking time allowed for passengers

sl_mip_gap the mip optimization gap for the solver, 0.1 equals a gap of 10 % (-1=use default value).

sl_model the model to use. For an overview on all models, see Section 3.1.

sl_radius the covering radius of a stop

sl_solver determine the solver to be used. Note that not all solvers are supported by all models.

sl_threads determine the maximal number of threads to use for the solver (-1=use default value, i.e., no restriction). Note that this will only be used for a possible solver integration of the chosen model, not for the rest of the algorithm.

sl_timelimit the time limit for the solver in seconds (-1=use default value).

sl_write_lp_file whether to write the lp file of the model to solve

7.3 OD

od_draw_conversion_factor scaling factor for the graph visualization of the OD data.

od_visualization_lower_bound percentage of the maximum OD pair value. Lower valued OD pairs will not be displayed in the graph visualization of the OD data.

od_visualization_upper_bound percentage of the maximum OD pair value. Higher valued OD pairs will not be displayed in the graph visualization of the OD data.

od_visualization_use_annotations whether to use annotations in the heatmap visualization of the OD data.

od_visualization_use_edge_colors whether to use the edge color (instead of the edge width) to indicate passenger numbers in the graph visualization of the OD data.

od_visualization_use_heatmap whether to use a heatmap (instead of a graph) for the visualization of the OD data

od_visualization_use_log_scale whether to visualize the logarithmic values of the passenger numbers.

od_visualization_max_edge_width the maximal edge width in points in the graph visualization where the edge width is used to visualize passenger numbers.

7.4 PTN

ptn_draw_use_coordinates whether stop-coordinates are used for plotting the PTN or stops are arranged automatically.

ptn_draw_interactive_graph_edge_labels whether edge labels are displayed in the interactive PTN visualizations

7.5 Line Planning

lc_budget the budget for the line concept, i.e., the maximal weighted sum of the line costs and the computed frequencies.

lc_common_frequency_divisor the common divisor of the frequencies, i.e., a frequency is only allowed if it is a multiple of this value. A value ≤ 0 will test any system frequency (except for 1) and output the best solution.

lc_direct_optimize_costs whether to additionally optimize the costs in the direct model, see Section 3.3.2. When set to **cv** true, the model will optimize a weighted sum of line costs and direct travelers and will use **lc_mult_relation** as a weight.

lc_eval_extended enables the extended evaluation. Needs an IP solver present. For more information, see Section 5.5.

- lc_maximal_frequency** the maximal frequency value allowed
- lc_mult_relation** weighting factor in a convex combination of costs and direct travelers. A value of 0 is equivalent to solving the direct travelers model while a value of 1 is equivalent to solving the cost model, therefore the value should be in [0, 1].
- lc_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10 % (-1=use default value).
- lc_model** the line planning model to use. For an overview of all models, see Section 2.3.
- lc_number_of_possible_frequencies** the maximal number of different frequency values allowed to use.
- lc_respect_fixed_lines** whether to respect fixed lines, i.e., lines with a given frequency
- lc_respect_forbidden_edges** whether to respect forbidden links, i.e., links in the PTN that may not be used by the public transport model currently optimized. This may e.g. be the case when optimizing a bus network and considering a PTN containing train tracks.
- lc_solver** determine the solver to be used. Note that not all solvers are supported by all models.
- lc_threads** determine the maximal number of threads to use for the solver (-1=use default value, i.e., no restriction). Note that this will only be used for a possible solver integration of the chosen model, not for the rest of the algorithm.
- lc_timelimit** the time limit for the solver in seconds (-1=use default value).
- lc_write_lp_file** whether to write the lp file of the model to solve

7.6 Ridepooling

- rpool_model** Model to use for ride pool generation. Available options are **all**, **configvalue-demand_heuristic** and **tree_based**.
- rpool_min_edges** Minimal number of edges in a ridepooling area. The default value is 1.
- rpool_max_edges** Maximal number of edges in a ridepooling area. The default value is 4.
- rpool_costs_fixed** Costs for one vehicle to operate in one ridepooling area.
- rpool_load_factor** Used by the demand heuristic. Only edges where the load is at most this parameter times the line vehicle capacity are added to ridepooling areas.
- rpool_nb_edges_to_trim** Number of edges to trim the lines in the existing linepool.
- rc_model** model to use to ride concept determination.
- rc_solver** specifies the solver to use the IP formulations.
- rc_timelimit** Timelimit in seconds for the solver. Value of -1 means no restriction.
- rc_mip_gap** MIP Gap for the solver. Value of -1 means no restriction.
- rc_threads** Number of threads to use for the solver. Value of -1 means no restriction.
- rc_write_lp_file** Whether to write the IP formulation as a file or not.
- rc_passengers_per_vehicle** Capacity of a single ridepooling vehicle.

7.7 Load Generation

- load_generator_add_additional_load** whether to add additional load per link, given in `filename_additional_load_file` (`basis/Additional-Load.giv`).
- load_generator_fixed_upper_frequency** the fixed upper frequency bound of a link after load generation. Whether this or a factor of the lower frequency bound is used is determined by `load_generator_fix_upper_frequency`
- load_generator_fix_upper_frequency** whether a fixed upper frequency bound (`true`) or a multiple of the lower bound should be used for a link after load generation.
- load_generator_lower_frequency_factor** the factor to multiply the minimal lower frequency bound (given by the capacity of the vehicle) to obtain the new lower frequency bound. The result is rounded up.
- load_generator_max_iteration** determines the number of iterations allowed before the algorithms terminates, if no convergence is observed
- load_generator_min_change_time_factor** the factor to weight the minimal change time (`ean_default_minimal_change_time`) to obtain the change objective function for routing. The change objective function will never be higher than the maximal change time (`ean_default_maximal_change_time`)
- load_generator_model** how to route the passengers
- LOAD_FROM_EAN** use the current weights in the EAN to determine the weights on the PTN links. The EAN has to be present.
 - LOAD_FROM_PTN** determine new passenger routes based on the other parameters given.
 - spanners** determine new passenger routes using a MIP formulation considering building and routing costs and maximum detour factors.
- load_generator_number_of_shortest_paths** the number of shortest paths to use. For every passenger, the given number of shortest paths are computed and the passengers are distributed with a logit model using `load_generator_sp_distribution_factor` on their different paths
- load_generator_scaling_factor** the factor for the reward or reduction cost factor in the objective function when `load_generator_type` is set to `REWARD` or `REDUCTION`. A higher value will result in larger detours for the passengers.
- load_generator_sp_distribution_factor** the parameter for the logit model used to distribute the passenger when `load_generator_number_of_shortest_paths` is bigger than 1.
- load_generator_type** the different ptn load generator types.
- SP** use the travel time shortest paths for the passengers, depending on the travel time approximation used.
 - REDUCTION** adds a penalty depending on the usage of the edge in the PTN (high penalty for low usage)
 - REWARD** reward an edge more, if less passengers are needed to fill the next vehicle on the edge
- For a more detailed description of the different models, see [8].
- load_generator_use_cg** whether to use a change and go network for routing. This includes knowledge of the line pool, allowing to consider transfers. The line pool needs to be present!

- load_generator_upper_frequency_factor** the factor to multiply the lower frequency bound to obtain the new upper frequency bound. The result is rounded up. Whether this or a fixed bound is used depends on **load_generator_fix_upper_frequency**.
- load_generator_building_cost** sets building cost as the objective (obj) or a constraint (cons). If building cost should not be considered in the model, set this to none.
- load_generator_travel_time** sets total travel time as the objective (obj) or a constraint (cons). If total travel time should not be considered in the model, set this to none.
- load_generator_detour_factor** sets maximum detour factor as the objective (obj) or a constraint (cons). If maximum detour factor should not be considered in the model, set this to none.
- load_generator_max_building_cost** upper bound for building cost used if **load_generator_building_cost** is set to cons. Note that this value is converted according to **gen_conversion_length**.
- load_generator_max_travel_time** upper bound for total travel time used if **load_generator_travel_time** is set to cons. Note that this value is converted according to **gen_conversion_length**.
- load_generator_max_detour** upper bound for maximum detour factor used if **load_generator_detour_factor** is set to cons.
- load_generator_gen_cuts** when this is set to true, so-called valid inequalities are added to the spanner MIP. These have no effect on the solution, but tend to greatly improve computational performance.
- load_generator_remove_unused_edges** when this is set to true, after solving the model, the unused edges are removed from **default_edges_file** (basis/Edge.giv).
- load_generator_solver** determine the solver to be used.
- load_generator_threads** determine the maximal number of threads to use for the solver (-1=use default value, i.e., no restriction). Note that this will only be used for a possible solver integration of the chosen model, not for the rest of the algorithm.
- load_generator_timelimit** the time limit for the solver in seconds (-1=use default value).
- load_generator_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10 % (-1=use default value).
- load_generator_write_lp_file** whether to write the lp file of the model to solve.

7.8 Load Visualization

- filename_loads_graph_file** filename under which the load visualization is saved.
- loads_draw_conversion_factor** scaling factor for the load visualization.
- loads_graph_lower_bound** percentage of the maximum load. Lower loads will not be displayed in the load visualization.
- loads_graph_max_edge_width** maximum edge width used in the load visualization which uses edge width as a medium to display the traffic load.
- loads_graph_upper_bound** percentage of the maximum load. Higher loads will not be displayed in the load.
- loads_graph_use_edge_color** whether to use the edge color to represent the traffic loads. Otherwise edge width will be used.

7.9 Periodic EAN

- CK** **ean_algorithm_shortest_paths** the algorithm to use to compute the shortest paths in the ean. Choices are **CV** JOHNSON, **CV** FLOYD, **CV** FIBONACCI_HEAP and **CV** TREE_MAP_QUEUE.
- CK** **ean_change_penalty** the change penalty for routing, i.e., a penalty for each transfer a passenger needs to take during their journey. Given in time units.
- CK** **ean_construction_skip_passenger_distribution** whether to skip the initial passenger distribution during ean construction.
- CK** **ean_construction_target_model_frequency** whether to include the frequency of lines only as attributes in the ean (**CV** FREQUENCY_AS_ATTRIBUTE) or include multiple frequency repetitions, connected by synchronization activities (**CV** FREQUENCY_AS_MULTIPLICITY). Note that **CV** FREQUENCY_AS_ATTRIBUTE can not be handled by all timetabling algorithms.
- CK** **ean_construction_target_model_headway** how to model headways in the ean. The following options are available:
- CV** **NO_HEADWAYS** create new headway activities
 - CV** **SIMPLE** creates a headway between every two departures from the same station using the same link
 - CV** **PRODUCT_OF_FREQUENCIES** When using **CV** FREQUENCY_AS_ATTRIBUTE as a **CK** ean_construction_target_model_frequency, this will create all the necessary headway activities between the corresponding departure events, i.e., will include multiple headway activities for lines with frequency > 1. Note that this is not necessary when **CK** ean_construction_target_model_frequency is set to **CV** FREQUENCY_AS_MULTIPLICITY, in this case this model is the same as **CV** SIMPLE.
 - CV** **LCM_OF_FREQUENCIES** This will create as many headway activities as the least common multiple of the two corresponding line frequencies between all departures from the same station using the same link.
 - CV** **LCM_REPRESENTATION** For headway creation, this behaves the same as **CV** SIMPLE but some timetabling models will respect these headways the same as **CV** LCM_OF_FREQUENCIES later on.
- CK** **ean_default_maximal_change_time** the default maximal change time at a station
- CK** **ean_default_maximal_waiting_time** the default maximal waiting time at a station
- CK** **ean_default_minimal_change_time** the default minimal change time at a station
- CK** **ean_default_minimal_waiting_time** the default minimal waiting time at a station
- CK** **ean_discard_unused_change_activities** when set to **CV** true, this will remove all change activities from the ean that do not have a positive weight after the initial passenger distribution
- CK** **ean_dump_initial_duration_assumption** when set to **CV** true, this will output the initial duration assumption of every activity, i.e., the computed duration of every activity in the initial passenger distribution. **CK** filename_initial_duration_assumption (**Fi** timetabling/Initial-duration-assumption-periodic.giv) will be used for output.
- CK** **ean_individual_station_limits** when set to **CV** true, individual station limits for change and waiting time will be used. For information on how to give these limits, see the documentation for **CK** filename_station_limit_file (**Fi** basis/Station-Limits.giv).

ean_initial_duration_assumption_model How to compute the initial duration assumption. The following options are available:

AUTOMATIC fully automated initial durations, based on **ean_model_weight_change**, **ean_model_weight_drive** and **ean_model_weight_wait**.

SEMI_AUTOMATIC the initial duration for individual activities can be given. For information on how to give these durations, see the documentation for **filename_initial_duration_assumption** (**timetabling/Initial-duration-assumption-periodic.giv**).

ean_model_weight_change determines how to estimate the transfer time between two lines without a given timetable. For a transfer between the lines l_1 and l_2 , let f_1 and f_2 be the respective frequencies and T the **period_length**. The following options are available:

FORMULA_1 $\frac{T}{f_1} + \frac{T}{f_2} +$ **ean_change_penalty**.

FORMULA_2 $\frac{T}{2 \cdot f_1 \cdot f_2} +$ **ean_change_penalty**.

FORMULA_3 $\frac{T}{2 \cdot f_2} +$ **ean_change_penalty**.

MINIMAL_CHANGING_TIME **ean_default_minimal_change_time** + **ean_change_penalty**.

ean_model_weight_drive determines how to estimate the drive time on an infrastructure edge without a given timetable. The following options are available:

AVERAGE_DRIVING_TIME using the average between minimal and maximal travel time of the infrastructure edge.

EDGE_LENGTH using the edge length of the infrastructure edge.

MINIMAL_DRIVING_TIME using the minimal travel time of the infrastructure edge.

MAXIMAL_DRIVING_TIME using the maximal travel time of the infrastructure edge.

ean_model_weight_wait determines how to estimate the waiting time when traversing a stop in a line without a given timetable. The following options are available:

AVERAGE_WAITING_TIME using the average of **ean_default_minimal_waiting_time** and **ean_default_maximal_waiting_time**.

MAXIMAL_WAITING_TIME using **ean_default_maximal_waiting_time**.

MINIMAL_WAITING_TIME using **ean_default_minimal_waiting_time**.

ZERO_COST assume the waiting time to be 0.

ean_random_shortest_paths

ean_use_walking whether to allow walking transfers in the EAN

7.10 Debug

debug_paths_in_ptn when set to **true**, some routing methods will output the found ptn paths to **default_debug_od_link_paths_file** (**Debug/ODLinkPaths.dbg**)

debug_paths_in_ean when set to **true**, some routing methods will output the found ean paths to **default_debug_od_activity_paths_file** (**Debug/ODActivityPaths.dbg**).

7.11 Timetabling

- tim_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10% (-1=use default value).
- tim_model** the timetabling model to use. For an overview of all models, see Section 3.6
- tim_pesp_ip_solution_limit** limit the number of feasible solutions found. Only implemented in Gurobi. Set to 0 to deactivate.
- tim_pesp_ip_best_bound_stop** a best bound stop criterion, only implemented for Gurobi. For details, see Gurobi documentation. Set to 0 to deactivate.
- tim_pesp_ip_mip_focus** set the MIPFocus, only implemented for Gurobi. For details, see Gurobi documentation. Set to 0 to deactivate.
- tim_solver** the solver to use for timetabling. Which solvers are implemented depends on the chose **tim_model**, see the corresponding documentation.
- tim_threads** determine the maximal number of threads to use for the solver (-1=use default value, i.e., no restriction). Note that this will only be used for a possible solver integration of the chosen model, not for the rest of the algorithm.
- tim_timelimit** the time limit to use for the solver in seconds (-1 = use default value).
- tim_use_old_solution** whether to use the current solution as a starting solution, only implemented for Gurobi and the pesp ip.
- tim_write_lp_file** whether to write the lp file of the model to solve

7.12 Tariff Planning

- taf_model** either **flat**, **beeline_distance**, **network_distance** or **zone**. Determining the model used to calculate a new tariff.
- taf_objective** either **sum_absolute_deviation** or **max_absolute_deviation**. Determining whether the sum or maximum of absolute price deviations is minimized.
- taf_weights_objective** either **od**, **unit** or **reference-inverse**. Determining whether the price deviations in the objective are weighted by the OD data, by unit weights or by the inverse of the given reference prices.
- taf_zone_counting** either **single** or **multiple**. If **single**, then each zone is only counted once when determining the number of traversed zones of a path. If **multiple**, a zone is counted each time that it is entered.
- taf_zone_n_zones** positive integer number, specifies the maximum number of zones when calculating a new zone tariff.
- taf_zone_enforce_all_zones** boolean, determines whether exactly **taf_zone_n_zones**-many zones (**true**) or at most that many zones (**false**) must be determined.
- taf_zone_connected** boolean, specifies whether the subgraph of a zone, induced by the nodes assigned to the zone, needs to be connected (in case of a directed graph it is weakly connected).
- taf_zone_enforce_no_elongation** boolean, determining whether the no-elongation property must be satisfied. This property ensures, that it is never cheaper for passengers to buy tickets for more zones than they actually want to travel through. Let p_k be the price of a path that uses k zones. The no-elongation property is satisfied if it holds that

$$p_k \leq p_{k+1} \quad \text{for all } k \in \{1, \dots, (\text{CK } \text{taf_zone_n_zones}) - 1\}.$$

taf_zone_enforce_no_stopover boolean, determining whether the no-stopover property must be satisfied. This property ensures that it is never cheaper for passengers to buy two separate tickets for one journey and combine them instead of buying one ticket for the whole journey. Let p_k be the price of a path that uses k zones. The no-stopover property in the case of single counting holds if

$$p_k \leq p_i + p_j \quad \text{for all } k \in \mathbb{N}_{\geq 1}, i, j \in \{1, \dots, k\} \text{ with } i + j \geq k + 1.$$

In the case of multiple counting the property holds if

$$p_k \leq p_i + p_j \quad \text{for all } k \in \mathbb{N}_{\geq 1}, i, j \in \{1, \dots, k\} \text{ with } i + j = k + 1.$$

taf_zone_symmetry_breaking determines which symmetry breaking model (see below) should be used. Possible values are A, B and NONE.

taf_routing_generation either fastest-paths, read-all or read-partial-fill. Determines which routing should be used, see Section 3.7.1.

taf_zone_only_zones boolean, specifies whether only zones based on given prices are computed.

taf_zone_only_prices boolean, specifies whether only prices based on given zones are computed.

taf_draw_zones boolean, specifies whether a PTN with nodes allocated to their zones should be drawn. By default false.

taf_draw_heatmap boolean, specifies whether a heatmap should be drawn.

taf_heatmap_mode either old, new or compare. Specifies which prices or price differences should be shown in a heatmap.

taf_heatmap_log_scale boolean, specifies whether or not the heatmap should use a logarithmic scale. By default false.

taf_heatmap_use_annotations boolean, specifies whether or not the heatmap should be annotated. By default false.

taf_evaluate_old_prices points towards a first price matrix for the evaluation and for the heatmap. By default it is the reference price matrix basis/Reference-Price-Matrix.giv.

taf_evaluate_new_prices points towards a second price matrix for the evaluation and the heatmap. By default it is the tariff price matrix tariff/Price-Matrix.taf.

taf_solver determine the solver to be used. Note that currently only Gurobi is supported.

taf_threads determine the maximum number of threads to use for the solver (-1=use default value, i.e., no restriction). Note that this will only be used for a possible solver integration of the chosen model, not for the rest of the algorithm.

taf_timelimit the time limit for the solver in seconds (-1=use default value).

taf_write_lp_file whether to write the lp file of the model to solve.

taf_mip_gap sets the MIP optimization gap for the solver. The solver will terminate with an optimal solution if the gap between lower and upper objective bound is less than this value times the absolute value of the incumbent objective value.

7.13 Vehicle Scheduling

- vs_depot_index** the stop index of the depot. Set to -1 to disable to consideration of a depot.
- vs_eval_cost_factor_empty_trips_duration** the weight factor for the duration of empty trips in the cost function for a vehicle schedule
- vs_eval_cost_factor_empty_trips_length** the weight factor for the length of empty trips in the cost function for a vehicle schedule
- vs_eval_cost_factor_full_trips_duration** the weight factor for the duration of services in the cost function for a vehicle schedule
- vs_eval_cost_factor_full_trips_length** the weight factor for the length of services in the cost function for a vehicle schedule
- vs_maximum_buffer_time** the maximal buffer time between the service of two trips
- vs_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10% (-1=use default value).
- vs_model**] the vehicle scheduling model to use. For an overview of all models, see Section 3.8
- vs_solver** the solver to use for vehicle scheduling. Which solvers are implemented depends on the chose **vs_model**, see the corresponding documentation.
- vs_timelimit** the time limit to use for the solver in seconds (-1 = use default value).
- vs_threads** determine the maximal number of threads to use for the solver (-1=use default value, i.e., no restriction). Note that this will only be used for a possible solver integration of the chosen model, not for the rest of the algorithm.
- vs_turn_over_time** the minimal time between two services, given in time units.
- vs_vehicle_costs** the costs of a vehicle
- vs_write_lp_file** whether to write the lp file of the model to solve

7.14 Delay Management

- DM_best_of_all_write_objectives** whether to write all objectives to a file, when **DM_method** **best-of-all** is used
- DM_debug** enable debug output
- DM_earliest_time** the start of the rollout period
- DM_enable_consistency_checks** enable consistency checks for the input data, i.e., 28800 is 08:00.
- DM_eval_extended** enable the extended evaluation
- DM_latest_time** the end of the rollout period, given in seconds after midnight, i.e., 28800 is 08:00.
- DM_method** the delay management model to use. For an overview of all models, see Section 3.9.
- DM_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10% (-1=use default value).
- DM_opt_method_for_heuristic** the optimization method to use for the heuristics.

- DM_solver** the solver to use for vehicle scheduling. Which solvers are implemented depends on the chose **DM_model**, see the corresponding documentation.
- DM_threads** determine the maximal number of threads to use for the solver (-1=use default value, i.e., no restriction). Note that this will only be used for a possible solver integration of the chosen model, not for the rest of the algorithm.
- DM_timelimit** the time limit to use for the solver in seconds (-1 = use default value).
- DM_write_lp_file** whether to write the lp file of the model to solve
- DM_verbose** enable verbose output

7.15 Dataset Generation

- dg_model** The model that should be used to create the new dataset. For a detailed description of the available algorithms see Section 4.1.

The following parameters are only valid for the Parametrized-City-Model:

- dg_param_city_number_subcenters** Number of subcenters surrounding the CBD. The PTN has $2n + 1$ nodes.
- dg_param_city_alpha** Trips proportion from periphery that go to the CBD.
- dg_param_city_beta** Trips proportion from periphery to own subcenter.
- dg_param_city_eta** Portion of displacement of the CBD from the center of the city in an axis CBD-subcenter.
- dg_param_city_Y** Total number of trips generated.
- dg_param_city_L** Distance from any subcenter to the geometrical center of the city.
- dg_param_city_g** Distance periphery-subcenter / Distance subcenter-CBD.
- dg_param_city_a** Trips proportion that depart from the periphery.
- dg_ring_number_of_rings** Number of concentric rings that are generated.
- dg_ring_nodes_per_ring** Number of nodes that each ring consists of
- dg_ring_length_1** If this boolean parameter is set to true, the lengths of all edges are equal to 1.
- dg_ring_radius** specifies the radius of the inner ring, i.e. the lengths of the edges from the center to the nodes of the inner ring.
- dg_ring_demand_type** specifies the method for the creation of the OD data.

7.16 Integrated Models

7.16.1 General

- int_solver** the solver to use. Which solvers are implemented depends on the chosen model, see the corresponding documentation.
- int_threads** determine the maximal number of threads to use for the solver (-1=use default value, i.e., no restriction). Note that this will only be used for a possible solver integration of the chosen model, not for the rest of the algorithm.

7.16.2 LinTimPass

- lin_tim_pass_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10% (-1=use default value).
- lin_tim_pass_timelimit** the time limit to use for the solver in seconds (-1 = use default value).
- lin_tim_pass_write_lp_file** whether to write the lp file of the model to solve.

7.16.3 LinTimPassVeh

- lin_tim_pass_veh_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10% (-1=use default value).
- lin_tim_pass_veh_timelimit** the time limit to use for the solver in seconds (-1 = use default value).
- lin_tim_pass_veh_write_lp_file** whether to write the lp file of the model to solve.

7.16.4 TimPass

- tim_pass_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10% (-1=use default value).
- tim_pass_timelimit** the time limit to use for the solver in seconds (-1 = use default value).
- tim_pass_write_lp_file** whether to write the lp file of the model to solve.

7.16.5 TimVeh

- tim_veh_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10% (-1=use default value).
- tim_veh_timelimit** the time limit to use for the solver in seconds (-1 = use default value).
- tim_veh_write_lp_file** whether to write the lp file of the model to solve.

7.16.6 TimVehToLin

- tim_veh_to_lin_mip_gap** the mip optimization gap for the solver, 0.1 equals a gap of 10% (-1=use default value).
- tim_veh_to_lin_timelimit** the time limit to use for the solver in seconds (-1 = use default value).
- tim_veh_to_lin_write_lp_file** whether to write the lp file of the model to solve.

7.17 TimPassLib

- timpasslib_import_timetable** whether timetable is imported.
- timpasslib_export_timetable** whether timetable is exported.

Chapter 8

In- and Output Data

This section will describe all files and their contents that are in- or outputs of the LINTIM algorithms.

8.1 Config

Config is the short form for *configuration* and an important tool in LINTIM. We will now have a look at the general structure of the LINTIM config files.

The LINTIM config is contained in several CSV files that have the syntax

```
config_key; config_value
```

It organizes those values that are parameters to the calculation. Typical examples are the period length, the vehicle capacity (if there is only one), which algorithm to use for a specific computation step, e.g. for timetabling and filenames as well and could thus look like

```
period_length; 60  
gen_passengers_per_vehicle; 100  
tim_model; MATCH
```

Besides key-value pairs the configuration may also include other config files with either the `[CK] include` or `[CK] include_if_exists` statement. Former states that the file must exist or else an exception is thrown, in latter case, if the file does not exist, it will not be included. This inclusion is recursive, i.e. files included in already included files are included as well.

If a certain config key occurs twice, the latter value overwrites the former, e.g.

```
period_length; 60  
period_length; 120
```

sets the `[CK] period_length` to 120. As a consequence, all values that belong to keys in an included file overwrite those defined before.

All keys demanded by programs are expected to exist, i.e., there are no in-program default values. Programs accessing config are expected to exit with an error message in case a key does not exist.

The meaning of the parameters is explained in the corresponding sections of this documentation.

Config has the following file hierarchy

`[Fi] /datasets/Global-Config.cnf` offers a default value for all config parameters that are not network specific, like `[CK] ptn_name` or `[CK] period_length`.

`[Fi] basis/Config.cnf` contains all the values specific to the dataset. Together with the global config this offers a value for all parameters. It includes the global config at the beginning, i.e., every parameter that was already defined in the global config will be overwritten. It roughly looks like


```

include; "../..//Global-Config.cnf"
ptn_name; "DATASET"
...
include_if_exists; "State-Config.cnf"
include_if_exists; "Private-Config.cnf"
include_if_exists; "After-Config.cnf"

```

CK `filename_state_config` (**Fi** `basis/State-Config.cnf`) is intended to allow programs to not only generate networks, but also to save and modify state information about them, e.g. whether the event activity network is modeled with `frequency_as_attribute` or `frequency_as_multiplicity` which is once set on construction and may be modified by a `PERIODIC ROLLOUT`. The network specific state is not part of the version control system, although there are state defaults in the global config.

Fi `basis/Private-Config.cnf` is used for user specific settings, e.g. for choosing a specific algorithm for solving or manipulating its parameters and is not part of the version control system. Note that if a value is defined in the config or state config as well as in the private config, the one given in the private config is used.

Fi `basis/After-Config.cnf` can be used for automation and is intended to be *thrown away* upon usage, unlike all other configurations. A script that automatically evaluates a wide range of configurations thus may overwrite the after config in every step. Make sure that at the end of the script, the after config is deleted again or else it still influences manual runs as it overwrites all other configs.

8.2 Statistic

The statistic file **CK** `default_statistic_file` (**Fi** `statistic/statistic.sta`) contains the outcome of the evaluation routines described in 5. The content is formatted as follows

```

statistic_key; statistik_value

```

where the statistic key described what is evaluated and the statistic value gives the corresponding value. Statistic files are intended to be modified, i.e., new entries are added but old entries are not deleted, although the statistic file itself may be deleted any time. Make sure that the entries are up to date, e.g. **R** `make tim-timetable-evaluate` is run after calculating a new timetable and before evaluating the statistic.

8.3 Basis

Files in the folder **Fo** `basis` describe the structure of the Public Transportation Network, the demand and the line pool with its corresponding costs.

8.3.1 Additional load

The file **CK** `filename_additional_load_file` (**Fi** `basis/Additional-Load.giv`) contains additional load on single PTN links. When **CK** `load_generator_add_additional_load` is set to **CV** `true`, these loads will be added to the corresponding links during load generation. For an undirected network, a link may be given in both directions, allowing for different additional load values for the different directions. Unmentioned links will be assumed to have no additional load. The columns of the csv file correspond to:

edge-id id of the PTN edge

left-stop-id the id of the left stop, i.e., the origin of the edge

right-stop-id the id of the right stop, i.e., the destination of the edge

additional-load the value of the additional load

8.3.2 Change station

The file `[CK] filename_change_station_file ([Fi] basis/Change-Stations.giv)` contains a list of change stations, i.e., a list of stops where passengers can transfer. The columns of the csv file correspond to:

stop-id id of the stop

8.3.3 Demand

The file `[CK] default_demand_file ([Fi] basis/Demand.giv)` contains the demand at specified locations. The columns of the csv file correspond to:

demand-id id of the demand point

short-name short name of the demand point

long-name log name of the demand point

x-coordinate x-coordinate of the demand point

y-coordinate y-coordinate of the demand point

demand demand at the demand point

Note: the distance between two demand points can be transformed to kilometers by multiplying with `[CK] gen_conversion_coordinates`.

8.3.4 Demand geo

The file `[CK] default_demand_coordinates_file ([Fi] basis/Demand.giv.geo)` gives the geographical coordinates (latitude and longitude) of the demand points. The columns of the csv file correspond to:

demand-id id of the demand point

latitude latitude of the demand point

longitude longitude of the demand point

8.3.5 Edge

The file `[CK] default_edges_file ([Fi] basis/Edge.giv)` contains information about the edges in the PTN. The columns of the csv file correspond to:

edge-id id of the edge

left-stop-id id of the left stop (source node in directed case)

right-stop-id id of the right stop (target node in directed case)

length length of the edge

lower-bound minimum time to traverse the edge in minutes

upper-bound maximum time to traverse the edge in minutes

Note: whether the edges are directed or undirected is defined by `[CK] ptn_is_undirected`.

Note: the length of an edge can be transformed to kilometers by multiplying with `[CK] gen_conversion_length`.

8.3.6 Edge forbidden

The file `filename_forbidden_links_file` (`basis/Edge-forbidden.giv`) contains information about the edges in the PTN that are forbidden, i.e., that may not be used by the public transport mode that is being planned. These edges should be a subset of the edges in `default_edges_file` (`basis/Edge.giv`). The columns of the csv file correspond to:

edge-id id of the edge

left-stop-id id of the left stop (source node in directed case)

right-stop-id id of the right stop (target node in directed case)

length length of the edge

lower-bound minimum time to traverse the edge in minutes

upper-bound maximum time to traverse the edge in minutes

Note: whether the edges are directed or undirected is defined by `ptn_is_undirected`.

Note: the length of an edge can be transformed to kilometers by multiplying with `gen_conversion_length`.

8.3.7 Edge infrastructure

The file `filename_infrastructure_edge_file` (`basis/Edge-Infrastructure.giv`) contains information about the infrastructure edges, i.e., edges that connect infrastructure nodes. The columns of the csv file correspond to:

edge-id id of the edge

left-node-id id of the left stop (source node in directed case)

right-node-id id of the right stop (target node in directed case)

length length of the edge

lower-bound minimum time to traverse the edge in minutes

upper-bound maximum time to traverse the edge in minutes

Note: whether the edges are directed or undirected is defined by `ptn_is_undirected`.

Note: the length of an edge can be transformed to kilometers by multiplying with `gen_conversion_length`.

8.3.8 Edge infrastructure forbidden

The file `filename_forbidden_infrastructure_edges_file` (`basis/Edge-Infrastructure-forbidden.giv`) contains information about the infrastructure edges that are forbidden, i.e., that may not be used by the public transport mode that is being planned. These edges should be a subset of the edges in `filename_infrastructure_edge_file` (`basis/Edge-Infrastructure.giv`). The columns of the csv file correspond to:

edge-id id of the edge

left-node-id id of the left node (source node in directed case)

right-node-id id of the right node (target node in directed case)

length length of the edge

lower-bound minimum time to traverse the edge in minutes

upper-bound maximum time to traverse the edge in minutes

Note: whether the edges are directed or undirected is defined by `[CK] ptn_is_undirected`.

Note: the length of an edge can be transformed to kilometers by multiplying with `[CK] gen_conversion_length`.

8.3.9 Edge walking

The file `[CK] filename_walking_edge_file ([Fi] basis/Edge-Walking.giv)` contains information about the possible walking edges, i.e., connections between infrastructure nodes that can directly be used for walking by the passengers. The columns of the csv file correspond to:

edge-id id of the edge

left-node-id id of the left node (source node in directed case)

right-node-id id of the right node (target node in directed case)

length length of the edge, given in seconds

Note: whether the edges are directed or undirected is defined by `[CK] sl_walking_is_directed`.

Note: when read by `LINTIM`, `[CK] sl_max_walking_time` will be respected, i.e., only edges with a length smaller than the given value will be read. A value of `[CV] -1` will disable this and allow all edges will be read.

Note: it is possible to preprocess the walking edges by using

`[R] make ptn-preprocess-walking`.

With this, walking edges will be filtered by `[CK] sl_max_walking_amount`, `[CK] sl_max_walking_ratio` (both per node with outgoing demand) and `[CK] sl_max_walking_time`, possibly reducing the size of the walking graph.

8.3.10 Existing stop

The file `[CK] default_existing_stop_file ([Fi] basis/Existing-Stop.giv)` contains information about already existing stops in the PTN. The columns of the csv file correspond to:

stop-id id of the stop

short-name short name of the stop

long-name log name of the stop

x-coordinate x-coordinate of the stop

y-coordinate y-coordinate of the stop

Note: the distance between two stops can be transformed to kilometers by multiplying with `[CK] gen_conversion_coordinates`.

8.3.11 Existing stop geo

The file `[CK] default_existing_stop_coordinates_file ([Fi] basis/Existing-Stop.giv.geo)` gives the geographical coordinates (latitude and longitude) of the already existing stops. The columns of the csv file correspond to:

stop-id id of the stop

latitude latitude of the stop

longitude longitude of the stop

8.3.12 Existing edge

The file `[CK] default_existing_edge_file ([Fi] basis/Existing-Edge.giv)` contains information about already existing edges in the PTN. The columns of the csv file correspond to:

edge-id id of the edge

left-stop-id id of the left stop (source node in directed case)

right-stop-id id of the right stop (target node in directed case)

length length of the edge

lower-bound minimum time to traverse the edge in minutes

upper-bound maximum time to traverse the edge in minutes

Note: whether the edges are directed or undirected is defined by `[CK] ptn_is_undirected`.

Note: the length of an edge can be transformed to kilometers by multiplying with `[CK] gen_conversion_length`.

8.3.13 Headway

The file `[CK] default_headways_file ([Fi] basis/Headway.giv)` contains information about the headway needed for the edges in the PTN. The columns of the csv file correspond to:

edge-id id of the edge

headway headway on the edge, i.e., the minimum time between two consecutive vehicles on this edge in minutes

8.3.14 Load

The file `[CK] default_loads_file ([Fi] basis/Load.giv)` contains information about the load and frequency constraints of the edges in the PTN. The columns of the csv file correspond to:

edge-id id of the edge

load load on the edge

lower-frequency minimal frequency all lines in the line concept have to add up to the edge

upper-frequency maximal frequency all lines in the line concept are allowed to add up to for the edge

8.3.15 Node

The file `[CK] filename_node_file ([Fi] basis/Node.giv)` contains information about infrastructure nodes. Infrastructure nodes are the smallest unit of nodes in LINTIM, they may e.g. represent crossings or (potential) stops. The columns of the csv file correspond to:

node-id the id of the node

name the name of the node

x-coordinate the x coordinate of the node

y-coordinate the y coordinate of the node

stop-possible? whether its possible for this node to be a stop

Note: x- and y-coordinate are assumed to be planar coordinates, i.e., will be directly used to compute the euclidean distance between stops. The distance between two stops can be transformed to kilometers by multiplying with `[CK] gen_conversion_coordinates`.

8.3.16 OD

The file `default_od_file` (`basis/OD.giv`) contains information about the passenger demand between all pairs of stops in the PTN. The columns of the csv file correspond to:

left-stop-id id of the stop the passengers start at

right-stop-id id of the stop the passengers travel to

customers number of passengers traveling

8.3.17 OD node

The file `filename_od_nodes_file` (`basis/OD-Node.giv`) contains information about the passenger demand between pairs of nodes in the infrastructure network. The columns of the csv file correspond to:

left-node-id id of the node the passengers start at

right-node-id id of the node the passengers travel to

customers number of passengers traveling

8.3.18 Pool

The file `default_pool_file` (`basis/Pool.giv`) contains information about the line pool. The columns of the csv file correspond to:

line-id id of the line

edge-order where the edge is in the line

edge-id id of the edge

8.3.19 Pool cost

The file `default_pool_cost_file` (`basis/Pool-Cost.giv`) contains information about the cost and length of lines in the line pool. The columns of the csv file correspond to:

line-id id of the line

length length of the line

cost cost of the line

Note: the length of a line can be transformed to kilometers by multiplying with `gen_conversion_length`.

8.3.20 Reference Price Matrix

The file `filename_tariff_reference_price_matrix_file` (`basis/Reference-Price-Matrix.giv`) contains the given reference prices for each OD pair. The columns of the csv file correspond to:

origin-id id of the origin of the OD pair,

destination-id id of the destination of the OD pair,

price reference price when travelling from the origin to the destination.

8.3.21 Restricted turns

The file `filename_turn_restrictions` (`basis/Restricted-Turns.giv`) contains information about restricted turns, i.e., pairs of link ids of the PTN that are not allowed to be traversed by a line directly after each other. The columns of the csv file correspond to:

first-edge-id the first edge id

second-edge-id the second edge id

Note: whether the information will be interpreted as directed is dependent on `ptn_is_undirected`.

8.3.22 Restricted turns infrastructure

The file `filename_turn_restrictions_infrastructure` (`basis/Restricted-Turns-Infrastructure.giv`) contains information about restricted turns, i.e., pairs of edge ids in the infrastructure network that are not allowed to be traversed by a line directly after each other. The columns of the csv file correspond to:

first-edge-id the first edge id

second-edge-id the second edge id

Note: whether the information will be interpreted as directed is dependent on `ptn_is_undirected`.

8.3.23 Ridepool

The file `filename_rpool_file` (`basis/Ridepool.giv`) contains a ridepool, i.e. areas in the PTN. The columns of the csv file correspond to:

area-id id of the area

edge-id id of an edge contained in the area

demand-factor edge-specific constants specifying the fraction of time where one vehicle in the area is used on this edge.

8.3.24 Routings

The files `filename_routing_ptn_input` (`basis/Routing-ptn.giv`) contains a routing in the PTN, i.e. for each node pair at most one path is specified as a list of nodes. The columns of the csv file correspond to:

origin-id id of the first node of the path,

destination-id id of the last node of the path,

node-ids path specified by the sequence of the stop-ids.

The parameter `filename_routing_ptn_output` (`basis/Routing-ptn.giv`) specifies the file, where the computed routing is written to.

8.3.25 Station limits

The file `filename_station_limit_file (basis/Station-Limits.giv)` contains information about individual station limits on wait or change times. The columns of the csv file correspond to:

stop-id the id of the stop

min-wait-time the minimal waiting time.

max-wait-time the maximal waiting time.

min-change-time the minimal change time.

max-change-time the maximal change time.

Note: every individual limit may be set to -1 if there is none. Then the corresponding default parameters will be used. The same holds for stops not present in this file.

8.3.26 Stop

The file `default_stops_file (basis/Stop.giv)` contains information about the stops in the PTN. The columns of the csv file correspond to:

stop-id id of the stop

short-name short name of the stop

long-name log name of the stop

x-coordinate x-coordinate of the stop

y-coordinate y-coordinate of the stop

Note: x- and y-coordinate are assumed to be planar coordinates, i.e., will be directly used to compute the euclidean distance between stops. The distance between two stops can be transformed to kilometers by multiplying with `gen_conversion_coordinates`.

8.3.27 Stop geo

The file `default_stops_coordinates_file (basis/Stop.giv.geo)` gives the geographical coordinates (latitude and longitude) of the stops. The columns of the csv file correspond to:

stop-id id of the stop

latitude latitude of the stop

longitude longitude of the stop

8.3.28 Terminals

The file `filename_terminals_file (basis/Terminals.giv)` gives the stop ids of terminals, i.e., stops where lines are allowed to terminate. The columns of the csv file correspond to:

stop-id id of the stop

Note: the stop ids should be a subset of the ptn stops, i.e., of `default_stops_file (basis/Stop.giv)`.

8.4 Line Planning

The folder `line-planning` contains information about the line concept.

8.4.1 Line concept

The file `default_lines_file` (line-planning/Line-Concept.lin) contains information about the line concept. The columns of the csv file correspond to:

line-id id of the line

edge-order where the edge is in the line

edge-id id of the edge

frequency frequency of the line. If this is zero, the line is not used in the line concept.

8.4.2 Fixed lines

The file `filename_lc_fixed_lines` (line-planning/Fixed-Lines.lin) contains information about the fixed lines that should be in the line concept. It can not be read/respected by all line planning methods, so see Section 3.3 for more information. The columns of the csv file correspond to:

line-id id of the line

edge-order where the edge is in the line

edge-id id of the edge

frequency frequency of the line. If this is zero, the line is not used in the line concept.

8.4.3 Line capacities

The file `filename_lc_fixed_line_capacities` (line-planning/Line-Capacities.lin) contains information about the capacities of the fixed lines that should be in the line concept. It can not be read/respected by all line planning methods, so see Section 3.3 for more information. The columns of the csv file correspond to:

line-id id of the line

capacity the capacity of the line

8.4.4 Rideconcept

The file `filename_rc_file` (line-planning/Ride-Concept.lin) contains a ride concept, i.e. a set of ridepooling areas with assigned number of vehicles for each area. The columns of the csv file correspond to:

area-id id of the area

edge-id id of an edge contained in the area

demand-factor edge-specific constants specifying the fraction of time where one vehicle in the area is used on this edge.

number-of-vehicles Number of vehicles operation in the area

8.5 Timetabling

The folder `timetabling` contains information about the periodic event-activity-network and the timetable.

8.5.1 Activities periodic

The file default_activities_periodic_file (timetabling/Activities-periodic.giv) contains information about activities in the periodic EAN. The columns of the csv file correspond to:

activity-id id of the activity

type type of the activity, can be **drive** for drive activities, **wait** for wait activities, **change** for transfers of passengers, **sync** for synchronization activities between different servings of a line with frequency greater than one or **turnaround** for turnaround activities, i.e., activities of vehicles serving one line after another

tail-event-id id of source event, i.e., the start of the activity

head-event-id id of target event, i.e., the end of the activity

lower-bound the minimal time for this activity, i.e., the minimal time duration needed between the corresponding source and target event to be feasible

upper-bound the maximal time for this activity, i.e., the maximal time duration allowed between the corresponding source and target event to be feasible

passengers the number of passengers using this activity

8.5.2 Events periodic

The file default_events_periodic_file (timetabling/Events-periodic.giv) contains information about events in the periodic EAN. The columns of the csv file correspond to:

event-id id of the event

type type of the event, can be **departure** for events which are departures of a line at a stop or **arrival** for events which are arrivals of a line at a stop

stop-id id of the corresponding stop

line-id id of the corresponding line

passengers number of passengers boarding/alighting at the event

line-direction direction of the line, > for forward direction (i.e., the direction given in the file Pool.giv) or < for the backward direction

line-freq-repetition repetition of the line, i.e., how often the line has already been used in the planning period

8.5.3 Fixed times

The file filename_tim_fixed_times (timetabling/Fixed-timetable-periodic.tim) gives restrictions on the allowed times for single events. Not all events need to be included in this file, only the ones with additional restrictions.

event-id the periodic event id

lower-bound the lower time bound on the event

upper-bound the upper time bound on the event

8.5.4 Initial duration assumptions

The file filename_initial_duration_assumption (timetabling/Initial-duration-assumption-periodic.giv) may contain a duration for each activity used in the initial passenger distribution of the ean creation. The columns of the csv file correspond to:

activity-id id of the activity

duration the duration to use for the passenger distribution

Note that ean_initial_duration_assumption_model needs to be set to SEMI_AUTOMATIC for this file to be read. Not all activities need to be present in the file, the duration of activities not present will be computed normally.

8.5.5 Timetable periodic

The file default_timetable_periodic_file (timetabling/Timetable-periodic.tim) contains a time for each event in the periodic EAN. The columns of the csv file correspond to:

event-id id of the event

time the periodic time of the event

8.5.6 Timetable for VISUM

The file default_timetable_visum_file (timetabling/Timetable-visum-nodes.tim) is an intermediate format for reading a LINTIM timetable into VISUM. For more information, see 4.12. The columns of the csv file correspond to:

line-id the line id

line-code the line code, i.e., a short name

direction the direction of the line

stop-order where the stop is in the line

stop-id the id of the stop

frequency the frequency of the line

departure_time the departure time at this stop

arrival_time the arrival time at this stop

line-freq-repetition the repetition of the line

8.6 Tariff Planning

8.6.1 Price Matrix

The file filename_tariff_price_matrix_file (tariff/Price-Matrix.taf) contains the (newly calculated) prices with respect to the specified model (taf_model) for all OD pairs. The columns of the csv file correspond to:

origin-id id of the origin of the OD pair,

destination-id id of the destination of the OD pair,

price price when travelling from the origin to the destination.

8.6.2 Zones

The file `filename_tariff_zone_file` (`tariff/Zones.taf`) contains the assignment of stops to their zones within the zone model `taf_model` zone. The columns of the csv file correspond to:

zone-id id of a zone,

stop-id id of the stop belonging to that zone.

8.6.3 Zone Prices

The file `filename_tariff_zone_price_file` (`tariff/Zone-Prices.taf`) contains the prices in zone model `taf_model` zone for traversing a certain number of zones. The price for traversing more zones than the maximum number of zones specified is just the price for traversing the maximum number of zones specified. The columns of the csv file correspond to:

n-traversed-zones number of traversed zones,

price price for traversing a given number of zones, i.e the price list.

8.7 Vehicle Scheduling

The folder `vehicle-scheduling` contains information about the vehicle tours in the dataset.

8.7.1 Vehicle schedules

The file `default_vehicle_schedule_file` (`vehicle-scheduling/Vehicle_Schedules.vs`) contains information regarding the scheduling of the vehicles. The columns of the csv file correspond to:

circulation-ID Id of the corresponding circulation

vehicle-ID Id of the vehicle

trip-number of this vehicle the trip number of the vehicle

type the type of the tour, can be `trip` for a line serving or `empty` for an empty trip

aperiodic-start-ID the aperiodic event id of the start event of this serving of the line

periodic-start-ID the periodic event id of the start event of this serving of the line

start-stop-id the stop id of the start of the line

start-time the starting time of this service of the line

aperiodic-end-ID the aperiodic event id of the end event of this serving of the line

periodic-end-ID the periodic event id of the end event of this serving of the line

end-stop-id the stop id of the end of the line

end-time the ending time of this service of the line

line the line id

8.8 Delay Management

The folder `delay-management` contains information about the aperiodic event-activity-network, timetable and delays with a disposition timetable

8.8.1 Events expanded

The file default_events_expanded_file (delay-management/Events-expanded.giv) contains information about events in the aperiodic EAN. The columns of the csv file correspond to:

event-id id of the event

periodic-id the corresponding periodic id

type type of the event, can be **departure** for events which are departures of a line at a stop or **arrival** for events which are arrivals of a line at a stop

time the time of the event

passengers number of passengers boarding/alighting at the event

stop-id id of the corresponding stop

8.8.2 Activities expanded

The file default_activities_expanded_file (delay-management/Activities-expanded.giv) contains information about activities in the aperiodic EAN. The columns of the csv file correspond to:

activity-id id of the activity

periodic-id the corresponding periodic id

type type of the activity, can be **drive** for drive activities, **wait** for wait activities, **change** for transfers of passengers, **sync** for synchronization activities between different servings of a line with frequency greater than one or **turnaround** for turnaround activities, i.e., activities of vehicles serving one line after another

tail-event-id id of source event, i.e., the start of the activity

head-event-id id of target event, i.e., the end of the activity

lower-bound the minimal time for this activity, i.e., the minimal time duration needed between the corresponding source and target event to be feasible

upper-bound the maximal time for this activity, i.e., the maximal time duration allowed between the corresponding source and target event to be feasible

passengers the number of passengers using this activity

8.8.3 Timetable expanded

The file default_timetable_expanded_file (delay-management/Timetable-expanded.tim) contains information about the aperiodic timetable, i.e., the time for each aperiodic event. The columns of the csv file correspond to:

event-id id of the event

time the time of the event

8.8.4 Timetable disposition

The file default_disposition_timetable_file (delay-management/Timetable-disposition.tim) contains information about the disposition timetable, i.e., the time for each aperiodic event in the given delay scenario. The columns of the csv file correspond to:

event-id id of the event

time the time of the event

8.8.5 Delays events

The file default_event_delays_file (delay-management/Delays-Events.giv) contains information about the delay induced at the events. The columns of the csv file correspond to:

ID the id of the delayed event

delay the delay, given in seconds

8.8.6 Delays activities

The file default_activity_delays_file (delay-management/Delays-Activities.giv) contains information about the delay induced at the activities. The columns of the csv file correspond to:

ID the id of the delayed activity

delay the delay, given in seconds

8.8.7 Trips

The file default_trips_file (delay-management/Trips.giv) contains information regarding the vehicle trips. A vehicle trips is the serving of a line by a vehicle, i.e., this file contains all line servings in the aperiodic EAN. The columns of the csv file correspond to:

aperiodic-start-ID the aperiodic event id of the start event of this serving of the line

periodic-start-ID the periodic event id of the start event of this serving of the line

start-stop-id the stop id of the start of the line

start-time the starting time of this service of the line

aperiodic-end-ID the aperiodic event id of the end event of this serving of the line

periodic-end-ID the periodic event id of the end event of this serving of the line

end-stop-id the stop id of the end of the line

end-time the ending time of this service of the line

line the line id

8.9 GTFS

Using

```
R make gtfs
```

will create all required gtfs files. For this, the stops (`CK default_stops_file (Fi basis/Stop.giv)`), the line concept (`CK default_lines_file (Fi line-planning/Line-Concept.lin)`), the aperiodic ean (`CK default_events_expanded_file (Fi delay-management/Events-expanded.giv)`), `CK default_activities_expanded_file (Fi delay-management/Activities-expanded.giv)`) and the trips (`CK default_trips_file (Fi delay-management/Trips.giv)`) will be read and the corresponding raw gtfs files will be written to `CK gtfs_output_path (Fi gtfs)`, i.e. the files

- `Fi agency.txt`,
- `Fi stops.txt`,
- `Fi routes.txt`,
- `Fi trips.txt`,
- `Fi stop_times.txt` and
- `Fi calendar.txt`.

Additionally, a zipped file containing the raw data will be created in `CK gtfs_output_path (Fi gtfs)`, named after `CK ptn_name`.

Chapter 9

Datasets

LINTIM provides many datasets to test and evaluate public transport planning algorithms. The following chapter should give an overview over the available datasets and the compatibility with the different planning steps.

9.1 Configuration Parameters for Datasets

There are some configuration parameters used per dataset and not per algorithm. These are set in the file `Fi basis/Config.cnf`.

- `CK gen_conversion_length`: conversion factor used to convert the edge length given in `CK default_edges_file (Fi basis/Edge.giv)` to kilometers.
- `CK gen_conversion_coordinates`: conversion factor used to convert the distance between two stations given in `CK default_stops_file (Fi basis/Stop.giv)` by the coordinates to kilometers.
- `CK gen_vehicle_speed`: speed of the vehicles in km/h.
- `CK ptn_name`: the name of the network
- `CK ptn_stop_waiting_time`: the time each vehicle has to stop at each stop in average. Used in shortest path computation during OD creation.
- `CK period_length`: the length of a period in time units
- `CK time_units_per_minutes`: the number of time units per minute
- `CK ean_default_minimal_waiting_time`: the lower bound for wait activities in the ean. Used during the creation of the ean.
- `CK ean_default_maximal_waiting_time`: the upper bound for wait activities in the ean. Used during the creation of the ean.
- `CK ean_default_minimal_change_time`: the lower bound for change activities in the ean. Used during the creation of the ean.
- `CK ean_default_maximal_change_time`: the upper bound for change activities in the ean. Used during the creation of the ean.
- `CK ean_change_penalty`: the penalty for using a change activity in the ean. Used for routing passengers in the ean and evaluating the perceived travel time.
- `CK gen_passengers_per_vehicle`: the maximal number of passengers per vehicle. Used in computing lower frequency bounds in preparation of line planning.

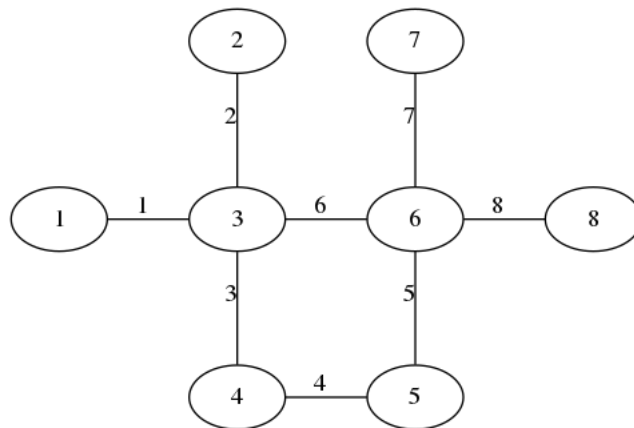


Figure 9.1: The PTN of the toy dataset

9.2 Artificial Datasets

There are two purely artificial datasets in `LINTIM`. These are small examples to test and understand new algorithms.

9.2.1 Toy

The toy dataset is purely designed for testing purposes. It contains 8 nodes, 8 edges and 22 OD pairs, consisting of 2622 passengers in total. An overview of the structure is given in Fig. 9.1.

Since the dataset does not contain the necessary information, stop location is not supported on this dataset out of the box.

9.2.2 Grid

The grid dataset is designed to be overseable, yet complex enough to contain complex effects. Therefore, the dataset contains a simple PTN structure but a reasonable demand structure designed by transportation planners, see [7]. It is part of the benchmark datasets found at [6].

The dataset contains 25 nodes, 40 edges and 567 OD pairs, consisting of 2546 passengers in total. An overview of the structure is given in Fig. 9.2. Since the dataset does not contain the necessary information, stop location is not supported on this dataset out of the box.

9.2.3 Ring

The ring dataset is a little bit larger than the grid dataset but still maintains a clear structure. It is part of the benchmark datasets found at [6].

The dataset contains 161 nodes, 320 edges and 25760 OD pairs, consisting of 2766.12 passengers in total. An overview of the structure is given in Fig. 9.3. Since the dataset does not contain the necessary information, stop location is not supported on this dataset out of the box.

9.3 Datasets based on real world data

9.3.1 Sioux Falls

The sioux falls dataset is a dataset often used in practical public transport planning. It was first introduced in [19] and is available at [42]. It is a representation of the city of Sioux Falls, South Dakote, USA. It is part of the benchmark datasets found at [6].

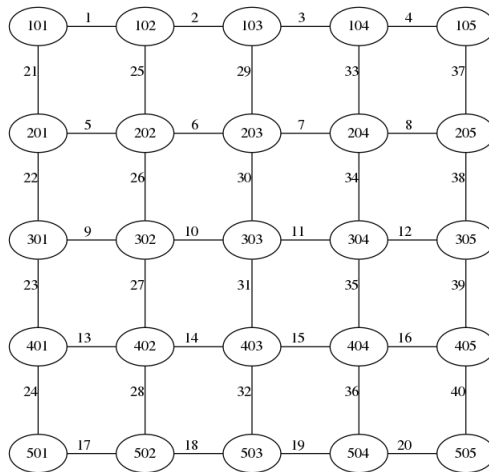


Figure 9.2: The PTN of the grid dataset

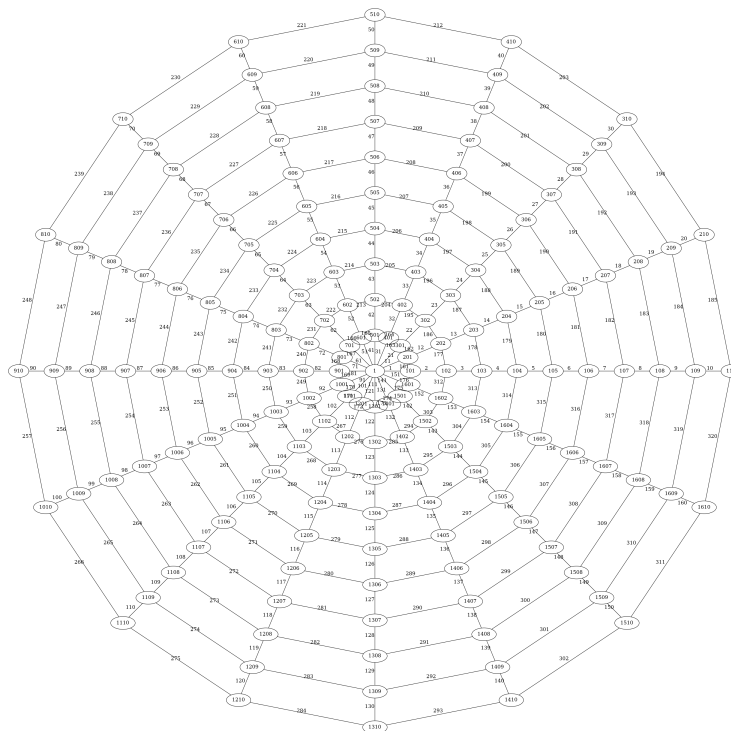


Figure 9.3: The PTN of the ring dataset

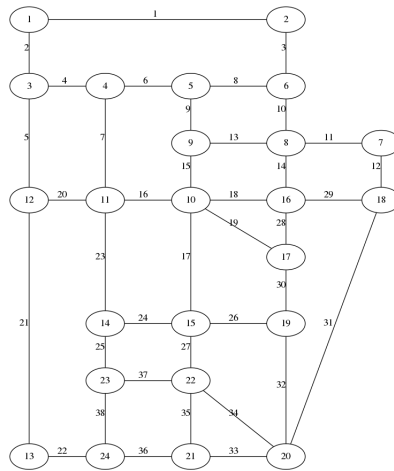


Figure 9.4: Infrastructure of the sioux falls dataset

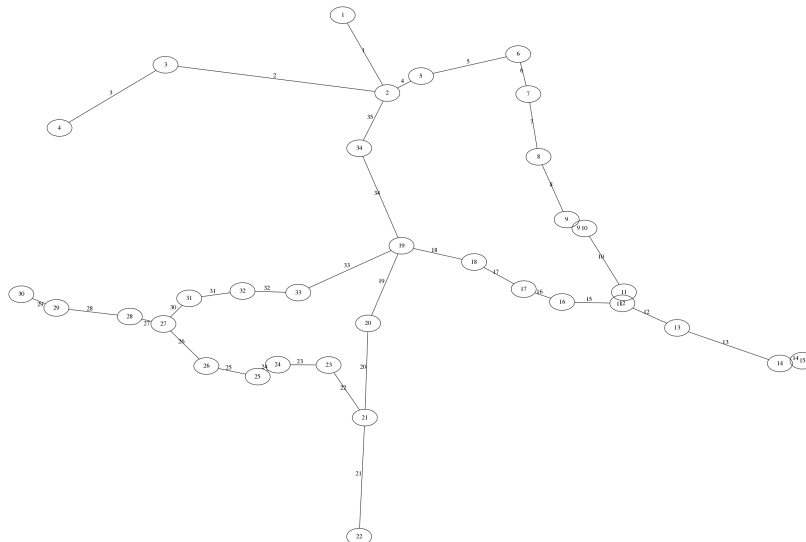


Figure 9.5: Existing infrastructure of the lower saxony dataset

The dataset contains 24 stops, 38 edges and 4114.57 passengers in 552 od pairs. An overview of the structure of the dataset is given in Fig. 9.4.

9.3.2 Lowsaxony

The lower saxony dataset was included to test the effects of stop location and line pool generation. It contains the regional railway data of lower saxony, a region in northern Germany.

The dataset contains 34 existing stops, 35 existing edges and 31 demand points. An overview of the structure given by the existing stops and edges is given in Fig. 9.5. To work with this dataset, you need to start with the stop location step.

9.3.3 Goevb

The goevb dataset represents the bus network in Göttingen, a city in the middle of Germany and home of the LINTIM project. It was included as part of a student project in 2011.

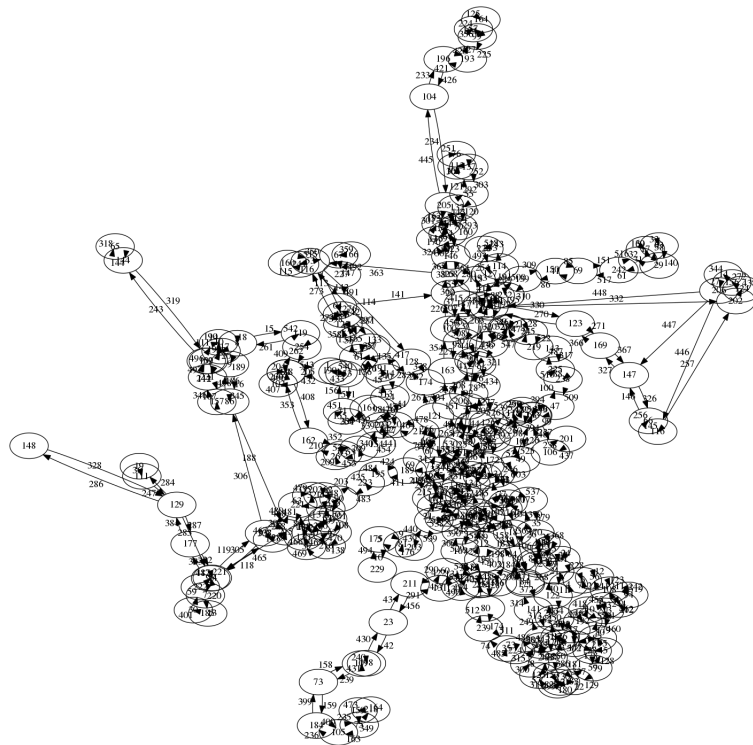


Figure 9.6: The PTN of the goevb dataset

The dataset contains 257 stops, 548 edges and 58226 OD pairs, consisting of 406146 passengers in total. An overview of the structure is given in Fig. 9.6. Since the dataset does not contain the necessary information, stop location is not supported on this dataset out of the box.

Note, that goevb is a directed network!

9.3.4 Athens

The athens dataset represents the metro system in Athens.

The dataset contains 51 stops, 52 edges and 2385 OD pairs, consisting of 63323 passengers in total. An overview of the structure is given in Fig. 9.7. Since the dataset does not contain the necessary information, stop location is not supported on this dataset out of the box.

9.3.5 Bahn-01

Currently not included in the release version of LinTim.

The bahn-01 dataset represents parts of the German railway network, including the long distance network. For larger datasets, see Sec. 9.3.6-9.3.8.

The dataset contains 250 stops, 326 edges and 48842 OD pairs, consisting of 3147382 passengers in total. An overview of the structure is given in Fig. 9.8. Since the dataset does not contain the necessary information, stop location is not supported on this dataset out of the box.

9.3.6 Bahn-02

Currently not included in the release version of LinTim.

The bahn-02 dataset represents parts of the German railway network, including the long distance network. For a smaller dataset see Sec. 9.3.5, for larger datasets, see Sec. 9.3.7 and 9.3.8.

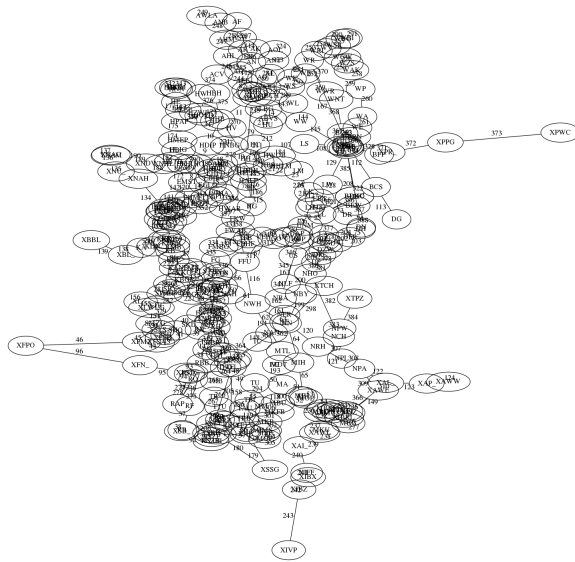


Figure 9.10: The PTN of the bahn-03 dataset

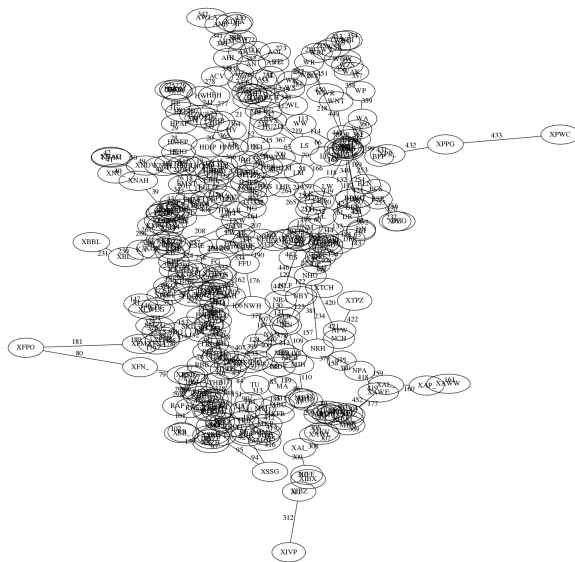


Figure 9.11: The PTN of the bahn-04 dataset

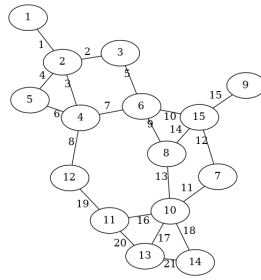


Figure 9.14: The PTN of the Mandl dataset

the local only default parameters in the `[Fi] basis/Config.cnf` file. For an explanation of the parameters, see Section 9.1.

Before running anything, you need to fill the new dataset with data. To see, which algorithm needs which data, see the respective section in this documentation. For information on the file structure, see Chapter 8.

9.4.1 Adding a dataset from PESLib

There is a helper method to import a PESLib dataset. PESLib ([11]) is a benchmark library for Periodic Event Scheduling Problems, based on timetabling problems in public transport planning.

To import a PESLib dataset, place the dataset file (e.g. `R1L1.txt`) into `[Fi]/src/tools/PESLib_import` and run e.g.

```
[R] python3 pesplib_import.py R1L1
```

there. This will create a new dataset folder with the given dataset name and all required files for timetabling in the `[Fi]/datasets`-directory.

9.4.2 Adding a dataset from TimPassLib

There is a helper method to import a TimPassLib dataset. TimPassLib, see [31, 32], is a library for integrated timetabling and passenger routing problems. To import a dataset, create a new dataset based on the template dataset and copy the files `Activities.csv`, `Events.csv`, `OD.csv` and `Config.csv` to

`[CK] filename_timpasslib_activities ([Fi] timpasslib/Activities.csv),`

`[CK] filename_timpasslib_events ([Fi] timpasslib/Events.csv),`

`[CK] filename_timpasslib_od ([Fi] timpasslib/OD.csv)` and

`[CK] filename_timpasslib_config ([Fi] timpasslib/Config.csv),` respectively. Depending on `[CK] timpasslib_import_timetable`, the timetable is imported, i.e., you might need to copy `Timetable.csv` to

`[CK] filename_timpasslib_timetable ([Fi] timpasslib/Timetable.csv).` Run

```
[R] timpasslib-import
```

such that the following files are created:

- `[CK] default_stops_file ([Fi] basis/Stop.giv),`
- `[CK] default_edges_file ([Fi] basis/Edge.giv),`
- `[CK] default_od_file ([Fi] basis/OD.giv),`
- `[CK] default_pool_file ([Fi] basis/Pool.giv),`
- `[CK] default_pool_cost_file ([Fi] basis/Pool-Cost.giv),`
- `[CK] default_lines_file ([Fi] line-planning/Line-Concept.lin),`

- `default_events_periodic_file` (`timetabling/Events-periodic.giv`),
- `default_activities_periodic_file` (`timetabling/Activities-periodic.giv`),
- `default_timetable_periodic_file` (`timetabling/Timetable-periodic.tim`) (if `timpasslib_import_timetable` is set to `true`),

Note that the TimPassLib files contain no information on the coordinates of stops, such that the coordinates (x-coordinate, y-coordinate) in `default_stops_file` will be set to (0,0). Similarly, in `default_edges_file`, the parameter `length` is set to zero and `lower-bound` and `upper-bound` are set according to `lower-bound` and `upper-bound` in `default_activities_periodic_file`.

Additionally, config parameters specified in `filename_timpasslib_config` will be used to update `Config.cnf`. Note that parameters in `filename_timpasslib_config` that are not already present in `Config.cnf` will not be added. These need to be added manually to the corresponding configuration file.

Exporting a dataset in TimPassLib format

A dataset can be exported in TimPassLib format by running

```
 timpasslib-import
```

such that the following files are created according to the specification of [31]:

- `filename_timpasslib_activities` (`timpasslib/Activities.csv`),
- `filename_timpasslib_events` (`timpasslib/Events.csv`),
- `filename_timpasslib_od` (`timpasslib/OD.csv`),
- `filename_timpasslib_config` (`timpasslib/Config.csv`) and
- `filename_timpasslib_timetable` (`timpasslib/Timetable.csv`) (depending on `timpasslib_export_timetable`).

Additionally, basic evaluations of the data set are computed and added to `default_statistic_file` (`statistic/statistic.sta`):

n_stations: The number of stations.

n_lines: The number of operated lines. Note that a line can have a frequency higher than one.

od_prop_entries_greater_zero: The number of OD pairs with $d_{st} > 0$.

od_prop_overall_sum: The total number of passengers.

n_events: The number of events in the event-activity network.

n_activities: The total number of activities in the event-activity network.

n_activities_fixed: The number of activities in the event-activity network with $\ell_a = u_a$.

n_activities_free: The number of activities in the event-activity network with $u_a - \ell_a = T - 1$.

n_activities_restricted: The number of activities in the event-activity network with $\ell_a < u_a < \ell_a + T - 1$.

9.4.3 Dataset generator

There is a make command to create new artificial datasets. To use it, navigate into the `Fi` /datasets- directory and run

```
R make dg-generate-dataset
```

This creates a new dataset as new subdirectory with the method specified by `CK` `dg_model`. For a detailed description of the available models see Section 4.1.

Chapter 10

LinTim Core

For allowing easier extensions of `LINTIM`, its core functionality is provided in two languages, namely Python (3) and Java. There is a version for C++ too, but it is deprecated.

In the following the vocabulary of Java is used, but the versions for Python is structured in the same way. The core is organized into several packages, which are briefly explained in the following sections. Note that for continuity all core libraries follow the naming convention for Java for their public API as far as possible. To create a javadoc version of the documentation run

```
R make docs
```

in the folder `Fo` /src/core/java. An HTML version of the documentation can then be found in `Fo` /src/core/java/docs.

10.1 Model

The package `model` consists of interfaces which represent basic concepts and classes which represent the basic objects used in public transport planning.

10.1.1 Interfaces

The following interfaces are given.

Graph with basic graph functionality

Node with basic node functionality

Edge with basic edge functionality, can be directed or undirected

Path with basic path functionality

OD structure to handle OD information

10.1.2 Classes

The following classes are given.

Stop representing a stop in a PTN, implementing `Node`

Link representing a link in a PTN, implementing `Edge`

InfrastructureNode representing a node in an infrastructure network, e.g., a possible stop location or an intersection, implementing `Node`

InfrastructureEdge representing an infrastructure edge between infrastructure nodes, e.g. a street or a track, implementing Edge

WalkingEdge representing a walking path between infrastructure nodes, implementing Edge

DemandPoint representing a demand point, i.e., the demand at a certain location

StationLimit representing an individual station limit for a stop, containing individual bounds on the transfer or waiting times

Line representing a line in the PTN

LinePool representing a line pool

ODPair representing an origin destination pair

PeriodicEvent representing an event in the periodic event activity network

PeriodicActivity representing an activity in the periodic event activity network

PeriodicHeadway representing a headway activity in the periodic event activity network

AperiodicEvent representing an event in the aperiodic event activity network

AperiodicActivity representing an activity in the aperiodic event activity network

AperiodicHeadway representing a headway activity in the aperiodic event activity network

Timetable representation of a timetable

PeriodicTimetable representation of a periodic timetable

Trip representing an aperiodic trip, e.g., a line serving

Routing representing a routing in a PTN or an EAN.

Zone representing a zone, i.e. a subset of the nodes of a PTN.

ZonePrices stores the prices for travelling through a specified number of zones.

PriceMatrix representing a matrix of prices for each node pair in the PTN.

VehicleTour collecting multiple trips to represent the tour of a vehicle throughout the day

Circulation collecting multiple vehicle tours to represent a circulation

RidepoolArea representing a ridepooling area

Ridepool representing a ridepool, i.e. a collection of ridepooling areas

10.1.3 Enumerations

The following enumerations are given.

EventType possible types of events

ActivityType possible types of activities

LineDirection possible direction of a line (FORWARDS, BACKWARDS)

TariffModelType possible types of tariff models.

TariffObjectiveType possible objectives in tariff planning.

TariffWeightType possible weight options in the objective for tariff planning.

TariffZoneCountingType possible counting modes in a zone tariff.

TariffZoneSymmetryOption possible options for symmetry breaking in the optimization of a zone tariff.

TariffRoutingType possible options for generating a routing in tariff planning.

10.1.4 Package `model.impl`

The package `model.impl` in the Java core contains different implementations of the interfaces, which might be useful in different scenarios.

SimpleMapGraph graph implementation based on Java Maps. Most of the times faster than an `ArrayListGraph`.
May not contain multiple nodes/edges with the same index.

ArrayListGraph graph implementation

LinkedListPath path implementation

MapOD OD implementation used for OD matrices with unknown amount of entries. In most cases the fastest.

FullOD OD implementation used for OD matrices with many entries

SparseOD OD implementation used for OD matrices with few entries

10.2 Input and Output

The package `io` contains reader and writer for all classes in `model` as well as the ones in `util` which need them.

10.3 Algorithm

The package `algorithm` contains implementation of algorithms working on `model` classes, which are needed at several places in `LINTM`.

Dijkstra shortest path implementation using Dijkstra's algorithm

10.4 Utility

The package `util` contains utility classes and enumerations.

Config a representation of the config

Statistic a representation of the statistic

Pair representation of a tuple consisting of 2 elements

LogLevel wrapper mapping different Java logging levels to the ones we are using

SolverType enumeration of different solver types

10.5 Solver

The package `solver` contains an abstract solver implementation, used to formulate a model once and switch the used solver easily. Currently only a small subset of all possible features is implemented, aimed towards high performance to avoid unnecessary overhead. For more information, see the corresponding Javadoc or documentation in the python code.

10.6 Exceptions

The following error catalog is used. All exceptions inherit from `LinTimException` such that logging is handled only once.

Input

- input file cannot be found: Error I1: File `<filename>` cannot be found.
- format of input files is wrong: Error I2: File `<filename>` is not formatted correctly: `<x>` columns given, `<y>` needed.
- inconsistency: Error I3: Column `<x>` of file `<filename>` should be of type `<type>` but entry in line `<line number>` is `<entry>`.
- inconsistent numbering: Error I4: Datatype `<data-type>` is not numbered consistently starting from 1, but `<algorithm-name>` needs that.

Output

- output cannot be written: Error O1: File `<filename>` cannot be written.
- no output is produced: Error O2: Algorithm `<algo>` did not terminate correctly, no output will be produced.

Config parameters

- file not found: Error C1: No config file can be found.
- existence: Error C2: Config parameter `<configkey>` does not exist.
- type: Error C3: Config parameter `<configkey>` should be of type `<type>` but is `<configparameter>`.
- file name not given: Error C4: No config file name given.
- invalid value: Error C5: Value(s) of config parameter(s) `<configparameters>` are invalid or incompatible in this context.

Algorithms

- stopping criterion reached: Error A1: Stopping criterion of algorithm `<algo>` reached without finding a feasible/optimal solution.
- infeasible parameter setting: Error A2: Algorithm `<algo>` cannot be run with parameter setting `<configkey>`; `<configparameter>`.
- in Dijkstra, distance was queried before computation: Error A3: Distance to `<node>` was queried before computation
- in Dijkstra, path was queried before computation: Error A4: Path to `<node>` was queried before computation
- in Dijkstra, algo was called with node, that was not in the graph, when the class was initialized: Error A5: Usage of unknown node `<node>`. This may happen, when the graph was altered after initialization

- in Dijkstra, there is an edge with negative length: Error A6: Edge <edge> has negative length <length>. Dijkstra cannot work reliably with negative edge length.
- in Dijkstra, if the network is not connected: Error A7: Node <sourceNode> is not connected to node <targetNode>, but a shortest path was queried. This may happen during computation of a shortest path or when computing all shortest paths starting from a specific node.

Graphs

- multiple nodes with same index: Error G1: Node with id <node id> already exists.
- multiple edges with same index: Error G2: Edge with id <edge id> already exists.
- left or right node of edge does not exist: Error G3: Edge <edge id> is incident to node <node id> but node <node id> does not exist.
- edge between two nodes does not exist: Error G4: Edge between <node id> and <node id> does not exist.

Lines

- link cannot be added to line: Error L1: Link <link id> cannot be added to line <line id>.
- line contains a circle: Error L2: Line <line id> contains a circle.
- line is no path: Error L3: Line <line id> is no path.

Routings

- no path specified between two nodes: Error R1: No path available from <node id> to <node id>
- path in routing is inconsistent: Error R2: The given path from <node id> to <node id> is not consistent.

Data inconsistency

- periodic event to aperiodic event does not exist: Error D1: Periodic event <event id> to aperiodic event <event id> does not exist.
- periodic activity to aperiodic activity does not exist: Error D2: Periodic activity <activity id> to aperiodic activity <activity id> does not exist.
- index not found: Error D3: <Element> with index <index> not found.
- illegal event type: Error D4: <Event type> of event <event id> is no legal event type.
- illegal activity type: Error D5: <Activity type> of activity <activity id> is no legal activity type.
- illegal line direction: Error D6: <Line direction> of event <event id> is no legal line direction.
- number of lines in Pool-Cost-file does not match number of lines in the linepool: Error D7: Read <number of> entries in the line cost file <filename>, but <number of> lines are in the line pool.

- multiple paths given for one node pair in a routing: Error D8: There are multiple paths given from <node id> to <node id>.
- Path in a routing is inconsistent: Error D9: The path from <node id> to <node id> is not valid.
- Routing is incomplete, but complete routing is needed: Error D10: The given routing is incomplete, but a complete routing is needed.
- Stop is assigned to no zone: Error D11: The stop <stop id> is not assigned to any zone.
- Stop is assigned to multiple zones: Error D12: The stop <stop id> is assigned to more than one zone.
- price matrix does not contain a price for all node pairs with distinct origin and destination node: Error D13: There is no price specified from <node id> to <node id>.
- zone price list is inconsistent: Error D14: Zone price file is inconsistent
- Ridepooling area is not strongly connected: Error D15: Ridepool area with id <area id> is not strongly connected
- Ridepooling area is inconsistent: Error D16: Ridepool area with id <area id> can not be created

Solver

- solver not supported: Error S1: Solver <solver name> not supported for algorithm <algo>.
- Gurobi Error: Error S2: Gurobi returned the following error: <exception.toString()>
- Cplex Error: Error S3: Cplex returned the following error: <exception.toString()>
- Cplex Error: Error S4: The solver <solver name> is not yet implemented in the core solver library.
- Attribute not implemented: Error S5: Attribute <attribute name> is not implemented for <solver name> yet.
- Parameter not implemented: Error S6: The parameter <parameter name> is not implemented for <solver name> yet.
- Variable type not implemented: Error S7: The variable type <variable type> is not implemented for <solver name> yet.
- Invalid call: There was an invalid call, e.g., reading variables of an infeasible model. Please check the text for further information. Error S8: <error message>
- Glpk Error: Error S9: Glpk returned the following error: <exception.toString()>
- Solver found no feasible solution: Error S10: Solver found no feasible solution. Check solver output for further information.

Statistic

- type mismatch: Error ST1: Statistic key <key> should have type <type> but has value <value>.
- key not found: Error ST2: Statistic parameter <configkey> does not exist.

Chapter 11

Introduction to extending LinTim

11.1 Logging

The following guidelines govern the output expected from `LINTIM` programs.

11.1.1 Output from LinTim programs

Output from `LINTIM` programs must adhere to the formatting described here.

For software using a `LINTIM` core Library (Java, C++, ...), there are dedicated logging Classes to use for output.

These will default to write to `STDOUT`, and the Makefile invocations shall do so, but they can also be configured otherwise.

Software not using a `LINTIM` library should use `STDOUT` or a commonly used facility for its respective programming environment/language that can be configured for writing to `STDOUT`, so Makefile invocations can do so.

11.1.2 Log levels

The following Levels shall be used:

FATAL for errors that cancel the execution

ERROR for errors that are severe, but do not stop the program

WARN (a.k.a. warning) for messages from the program that need not be a real error, but may be of interest to the user (also hints for probably wrong configuration) because they might want to be cautious about it, as something is probably different from what they might expect

INFO for everything that happens as expected and is of interest to the end user

DEBUG for output that allows to see what's happening under the hood

In the output to `STDOUT` (be it configurable through a library or not), the loglevel must be written in capital letters, preceded by the current system time formatted as `YYYY-MM-DD HH:mm:ss` at the beginning of the line, followed by a colon, a space, and the actual message. (Only) `DEBUG` messages may additionally contain hints to the source code like the classname, source code line, and/or stack traces of Exceptions, etc.. Multi-line messages are allowed for `DEBUG` messages.

11.1.3 Error messages

The messages outlined in the Error catalog (Section 10.6) shall be used literally for their respective `FATAL`, `ERROR` or `WARN` messages. The level depends upon the severity for the respective program.

11.1.4 Info messages

The following INFO and DEBUG messages should be written at the beginning and end of the respective steps. If a step is not present in a particular program, the respective output can be omitted. Any introductory INFO message(s) (e.g. stating the program name and version) or nothing at all can be output at the beginning.

INFO: Begin reading configuration

DEBUG: Parameter <key> set to <value>

INFO: Finished reading configuration

INFO: Begin reading input data

DEBUG: Reading file <path/and/filename> (done automatically by the respective reader)

INFO: Finished reading input data

INFO: Begin execution

further DEBUG and INFO messages as you see fit

INFO: Finished execution

INFO: Begin writing output data

DEBUG: Writing to file <path/and/filename> or Appending to file <path/and/filename> (done automatically by the respective writer)

INFO: Finished writing output data

Whether the setup of a mathematical program for a solver is done during the reading step (maybe on the fly) or as part of the execution step is up to the author. Solvers may produce their own output to report progress. Whenever possible, the output of a solver shall be configured to go into the filename provided by the configuration key `CK solver_output_file`. (which may contain a relative or absolute path). If the key is the empty string or not set at all, solver output shall be printed to STDOUT, but not through the logging facility (or only at the DEBUG level). (Note: Solver output refers to the usual progress report, not to the results, i.e., values of variables in the solution. Still, intermediate or final results may or may not be part of the solver output.)

11.2 Cleaning

Due to the vast number of algorithms in LINTIM, manually cleaning the `Fo` src directory is tedious. Therefore, LINTIM provides an automatic capability to do so by running

```
R make clean-src
```

in a dataset-folder or

```
R make clean
```

in the `Fo` src directory. There are several file types cleaned automatically from all directories in `Fo` src (see `Fi` src/FILES_TO_CLEAN) but you may add additional files as well. To do so, create a file named `Fi` FILES_TO_CLEAN in the source directory of the algorithm and add all files that should be deleted, one per line. Glob patterns, e.g. `Fi` bin/* are supported.

Chapter 12

Continuous Integration

There are some continuous integration tests contained in LINTM. They can be found in the folder `[Fo]/ci`.

12.1 Running the tests

There are two possibilities, running all test cases and running a specific test.

For running all tests, run the script `[Fi]/ci/run_all_tests.sh`. This file will set some basic environment variables for the solvers and run every test separately. Failed tests will output their respective console log and the names of all failed tests will be collected in `[Fi]/ci/failed_tests`. Also, you may need to make sure, that the environment variables for running the necessary solvers are set for your system, see Chapter 1.2.

There is also the possibility to run a single test. For this, use `[Fi]/ci/run_ussingle_test.sh` with the corresponding test name as the first and only parameter.

Additionally, note that the tests are mostly regression tests, designed to find unintended changes on already implemented algorithms. Therefore, the results are based on running specific software versions on specific hardware. They are therefore likely to fail for you. On the other hand, the unit tests should work for every installation of LINTM. You can run them separately with `[Fi]/ci/run_unit_tests.sh`

12.2 Adding test cases

There is the possibility to add your own test cases. A test contains of four things, a list of LINTM commands to run, a dataset to run the commands on, a `[Fi] Private-Config.cnf` for configuration, and an expected statistic result.

To add your own test, copy the content of `[Fo]/ci/template` into a new subdirectory of `[Fo]/ci`. In there, the commands to run and the dataset can be changed by setting the corresponding variables in `[Fi] run.sh`. To add your own configuration parameters, adapt `[Fi] basis/Private-Config.cnf` in your test directory. This file will be copied in the given dataset before running the test commands.

For the expected results, add data into the file `[Fi] expected-statistic.sta` in your test directory. This file will be compared to the statistic file created by the test commands and will determine the success or the failure of the test. For a successful test, all statistic keys in the `[Fi] expected-statistic.sta` need to be contained in the produced statistic file and their values need to coincide. Note that the produced statistic file may contain more data, this will not cause the test to fail.

Every test will create a new version of the corresponding dataset, you may therefore not assume the dataset to differ from the currently committed version.

Chapter 13

Changelog

This section contains a brief changelog of the different versions. Note that the changelog is not complete and does only include the most important features. For a complete list of changes, use the version control system. The version numbers of LINTIM are based on the date of release and are not semantic.

2024.12

Added

- Added new load generation algorithm based on spanners in the PTN.
- Ridepooling: Added three models to generate ride pool, two models to solve the ridepooling and lineplanning problem simultaneously and corresponding Reader/Writer Classes.

Changed

- Updated code to Gurobi version 12. Gurobi versions older than version 11 are not supported anymore.
- For generated ring datasets, the stop-indices are shifted by 1 and now start at 1.

Fixed

- Line pool and line concept graphics are also created when the value of the configuration keys `CK default_pool_graph_file` and `CK default_line_graph_file` differ from the default setting.
- The configuration file given to `Fi src/line-planning/line-planning.sh` is passed over to the called java programs.
- Fixed line concept evaluation using change model `FORMULA_2` or `FORMULA_3`.

2024.08

Added

- Added functionality for tariff planning including optimizing and evaluating fares.
- New method to generate a line pool based on centers and periphery nodes.
- Added IP model for line planning with passenger routing to minimize the total (estimated) traveling time.
- Added Mandl dataset.

Fixed

- Line pool creators also add lines starting with an edge specified in backwards direction (of the line direction).
- Timetable evaluation parameter `SK tim_overcrowded_time_average` now also considers overcrowded wait activities.
- Correctly compute wait and transfer time in line concept evaluation using change-and-go graph.
- Specified change stations are taken into account when constructing the EAN.
- Upper frequency bounds computed via `CK load_generator_upper_frequency_factor` are rounded up.
- The terminal-to-terminal model for line pool generation uses the standard line cost computation method including the `CK lpool_costs_vehicles`.
- Line reader now uses `CK gen_conversion_length`.

2023.12

Added

- Functionality for creating artificial datasets, see Section 4.1.
- Import and export to TimPassLib format, see Section 9.4.2.
- Multi-commodity flow routing in order to evaluate a timetable, taking into account the capacities, see Section 5.7.1.
- Possibility to draw PTN without stop-coordinates.

Changed

- Allow non-integer activity weights for all activity buffer models.
- If OD pair is not present, it is interpreted as 0 in the python core.

Fixed

- Key feature computation for ml-robustness framework will now correctly respect the chosen routing window.
- Extended evaluation of line concept solve a multi-commodity flow problem instead of independent source problems for all origins.
- Timetable evaluation can now handle non-integer passengers.
- Fixed integer overflow in c++ core when writing large integers.
- Fixed bug in timetable evaluation that allowed routing of passengers on sync and headway activities.
- Periodic timetabling with constraint propagation terminates with an exception if activity file is missing.
- An error when reading a CSV file with the Java core is no longer always indicated by the message “File not found”. Message “Wrong encoding” was added.
- Fixed candidate set for stop location; will now correctly handle cases where demand points are exactly `CK sl_radius` distanced from possible stop locations.

- Added `CK` `gen_conversion_length` and `CK` `gen_conversion_coordinates` to core writers, so that a the values are not changed when a graph is only read and written again.
- Use `CK` `ean_change_penalty` under all settings of the configuration parameter `CK` `ean_model_weight_change`.

2022.08

Added

- Possibility to read stop geo coordinates in Java and Python core libraries
- Interactive visualization of a PTN, see Section 4.11.1
- Visualization of the OD data via a graph or a heatmap, see Section 4.11.2
- Possibility to visualize the ptn load weights, see Section 4.11.3
- Installation script for installing dependencies automatically, see Section 1.3

Changed

- Remove Station-Distance requirement of `R` `make vs-add-circulations-to-ean`
- Updated JUnit-version from 4.12 to 4.13.2
- Update Java core cplex interface to CPLEX 20.1
- Now most config parameters are case insensitive
- Rewrite line concept evaluation for better performance and more evaluations. Note that the names of some statistic entries changed to be more clear. For all current evaluations, see Section 5.5.

Fixed

- Python Core: Prevent overwriting statistic when trying to append
- Python Core: Correctly parse the entries in stop-possible? for infrastructure nodes
- Python & Java Core: Dijkstra will now return copies of the computed paths to prevent accidental changes by the user.
- Fix solver core interface of the stop location travel time model
- Fix solver dependency of extended line planning cost model
- Line planning direct model can now handle non-zero diagonal entries in the od matrix

2021.12

Added

- Added more integer programming solver support. For an overview which solvers are support by which algorithms, see Section 6. For more information on how to combine solvers with `LNTIM`, see Section 1.2.1.
- Robust integrated planning based on machine learning predictions. For more information, see Section 3.10.5.
- Possibility to run `LNTIM` an ARM-based cpus, e.g. Apple-M1

Fixed

- Add java core dependency installation for terminal-to-terminal line pool generation
- Fix wrong make target for ean passenger reroute
- Fix missing build files for line pool drawing
- Line pool cost computation will now scale the ptn edges according to `CK` `gen_conversion_length`
- Will now read `CK` `ptn_stop_waiting_time` for the ptn evaluation
- The vehicle-based term of the line costs now accounts for undirected lines as well

Removed

- Possibility to run `LINTIM` on i586 cpus.

2021.10

Added

- Ability to respect additional load per link in load generation, see Section 7.7
- Export to GTFS, see Section 8.9
- Cycle base formulation for periodic timetabling, see Section 3.6.6
- Phase 1 simplex for periodic timetabling, see Section 3.6.7
- Visum-Interface to import datasets from PTV Visum, see Section 4.12. This includes several additions to `LINTIM`:
 - An infrastructure model, more detailed than the current PTN representations, see e.g. Section 3.1.2
 - Possibility of passengers to walk, see e.g. Section 4.5 and Section 4.2.4
 - Respecting transfer stations and line terminals, see e.g. Section 4.5 and Section 3.2.1
 - Forbidding edges in line planning, see e.g. Section 3.3.1

Changed

- Bump used JGraphT version, now JGraphT 1.5 and JHeaps 0.13 are required
- Java 11 is now required
- Maven (≥ 4) is now required
- Rewrite several ip models, using a common naming scheme for solver parameters and align the output of the programs to the rest of `LINTIM`

Fixed

- The rollout step will not read the headways anymore if they are not needed
- Python Core now reads directed ptns correctly
- DM extended evaluation now computes average values correctly
- Rolling out passenger paths now works on aperiodic eans without changes

- PTN load generator will now compute correct variable upper frequency bounds for very small load values
- Rolling out passenger paths does not allow headways in passenger paths anymore
- Fixed Big-M-value for DM1

2020.12

Added

- Additional IP parameters for Gurobi
- Dataset ring

Changed

- Python Core: Replaced usage of DictGraph by SimpleDictGraph to improve performance
- Core: StatisticWriter will default to appending to the file on disc instead of overwriting
- Line planning model direct is now allowed a non-integer budget restriction
- Remove goblin dependency from periodic modulo simplex, use gurobi now instead
- Allow periodic timetable evaluation without an od matrix present

Fixed

- make ean-add-simple-vs will now respect the parameter time_units_per_minute
- Line Planning method cost_restricting_frequencies can now be compiled with only one of the supported solvers installed
- Python core will use default statistic for reading if none is given
- Fixed bug in cycle base version of integrated timetabling and passenger routing model
- Adapted ean_change_penalty for time_units_per_minute in dataset athens
- Equals method in periodic and aperiodic ean now working in python core
- Suppress double logging/console output when using the core gurobi solver interface with gurobi 9
- Python core vehicle schedule writer reads correct default config key for the vehicle schedule file
- make ean-add-simple-vs will now throw an error when run on a directed ptn
- time_units_per_minute are now consistently handled in all vehicle scheduling methods

2020.02

Added

- Sioux Falls dataset
- Models for integrated planning
 - Integrated timetabling and passenger routing
 - Integrated line planning, timetabling and passenger routing
 - Integrated timetabling and vehicle scheduling

- Integrated line planning, timetabling, passenger routing and vehicle scheduling
- Computing a new timetable for given line plan and vehicle schedule
- Respect fixed lines in line planning
- Respect fixed lines in timetabling
- Modulo Simplex algorithm for timetabling
- Full support for running under Windows
- Import of VISUM datasets
- New Python core graph implementation
- Automatic cleaning of src folders
- Robustness checks for delay management

Changed

- The export format to visum does now include the line repetition

Deprecated

- the cpp core will not be maintained any more and will be removed in a future version

2018.06

First release version

Bibliography

- [1] S. H. Bull, R. M. Lusby, and J. Larsen, *An optimization based method for line planning to minimize travel time*, Proceedings of the 13th conference on advanced systems in public transport (CASPT), 2015.
- [2] S. Bunte and N. Kliewer, *An overview on vehicle scheduling models*, Public Transport **1** 2009, no. 4, 299–317.
- [3] M. Bussieck, *Optimal lines in public rail transport*, Ph.D. Thesis, 1998.
- [4] E. Carrizosa, J. Harbering, and A. Schöbel, *Minimizing the passengers’ traveling time in the stop location problem*, Journal of the Operational Research Society **67** 2016, no. 10, 1325–1337.
- [5] A. Fielbaum, S. Jara-Diaz, and A. Gschwender, *A parametric description of cities for the normative analysis of transport systems*, Networks and Spatial Economics **17** 2017, no. 2, 343–365.
- [6] *Collection of open source public transport networks by DFG Research Unit “FOR 2083: Integrated Planning For Public Transportation”*, 2018. <https://github.com/FOR2083/PublicTransportNetworks>.
- [7] M. Friedrich, M. Hartl, A. Schiewe, and A. Schöbel, *Angebotsplanung im öffentlichen Verkehr - Planerische und algorithmische Lösungen*, Heureka, 2017.
- [8] ———, *Integrating passengers’ assignment in cost-optimal line planning*, 17th workshop on algorithmic approaches for transportation modelling, optimization, and systems (ATMOS 2017), 2017, pp. 1–16.
- [9] ———, *System headways in line planning*, CASPT 2018, 2018.
- [10] P. Gattermann, J. Harbering, and A. Schöbel, *Line pool generation*, Public Transport **9** 2017, 7–32.
- [11] M. Goerigk, *PESPlib*. <https://timpasslib.net/pesplib.html>.
- [12] ———, *Verallgemeinerte Schnittheuristiken in der periodischen Fahrplangestaltung*, Master’s Thesis, 2009.
- [13] M. Goerigk and A. Schöbel, *Improving the modulo simplex algorithm for large-scale periodic timetabling*, Computers & Operations Research **40** 2013, no. 5, 1363–1370.
- [14] M. Goerigk, A. Schöbel, and F. Spühler, *A phase I simplex method for finding feasible periodic timetables*, 21st symposium on algorithmic approaches for transportation modelling, optimization, and systems (ATMOS 2021), 2021, pp. 6:1–6:13.
- [15] J. Harbering, *Delay resistant line planning with a view towards passenger transfers*, TOP **25** 2017, 467–496.
- [16] I. Heinrich, O. Herrala, P. Schiewe, and T. Terho, *Using light spanning graphs for passenger assignment in public transport*, 23rd symposium on algorithmic approaches for transportation modelling, optimization, and systems (atmos 2023), 2023, pp. 2:1–2:16.
- [17] A. Kaufmann, *Column generation for line planning with minimal traveling time*, Master’s Thesis, 2016.
- [18] M. Lachmann, *Vehicle scheduling based on a line plan only*, Master’s Thesis, 2016.
- [19] L. J. LeBlanc, E. K. Morlok, and W. P. Pierskalla, *An efficient approach to solving the road network equilibrium traffic assignment problem*, Transportation Research **9** 1975, no. 5, 309–318.
- [20] C. Mandl, *Applied network optimization*, Academic Press, New York, 1979.
- [21] M. Müller-Hannemann, R. Rückert, A. Schiewe, and A. Schöbel, *Framework for generating machine learning models for robustness*, 2021. available at <https://gitlab.rlp.net/for2083/framework-for-generating-machine-learning-models-for-robustness>.
- [22] M. Müller-Hannemann, R. Rückert, A. Schiewe, and A. Schöbel, *Towards Improved Robustness of Public Transport by a Machine-Learned Oracle*, 21st symposium on algorithmic approaches for transportation modelling, optimization, and systems (atmos 2021), 2021, pp. 3:1–3:20.
- [23] C. L. Mumford, *Research on the urban transit routing problem (bus routing)*, 2016. <https://users.cs.cf.ac.uk/C.L.Mumford/Research%20Topics/UTRP/Outline.html#Xmumford2013>.
- [24] J. Pätzold, A. Schiewe, P. Schiewe, and A. Schöbel, *Look-ahead approaches for integrated planning in public transportation*, 17th workshop on algorithmic approaches for transportation modelling, optimization, and systems (ATMOS 2017), 2017, pp. 17:1–17:16.
- [25] J. Pätzold and A. Schöbel, *A matching approach for periodic timetabling*, 16th workshop on algorithmic approaches for transportation modelling, optimization, and systems (ATMOS 2016), 2016, pp. 1:1–1:15.
- [26] PTV AG, *Visum 17 user manual*, 2018.

- [27] M. Schachtebeck, *Delay management in public transportation: Capacities, robustness, and integration*, Ph.D. Thesis, 2010.
- [28] M. Schachtebeck and A. Schöbel, *To wait or not to wait—and who goes first? Delay management with priority decisions*, *Transportation Science* **44** 2010, no. 3, 307–321.
- [29] A. Schiewe and P. Schiewe, *An iterative approach for integrated planning in public transportation*, Georg-August-Universität Göttingen, 2018. Working Paper.
- [30] P. Schiewe, *Integrated optimization in public transport planning*, Ph.D. Thesis, 2018.
- [31] P. Schiewe, M. Goerigk, and N. Lindner, *TimPassLib*. <https://timpasslib.net/>.
- [32] ———, *Introducing TimPassLib – a library for integrated periodic timetabling and passenger routing*, *Operations Research Forum* **4** 2023, no. 3.
- [33] A. Schöbel, *Optimization models in public transportation*, 2004.
- [34] ———, *Optimization in public transportation. stop location, delay management and tariff planning from a customer-oriented point of view*, *Optimization and Its Applications*, Springer, New York, 2006.
- [35] ———, *Integer programming approaches for solving the delay management problem*, *Algorithmic methods for railway optimization*, 2007, pp. 145–170.
- [36] ———, *Line planning in public transportation: models and methods*, *OR Spectrum* **34** 2012, no. 3, 491–510.
- [37] ———, *An eigenmodel for iterative line planning, timetabling and vehicle scheduling in public transportation*, *Transportation Research C* **74** 2017, 348–365.
- [38] A. Schöbel, H. W. Hamacher, A. Liebers, and D. Wagner, *The continuous stop location problem in public transportation*, *Asia-Pacific Journal of Operational Research* **26** 2009, no. 1, 13–30.
- [39] A. Schöbel and S. Scholl, *Line planning with minimal traveling time*, 5th workshop on algorithmic methods and models for optimization of railways, 2006.
- [40] A. Schöbel and S. Schwarze, *Finding delay-resistant line concepts using a game-theoretic approach*, *Netnomics* **14** 2013, no. 3, 95–117.
- [41] P. Serafini and W. Ukovich, *A mathematical model for periodic scheduling problems*, *SIAM Journal on Discrete Mathematics* **2** 1989, no. 4, 550–581.
- [42] B. Stabler, *Sioux falls - github*, 2018. available at <https://github.com/bstabler/TransportationNetworks/tree/master/SiouxFalls>.
- [43] A. Uffmann, *Umlaufplanung mit dem Kanalmodell*, Master’s Thesis, 2010.